

Nonlinear approximation in bounded orthonormal product bases

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joint work with Lutz Kämmerer and Daniel Potts

Workshop on Mathematical Signal and Image Analysis

21.03.2023



UNIVERSITY OF TECHNOLOGY
IN THE EUROPEAN CAPITAL OF CULTURE
CHEMNITZ

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② The algorithm

Projected coefficients

The dimension-incremental method

③ Theoretical detection guarantee for function approximation

④ Numerical examples

10-dimensional periodic test function

9-dimensional non-periodic test function

⑤ Conclusion

General aim in the Fourier setting

Approximation (by using samples) of a (smooth) high-dimensional function $f \in L_2(\mathbb{T}^d, \mu)$

by using samples

black-box sampling, so we choose the sampling nodes \mathbf{x} adaptively

Hilbert space $L_2(\mathbb{T}^d, \mu)$

domain $\mathbb{T}^d = \times_{j=1}^d \mathbb{T} \subset \mathbb{R}^d$

measure $\mu = \times_{j=1}^d \mu_j$

basis $\Phi_{\mathbf{k}}(\mathbf{x}) = \prod_{j=1}^d e^{2\pi i k_j x_j}$

smooth function f

$$f(\mathbf{x}) := \sum_{\mathbf{k} \in \mathbb{Z}^d} c_{\mathbf{k}} e^{2\pi i \mathbf{k} \cdot \mathbf{x}}$$

Truncation and approximation

$$S_I^A f(\mathbf{x}) := \sum_{\mathbf{k} \in I} \hat{f}_{\mathbf{k}} e^{2\pi i \mathbf{k} \cdot \mathbf{x}}$$

To compute

- ▶ approximated coefficients $\hat{f}_{\mathbf{k}} \approx c_{\mathbf{k}} \forall \mathbf{k} \in I$
- ▶ suitable sparse index set $I \subset \mathbb{Z}^d$

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General aim for an arbitrary bounded orthonormal product basis (BOPB)

Approximation (by using samples) of a (smooth) high-dimensional function $f \in L_2(\mathcal{D}, \mu)$

by using samples

black-box sampling, so we choose the sampling nodes \mathbf{x} adaptively

Hilbert space $L_2(\mathcal{D}, \mu)$

domain $\mathcal{D} = \times_{j=1}^d \mathcal{D}_j \subset \mathbb{R}^d$

measure $\mu = \times_{j=1}^d \mu_j$

basis $\Phi_{\mathbf{k}}(\mathbf{x}) = \prod_{j=1}^d \phi_{j,k_j}(x_j)$

smooth function f

$$f(\mathbf{x}) := \sum_{\mathbf{k} \in \mathbb{N}^d} c_{\mathbf{k}} \Phi_{\mathbf{k}}(\mathbf{x})$$

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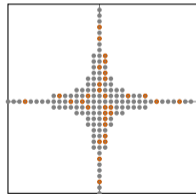
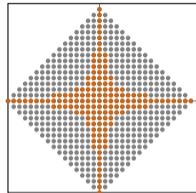
- ▶ approximated coefficients $\hat{f}_{\mathbf{k}} \approx c_{\mathbf{k}} \forall \mathbf{k} \in I$
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Problem: How to find a good, s -sparse index set I ?

- ▶ first idea:
 - ▶ choose a large search space $\Gamma \supset I, |\Gamma| \gg |I|$
 - ▶ compute all coefficients
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- ▶ but: unfeasible in practice for large Γ (\rightarrow “curse of dimensionality”)
- ▶ better idea: use a dimension-incremental approach

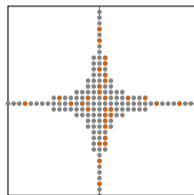
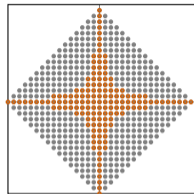
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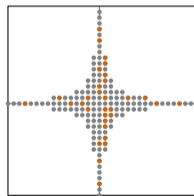
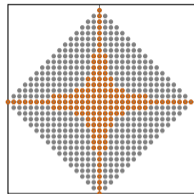
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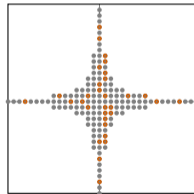
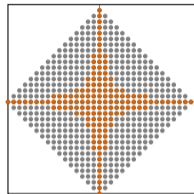
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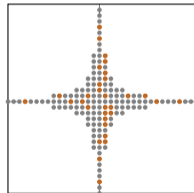
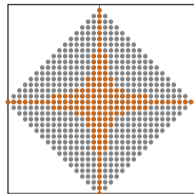
Question: Can we ensure to detect all important indices?



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Question: Can we ensure to detect all important indices? \rightarrow YES!



Previous works on sparse, high-dimensional approximation:

- ▶ our dimension-incremental approach in the Fourier setting:
 - ▶ sparse FFT using rank-1 lattices
 [Potts, Volkmer '16], [Kämmerer, Potts, Volkmer '21], [Kämmerer, Kraher, Volkmer '22]
 - ▶ application to differential equations with high-dimensional random coefficients
 [Bochmann, Kämmerer, Potts '20], [Kämmerer, Potts, T. '22]
- ▶ other approximation methods in the Fourier setting:
 e.g. [Iwen '13], [Indyk, Kapralov '14], [Choi, Christlieb, Wang '19], ...
- ▶ sparse high-dimensional approximation in more general bases:
 - ▶ sparse polynomial chaos expansions
 literature survey: [Lüthen, Marelli, Sudret '21], basis-adaptive: [Lüthen, Marelli, Sudret '21]
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⑤ Conclusion

Projected coefficients for the dimensions $\{1, \dots, t\}$

$$c_{\{1, \dots, t\}, \mathbf{k}}(\tilde{\mathbf{x}}) := \int_{\mathcal{D}_{\{1, \dots, t\}}} f(\boldsymbol{\xi}, \tilde{\mathbf{x}}) \overline{\Phi_{\{1, \dots, t\}, \mathbf{k}}(\boldsymbol{\xi})} d\mu_{\{1, \dots, t\}}(\boldsymbol{\xi}) \quad \forall \mathbf{k} \in \mathbb{N}^t, \forall \tilde{\mathbf{x}} \in \mathcal{D}_{\{t+1, \dots, d\}}$$

Projected coefficients are an indicator for the importance of the indices (\mathbf{k}, \mathbf{h}) with arbitrary \mathbf{h} , since

$$c_{\mathbf{u}, \mathbf{k}}(\tilde{\mathbf{x}}) = \dots = \sum_{\mathbf{h} \in \mathbb{N}^{d-t}} c_{(\mathbf{k}, \mathbf{h})} \Phi_{\{t+1, \dots, d\}, (\mathbf{k}, \mathbf{h})}(\tilde{\mathbf{x}}).$$

Example: Fourier setting with $d = 3$ and $t = 2$

$$c_{\{1, 2\}, \mathbf{k}}(\tilde{\mathbf{x}}) := \int_{\mathbb{T}^2} f(\boldsymbol{\xi}, \tilde{\mathbf{x}}) e^{-2\pi i \mathbf{k} \cdot \boldsymbol{\xi}} d\boldsymbol{\xi} = \sum_{\mathbf{h} \in \mathbb{Z}} c_{(\mathbf{k}, \mathbf{h})} e^{2\pi i \mathbf{h} \tilde{\mathbf{x}}} \quad \forall \mathbf{k} \in \mathbb{Z}^2, \forall \tilde{\mathbf{x}} \in \mathbb{T}$$

→ Works analogously for single dimensions $\{t\}$ and arbitrary $\mathbf{u} \subset \{1, \dots, d\}$.

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Key idea of the dimension-incremental algorithm

- ▶ Construct a candidate set with possible indices in the current dimension(s).
- ▶ Compute approximations of the projected coefficients (using samples).
- ▶ Build the index set:

$$\begin{cases} \text{projected coefficient large :} & \text{keep index} \\ \text{projected coefficient small :} & \text{neglect index} \end{cases}$$
- Increase the dimension by combining different detected index sets of smaller dimension.

Approximation of the projected coefficients via, e.g., cubature formulas (nodes ξ_j , weights w_j)

$$\begin{aligned} c_{\{1,\dots,t\},k}(\tilde{\mathbf{x}}) &:= \int_{\mathcal{D}_{\{1,\dots,t\}}} f(\boldsymbol{\xi}, \tilde{\mathbf{x}}) \overline{\Phi_{\{1,\dots,t\},k}(\boldsymbol{\xi})} d\mu_{\{1,\dots,t\}}(\boldsymbol{\xi}) \\ &\approx \sum_{j=1}^M w_j f(\boldsymbol{\xi}_j, \tilde{\mathbf{x}}) \overline{\Phi_{\{1,\dots,t\},k}(\boldsymbol{\xi}_j)} =: \hat{f}_{\{1,\dots,t\},k}(\tilde{\mathbf{x}}) \end{aligned}$$

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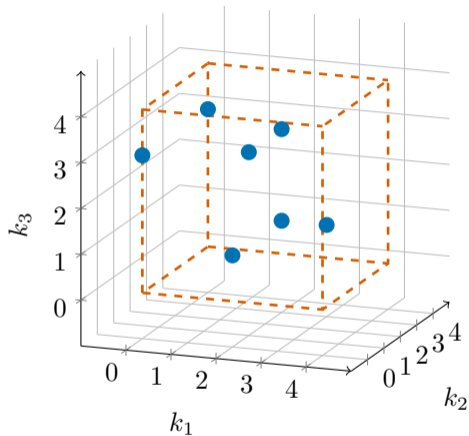
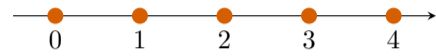
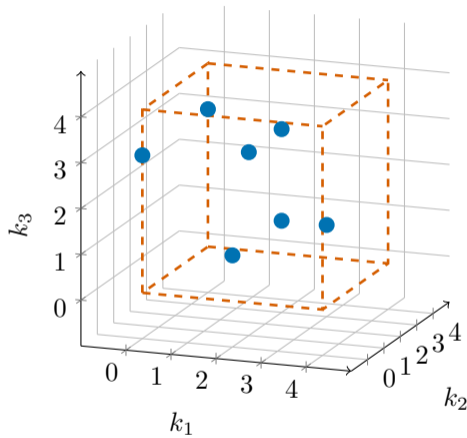
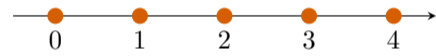


Figure: The desired but unknown index set I and the search space $\Gamma = \{0, 1, 2, 3, 4\}^3$.



$$k_1 \in \mathcal{P}_{\{1\}}(\Gamma)$$



$$k_2 \in \mathcal{P}_{\{2\}}(\Gamma)$$



$$k_3 \in \mathcal{P}_{\{3\}}(\Gamma)$$

Figure: The desired but unknown index set I and the search space $\Gamma = \{0, 1, 2, 3, 4\}^3$.

Figure: The one-dimensional candidate sets $\mathcal{P}_{\{j\}}(\Gamma)$.

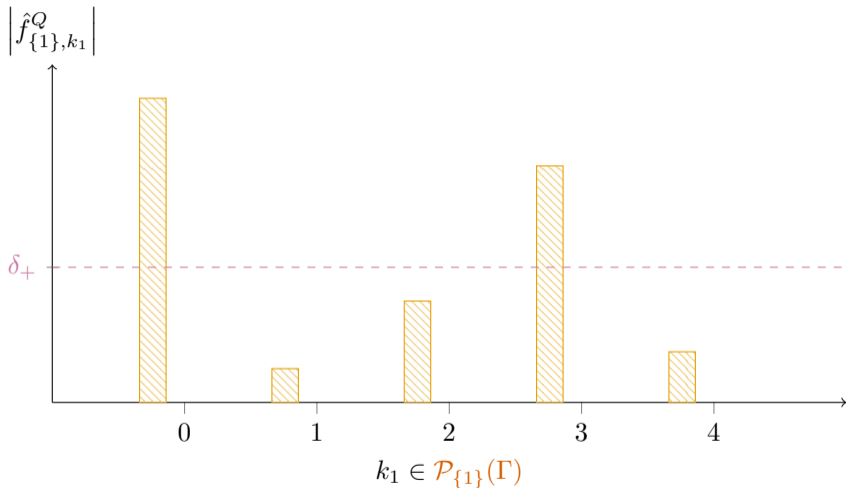


Figure: The one-dimensional detection step in the first dimension with $r = 3$ detection iterations.

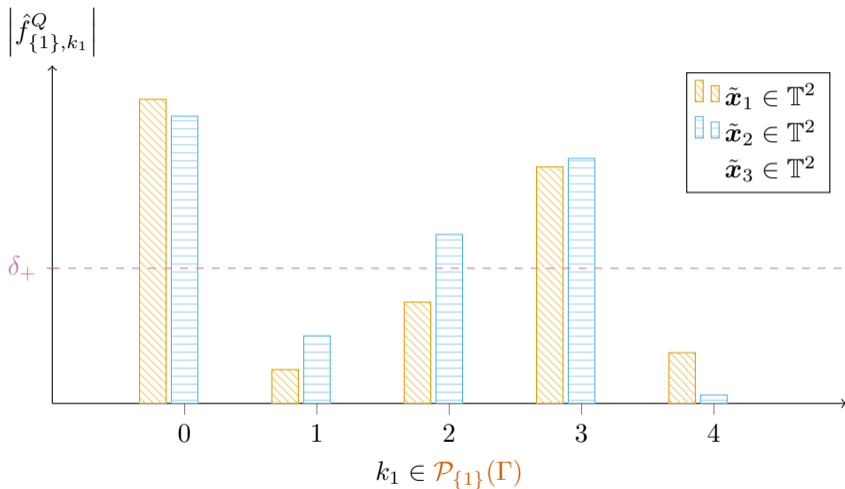


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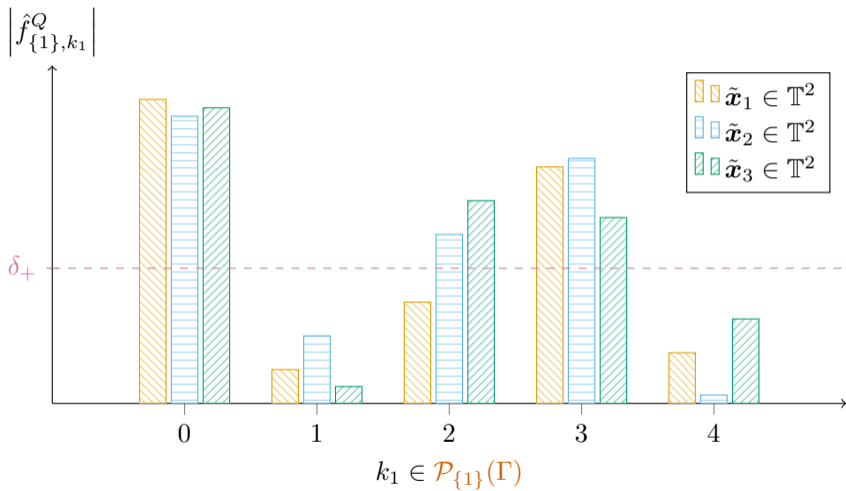


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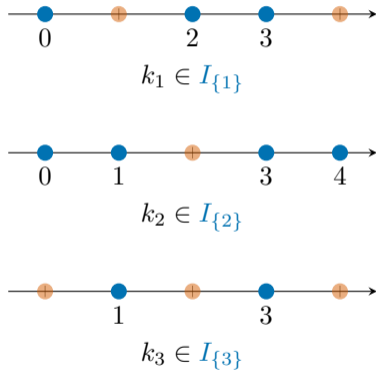


Figure: The one-dimensional index sets $I_{\{j\}}$.

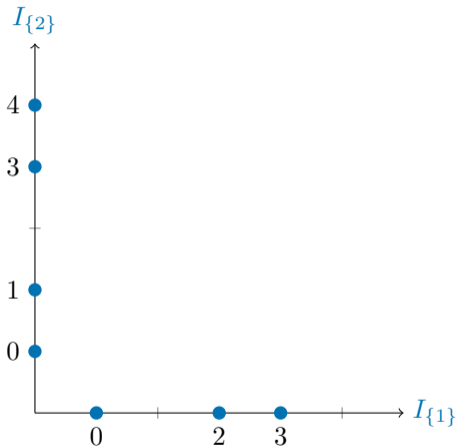
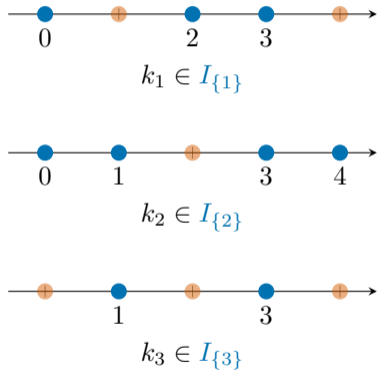


Figure: The one-dimensional index sets $I_{\{j\}}$.

Figure: The two-dimensional candidate set $I_{\{1\}} \times I_{\{2\}}$.

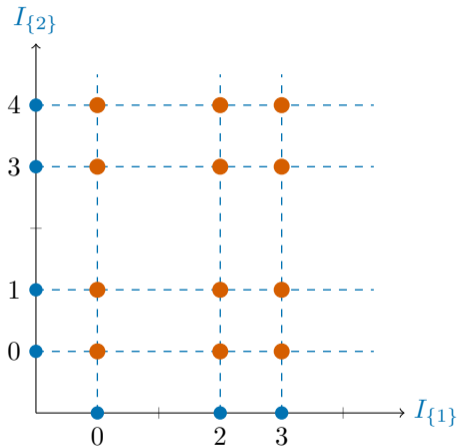
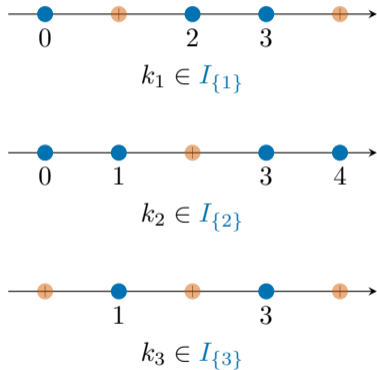


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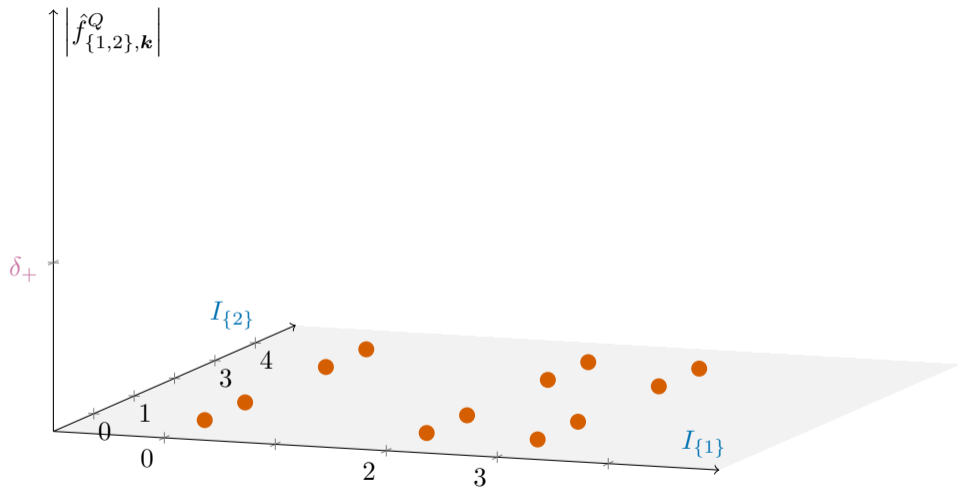


Figure: The two-dimensional detection step in the dimensions $\{1, 2\}$. (Only $r = 1$ detection iteration shown.)

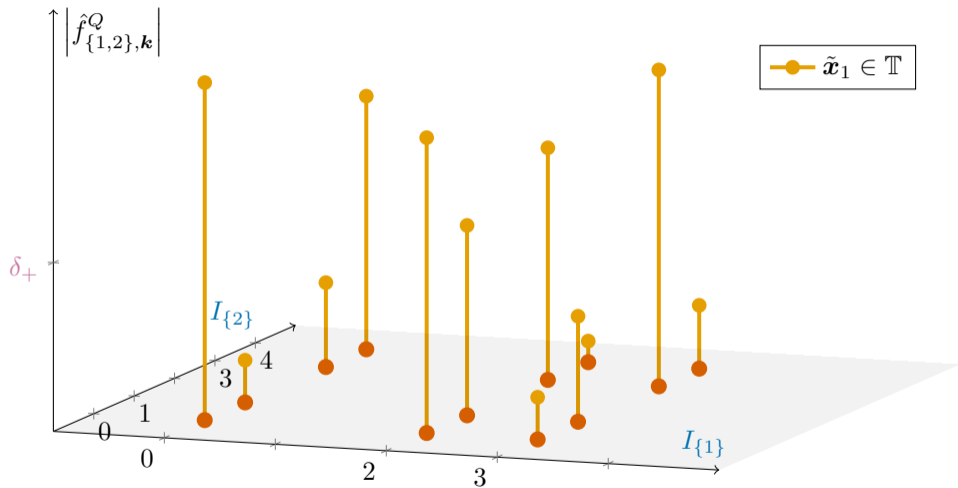


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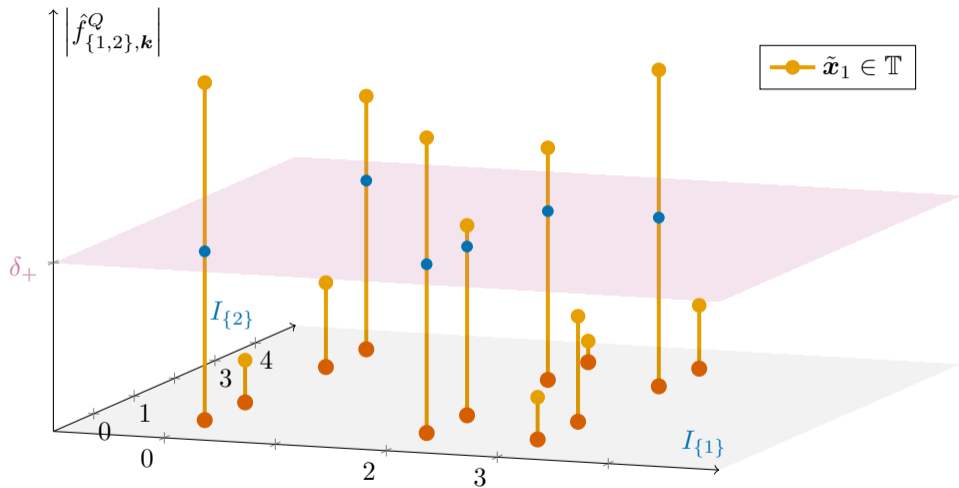


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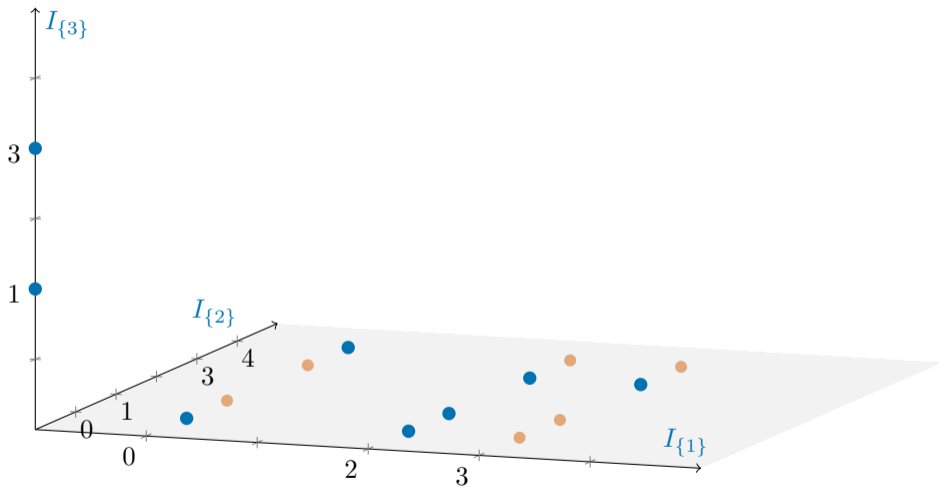


Figure: The two-dimensional index set $I_{\{1,2\}}$ and the three-dimensional candidate set $I_{\{1,2\}} \times I_{\{3\}}$.

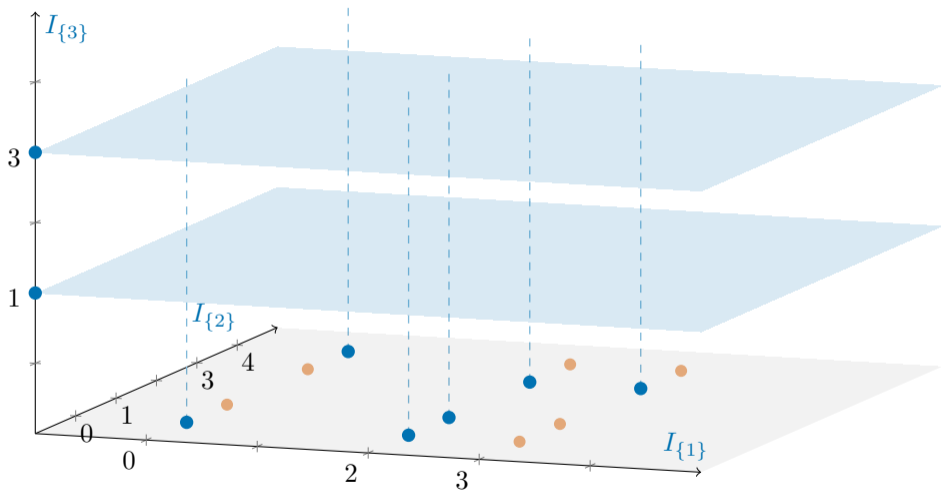


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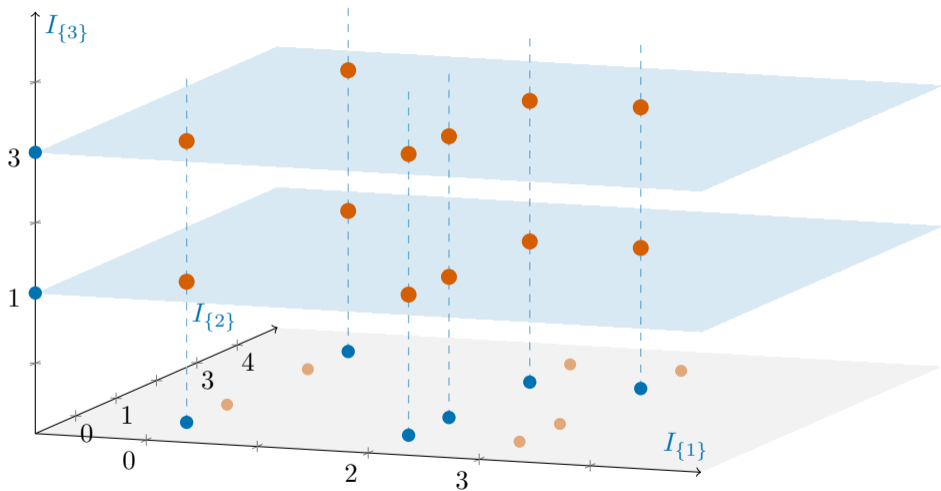


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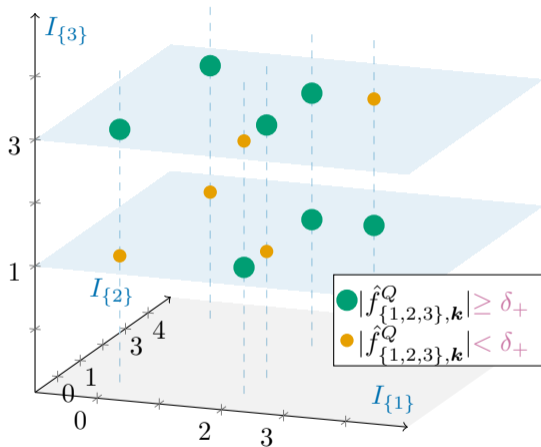


Figure: The three-dimensional detection step with the detected index set $I_{\{1,2,3\}}$.

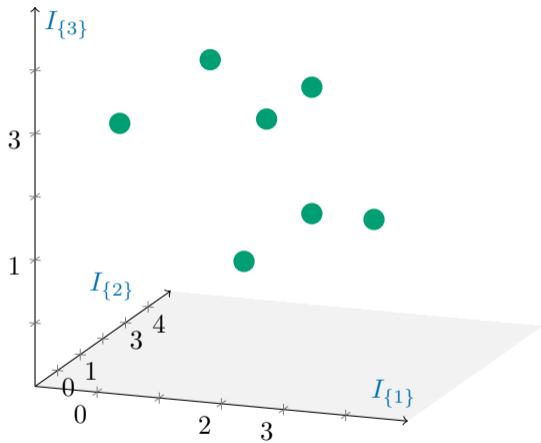


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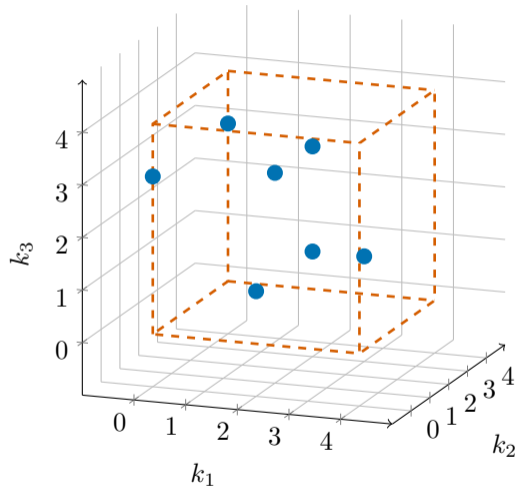


Figure: The correctly detected index set $I = I_{\{1,2,3\}}$.

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Theorem [Kämmerer, Potts, T. '22]

- ▶ given: function $f \in L_2(\mathcal{D}, \mu)$, BOPB constant B , threshold $\delta > 0$, failure prob. $\varepsilon \in (0, 1)$
- ▶ index set: $I_{3\delta} := \{\mathbf{k} \in \mathbb{N}^d : |c_{\mathbf{k}}| \geq 3\delta\}$
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Improvements: works for **function approximation** in **any** BOPB!

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Test function (from [Potts, Volkmer '16], [Kämmerer, Kraemer, Volkmer '21])

$$f(\mathbf{x}) := \prod_{j \in \{1,3,8\}} N_2(x_j) + \prod_{j \in \{2,5,6,10\}} N_4(x_j) + \prod_{j \in \{4,7,9\}} N_6(x_j)$$

$f \in L_2(\mathbb{T}^{10}, \mu)$, with the B-Spline of order $m \in \mathbb{N}$

$$N_m(x) := C_m \sum_{k \in \mathbb{Z}} \operatorname{sinc} \left(\frac{\pi}{m} k \right)^m (-1)^k e^{2\pi i k x}$$

Parameters

- ▶ hyperbolic cross search space:

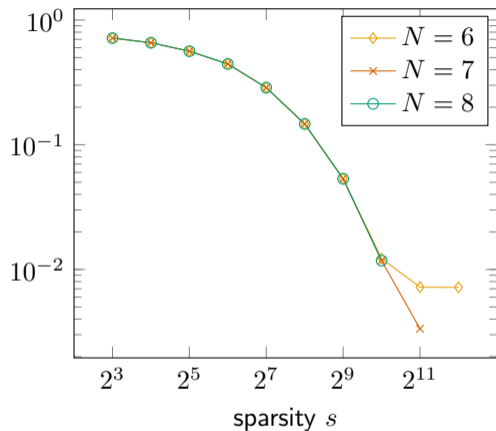
$$\Gamma = \left\{ \mathbf{k} \in \mathbb{Z}^d : \prod_{j=1}^d \max(1, 2|k_j|) \leq 2^N \right\}$$

- ▶ detection iterations $r = 5$
- ▶ detection threshold $\delta_+ = 10^{-12}$

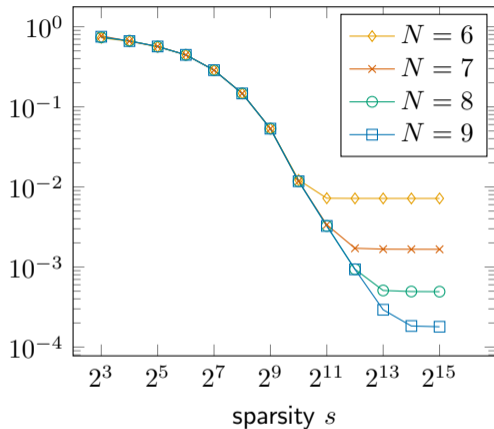
Approximation error (“How good is I ?”)

- ▶ rel. $L_2(\mathbb{T}^{10}, \mu)$ error:

$$\frac{\|f - S_I f\|_{L_2(\mathbb{T}^{10})}}{\|f\|_{L_2(\mathbb{T}^{10})}}$$

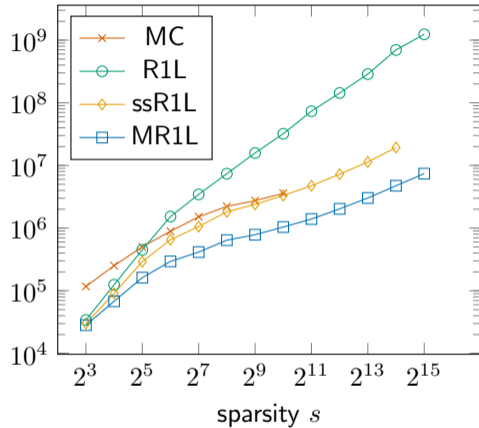


(a) approx. error for MC

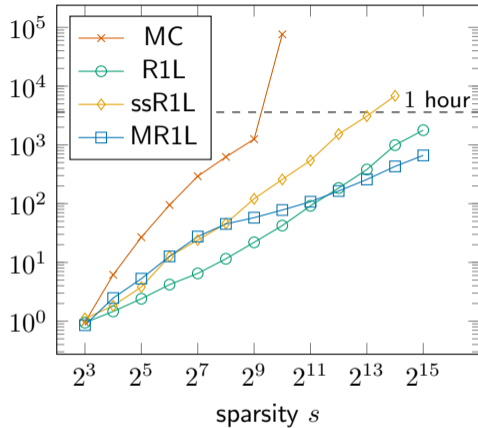


(b) approx. error for R1L, ssR1L & MR1L

Figure: Approximation results for the 10-dimensional periodic test function



(a) amount of samples for $N = 8$



(b) computation time (in seconds) for $N = 8$

Figure: Approximation results for the 10-dimensional periodic test function

Test function (from [Volkmer '17], [Potts, Volkmer '17])

$$f(\mathbf{x}) := \prod_{j \in \{1,3,4,7\}} B_2(x_j) + \prod_{j \in \{2,5,6,8,9\}} B_4(x_j)$$

$f \in L_2([-1, 1]^9, \mu_{\text{Cheb}})$, with B_2 and B_4 shifted, scaled and dilated B-Splines of order 2 and 4

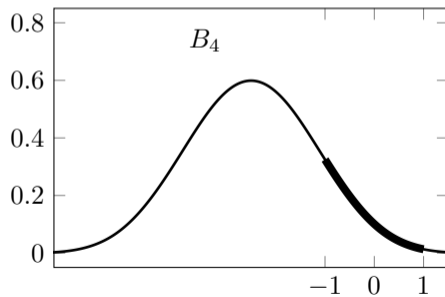
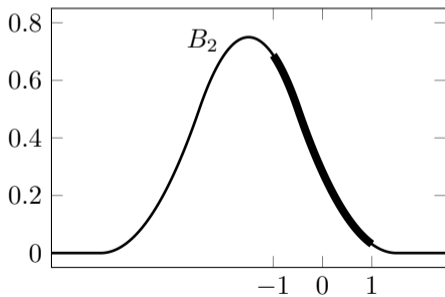
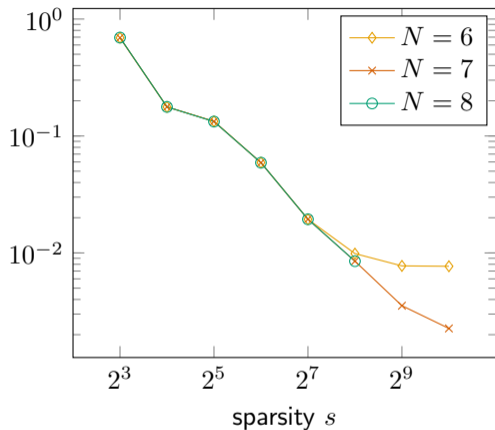
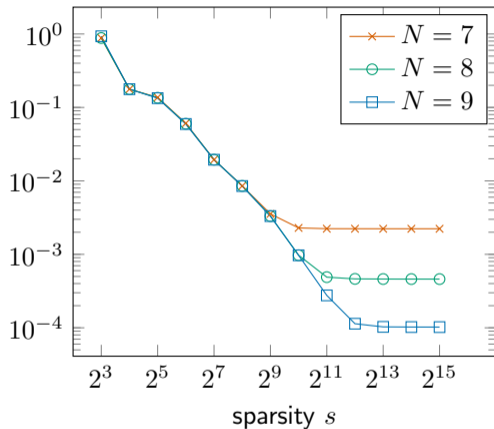


Figure: B-Splines B_2 and B_4 and the considered domain $[-1, 1]$

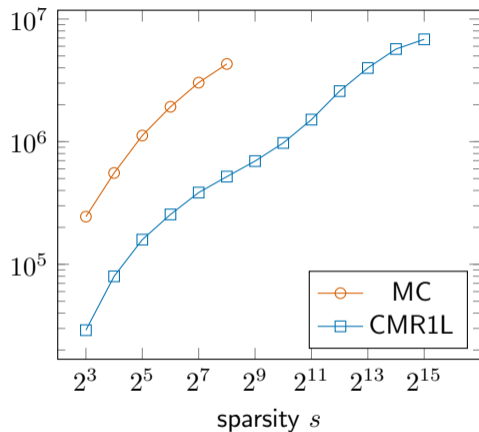


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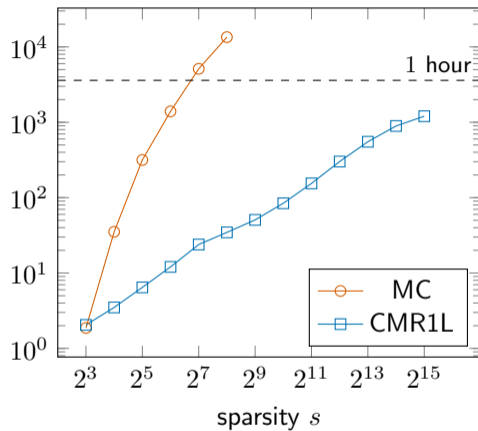


(b) approx. error for CMR1L

Figure: Approximation results for the 9-dimensional non-periodic test function



(a) amount of samples for $N = 8$



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② The algorithm

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③ Theorectical detection guarantee for function approximation

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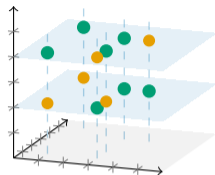
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⑤ Conclusion

► What did we do?

- generalization of the dimension-incremental method to arbitrary BOPB
- generalized algorithm also works
 - for various search spaces Γ
 - with other dimension-incremental strategies (dyadic, data-driven, ...)
 - ...



► What did we show?

- first theoretical detection guarantee of the dimension-incremental method for function approximation
- proof technique can be generalized for other reconstruction methods

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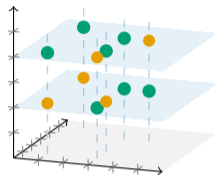
- promising numerical tests with good approximations
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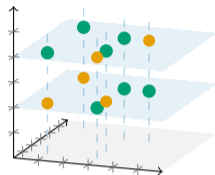
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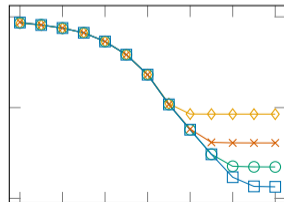
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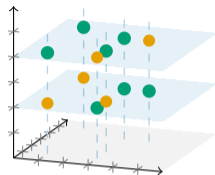
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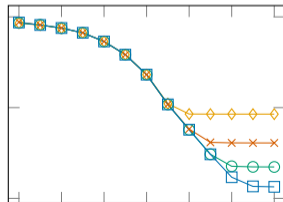
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Questions? Ideas? Suggestions?