



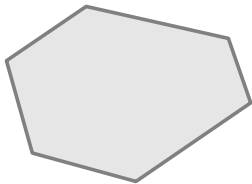
GEOMETRIC EXPANSION PROPERTIES OF POLYTOPE SKELETA

Martin Winter

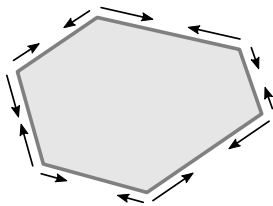
University of Warwick

20. March, 2022

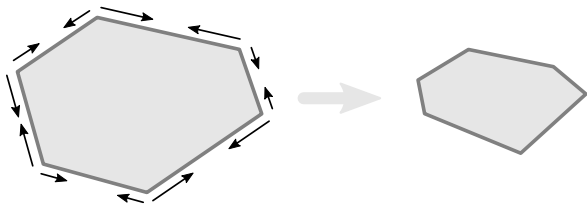
AN INTUITIVE PROBLEM



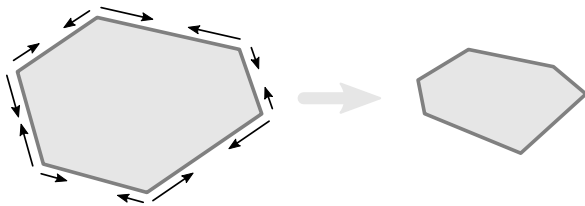
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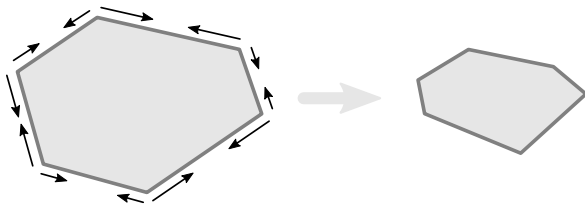


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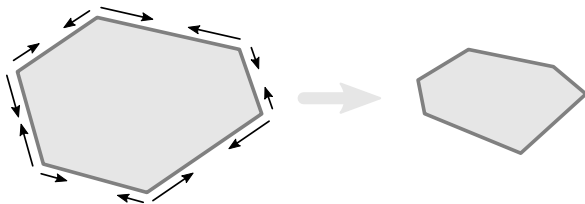
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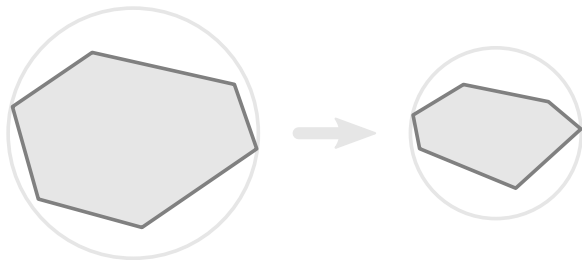
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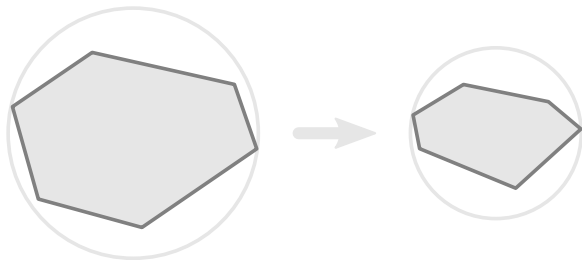
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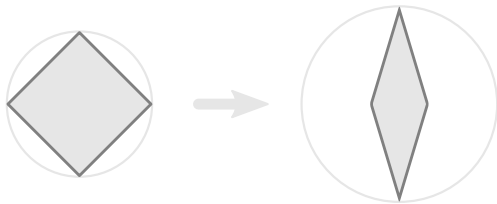


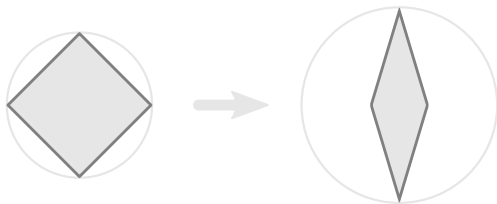
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- ▶ volume

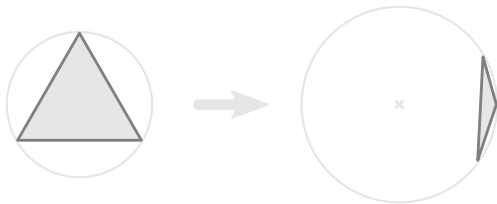


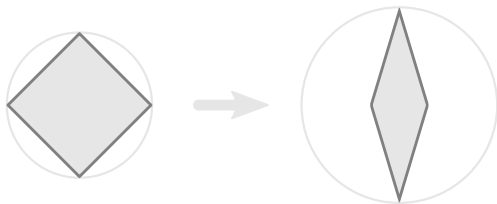


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Idea: what if the circumcenter lies inside the polytope?

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then the circumradius of Q is at most as large as the circumradius of P .

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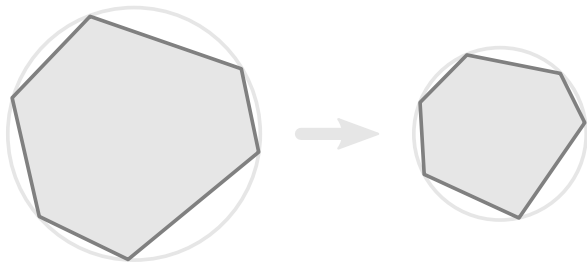
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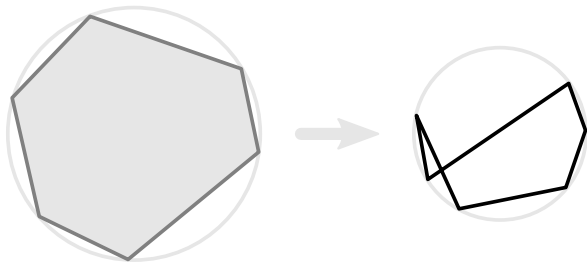
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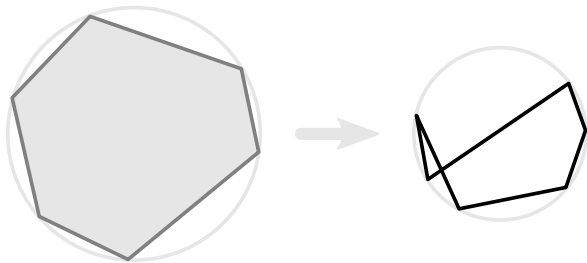
WHY STOP THERE?



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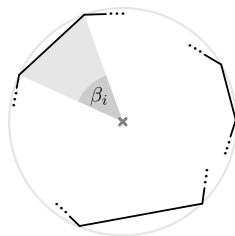
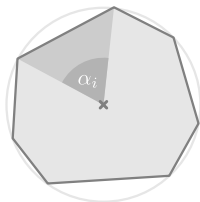


graph embedding $w: V \rightarrow \mathbb{R}^d$.

CASE $d = 2$: POLYGON SKELETONS

Proof sketch.

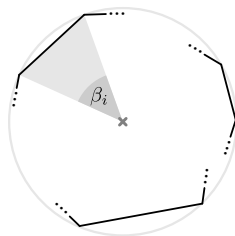
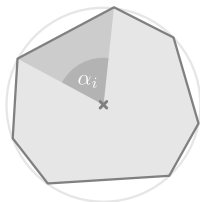
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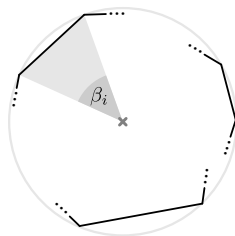
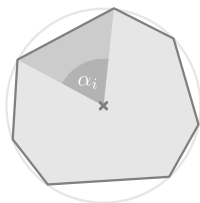


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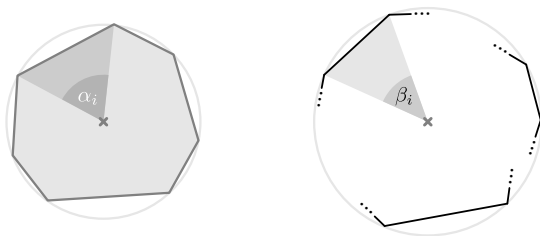
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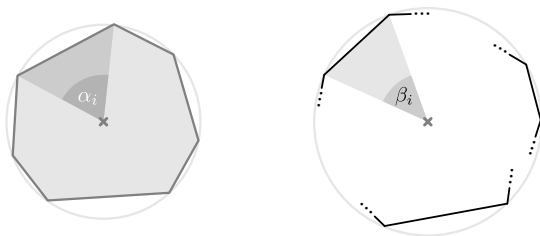
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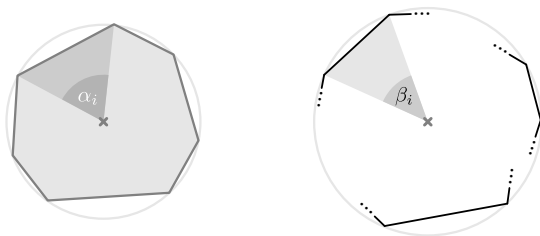
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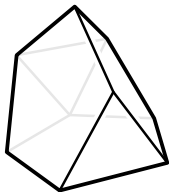
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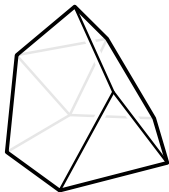
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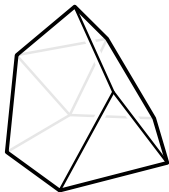
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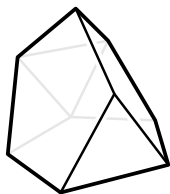
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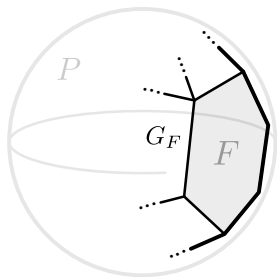
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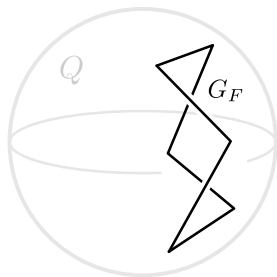
Yes, for matroid base polytopes



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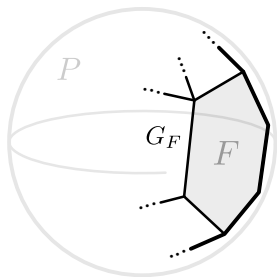


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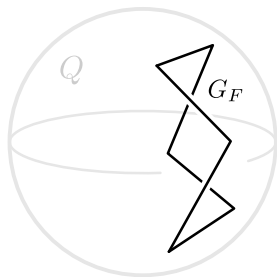


G_F "weirdly" embedded in Q

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The embedding of G_F in Q must be "less expanded" than F .

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We want to show: polytope edges are enough!

ANOTHER MEASURE FOR EXPANSION

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Definition.

An m -centered $v: V \rightarrow \mathbb{R}^d$ is m -**expanded** if for every m -centered $w: V \rightarrow \mathbb{R}^f$ with

$$\|w_i - w_j\| \leq \|v_i - v_j\|, \quad \text{for all } ij \in E$$

holds $\|w\|_m \leq \|v\|_m$.

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Example: (negated) Colin de Verdière matrices.

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Two facts about promising embeddings:

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If $0 \in P$ then its skeleton is a promising embedding for some matrix $M \in \mathbb{R}^{n \times n}$.

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- ▶ w has the same edge-lengths as v , and
- ▶ w is a linear transformation of v .

SO, WHAT ABOUT INSCRIBED POLYTOPES?

If P is inscribed with center $0 \in P$, then

$$\|v\|_m^2 = \sum_i m_i \|v_i\|^2 = r(v)^2 \cdot \underbrace{\left(\sum_i m_i \right)}_{\mu}.$$

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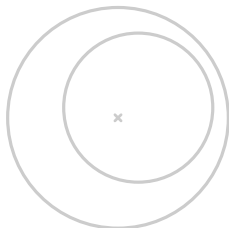
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SO, WHAT ABOUT INSCRIBED POLYTOPES?

Theorem. (W., 2022+)

Given an inscribed polytope $P \in \mathbb{R}^d$ and a centrally symmetric inscribed embedding $w: V \rightarrow \mathbb{R}^f$ so that

- ▶ each edge in w is at most as long as its counterpart in P ,

then $r(w) \leq r(P)$.

Thank you.

