

Efficient inversion of the multivariate nonequispaced fast Fourier transform

Task

Nonequispaced fast Fourier transform (NFFT)

Fast algorithm to evaluate a trigonometric polynomial

$$f(\mathbf{x}) = \sum_{\mathbf{k} \in \mathcal{I}_M} \hat{f}_{\mathbf{k}} e^{2\pi i \mathbf{k} \mathbf{x}}$$

at nonequispaced nodes $\mathbf{x}_j \in \mathbb{T}^d \cong [-\frac{1}{2}, \frac{1}{2}]^d$, $j = 1, \dots, N$, $d \in \mathbb{N}$, with index set $\mathcal{I}_M := \{\mathbf{k} \in \mathbb{Z}^d : -\frac{M}{2} \leq k_t < \frac{M}{2}, t = 1, \dots, d\}$

- ▶ equispaced grid \mathbf{x}_j and $M^d = |\mathcal{I}_M| = N$
⇒ **FFT**: $\mathcal{O}(|\mathcal{I}_M| \log(|\mathcal{I}_M|))$

- ▶ **Complexity**: $\mathcal{O}(|\mathcal{I}_M| \log(|\mathcal{I}_M|) + N)$

- ▶ Ajoint problem:

$$h_{\mathbf{k}} = \sum_{j=1}^N f_j e^{-2\pi i \mathbf{k} \mathbf{x}_j}, \quad \mathbf{k} \in \mathcal{I}_M$$

- ▶ nonequispaced Fourier matrix $\mathbf{A} := (e^{2\pi i \mathbf{k} \mathbf{x}_j})_{j=1, \mathbf{k} \in \mathcal{I}_M}^N$

- ▶ Factorizations: $\mathbf{A} \approx \mathbf{BFD}$ and $\mathbf{A}^* \approx \mathbf{D}^* \mathbf{F}^* \mathbf{B}^*$
sparse FFT diagonal

Inversion (iNFFT)

Given: $\mathbf{f} := (f(\mathbf{x}_j))_{j=1}^N$

Find: $\hat{\mathbf{f}} := (\hat{f}_{\mathbf{k}})_{\mathbf{k} \in \mathcal{I}_M}$, $\hat{f}_{\mathbf{k}} \in \mathbb{C}$

Motivation: FFT is invertible
various applications: MRI, solution of PDEs, ...

Challenge: in general $N \neq |\mathcal{I}_M|$

Now: new direct method for $d \geq 1$

Problems

- (1) Solve

$$\mathbf{A} \hat{\mathbf{f}} = \mathbf{f}, \quad \Rightarrow \text{inverse NFFT (iNFFT)}$$

given: \mathbf{f} , find: $\hat{\mathbf{f}}$.

- (2) Solve

$$\mathbf{A}^* \mathbf{f} = \mathbf{h}, \quad \Rightarrow \text{inverse adjoint NFFT (iNFFT*)}$$

given: \mathbf{h} , find: \mathbf{f} .

Preliminary considerations

- ▶ For illustration consider equispaced nodes

$$\mathbf{x}_j = \frac{1}{n} \mathbf{j} \in \mathbb{T}^d, \quad \mathbf{j} \in \mathcal{I}_n, \quad \mathbf{n} := (n, \dots, n)^T, \quad \text{with } |\mathcal{I}_n| = N.$$

Obtain

$$\mathbf{A} = (e^{2\pi i \mathbf{k} \frac{\mathbf{j}}{n}})_{\mathbf{j} \in \mathcal{I}_n, \mathbf{k} \in \mathcal{I}_M} \quad \text{and} \quad \mathbf{A}^* = (e^{-2\pi i \mathbf{k} \frac{\mathbf{j}}{n}})_{\mathbf{k} \in \mathcal{I}_M, \mathbf{j} \in \mathcal{I}_n}$$

and hence matrix products

$$\mathbf{A}^* \mathbf{A} = N \mathbf{I}_{|\mathcal{I}_M|} \quad \text{for } N \geq |\mathcal{I}_M|,$$

$$\mathbf{A} \mathbf{A}^* = |\mathcal{I}_M| \cdot \mathbf{I}_N \quad \text{for } N \leq |\mathcal{I}_M| \text{ with } N \mid |\mathcal{I}_M|.$$

⇒ In these special cases inversion given by composition of Fourier matrices.

- ▶ Whereas, for **nonequispaced nodes**

$$\mathbf{x}_j \in [-\frac{1}{2}, \frac{1}{2}]^d, \quad j = 1, \dots, N,$$

we have

$$\mathbf{A}^* \mathbf{A} \neq N \mathbf{I}_{|\mathcal{I}_M|} \quad \text{and} \quad \mathbf{A} \mathbf{A}^* \neq |\mathcal{I}_M| \cdot \mathbf{I}_N.$$

⇒ Look for a good approximation in general.

Matrix Approach

Basic idea

Considering

$$\mathbf{A} \hat{\mathbf{f}} = \mathbf{f}$$

we seek to find a matrix \mathbf{X} such that

$$\mathbf{X} \mathbf{A} \approx \mathbf{I}_{|\mathcal{I}_M|}$$

since then

$$\hat{\mathbf{f}} \approx \mathbf{X} \mathbf{A} \hat{\mathbf{f}} = \mathbf{X} \mathbf{f}.$$

Ideas (from equispaced case) to choose appropriate \mathbf{X} :

$$\mathbf{A}^* \mathbf{A} = N \mathbf{I}_{|\mathcal{I}_M|} \quad \text{and} \quad \mathbf{A}^* \approx \mathbf{D}^* \mathbf{F}^* \mathbf{B}^*$$

Aim:

- ↪ approximation of the form $\mathbf{D}^* \mathbf{F}^* \mathbf{B}^* \mathbf{A} \approx \mathbf{I}_{|\mathcal{I}_M|}$
- ↪ modification of matrix \mathbf{B}
- ↪ preserve band structure and arithmetic complexity

Optimization approach

Define $\tilde{\mathbf{h}} := \mathbf{D}^* \mathbf{F}^* \tilde{\mathbf{B}}^* \mathbf{f}$. It holds

$$\begin{aligned} \|\tilde{\mathbf{h}} - \hat{\mathbf{f}}\|_2 &= \|\mathbf{D}^* \mathbf{F}^* \tilde{\mathbf{B}}^* \mathbf{f} - \hat{\mathbf{f}}\|_2 = \|\mathbf{D}^* \mathbf{F}^* \tilde{\mathbf{B}}^* \mathbf{A} \hat{\mathbf{f}} - \hat{\mathbf{f}}\|_2 \\ &\leq \left\| \mathbf{D}^* \mathbf{F}^* \tilde{\mathbf{B}}^* \mathbf{A} - \mathbf{I}_{|\mathcal{I}_M|} \right\|_{\mathbb{F}} \|\hat{\mathbf{f}}\|_2. \end{aligned}$$

Optimization problem:

$$\text{Minimize}_{\tilde{\mathbf{B}}} \left\| \mathbf{D}^* \mathbf{F}^* \tilde{\mathbf{B}}^* \mathbf{A} - \mathbf{I}_{|\mathcal{I}_M|} \right\|_{\mathbb{F}}^2 = \left\| \mathbf{A}^* \tilde{\mathbf{B}} \mathbf{F} \mathbf{D} - \mathbf{I}_{|\mathcal{I}_M|} \right\|_{\mathbb{F}}^2$$

Suppose $\mathbf{A}^* \tilde{\mathbf{B}} \mathbf{F} \mathbf{D} \approx \mathbf{I}_{|\mathcal{I}_M|}$. ⇒ $\mathbf{A}^* \tilde{\mathbf{B}}$ is a pseudoinverse of $\mathbf{F} \mathbf{D}$.

Since $\mathbf{F}^* \mathbf{F} = |\mathcal{I}_{M_\sigma}| \mathbf{I}_{|\mathcal{I}_{M_\sigma}|}$, a pseudoinverse is given by $\frac{1}{|\mathcal{I}_{M_\sigma}|} \mathbf{D}^{-1} \mathbf{F}^*$.

$$\rightsquigarrow \text{Minimize}_{\tilde{\mathbf{B}}} \left\| \mathbf{A}^* \tilde{\mathbf{B}} - \frac{1}{|\mathcal{I}_{M_\sigma}|} \mathbf{D}^{-1} \mathbf{F}^* \right\|_{\mathbb{F}}^2 = \sum_{l \in \mathcal{I}_{M_\sigma}} \left\| \mathbf{A}_l^* \tilde{\mathbf{b}}_l - \frac{1}{|\mathcal{I}_{M_\sigma}|} \mathbf{D}^{-1} \mathbf{f}_l \right\|_2^2$$

⇒ inverse NFFT as well as inverse adjoint NFFT

Fast optimization of the matrix \mathbf{B}

For $d, N \in \mathbb{N}$ let $\mathbf{x}_j \in \mathbb{T}^d$, $j = 1, \dots, N$, as well as $\sigma \geq \mathbf{1}_d$ and $\mathbf{M}_\sigma := \sigma \odot \mathbf{M}$.

1. For $l \in \mathcal{I}_{M_\sigma}$:

- ▶ Determine the index set

$$\mathcal{I}_{M_\sigma, m}(l) := \{j \in \{1, \dots, N\} : \exists \mathbf{z} \in \mathbb{Z}^d \text{ with } -m\mathbf{1} \leq \mathbf{M}_\sigma \odot \mathbf{x}_j - l + \mathbf{z} \leq m\mathbf{1}\}.$$

- ▶ Compute the right side $\mathbf{v}_l := \frac{1}{|\mathcal{I}_{M_\sigma}|} \mathbf{H}_l^* \mathbf{D}^{-1} \mathbf{f}_l$.
- ▶ Determine $\mathbf{H}_l^* \mathbf{H}_l$.
- ▶ Solve for $\mathbf{H}_l^* \mathbf{H}_l \mathbf{b}_l$, i. e., compute

$$\mathbf{b}_l^{\text{opt}} := \frac{1}{|\mathcal{I}_{M_\sigma}|} (\mathbf{H}_l^* \mathbf{H}_l)^{-1} \mathbf{H}_l^* \mathbf{D}^{-1} \mathbf{f}_l, \quad l \in \mathcal{I}_{M_\sigma}.$$

2. Compose \mathbf{B}_{opt} column-wise of the $\mathbf{b}_l^{\text{opt}}$ observing the sparsity and periodicity.

Output: optimized matrix \mathbf{B}_{opt}

↪ $\mathcal{O}(|\mathcal{I}_M|)$

Workflow of iNFFT

For $d, N \in \mathbb{N}$ let $\mathbf{x}_j \in \mathbb{T}^d$, $j = 1, \dots, N$, as well as $\mathbf{f} \in \mathbb{C}^N$.

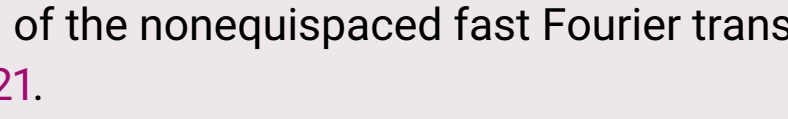
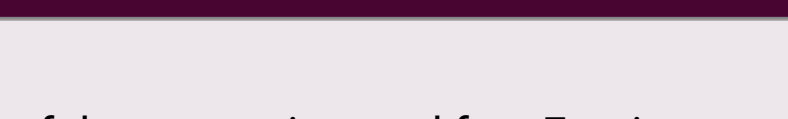
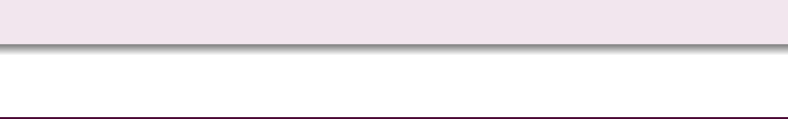
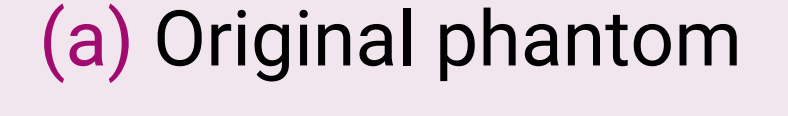
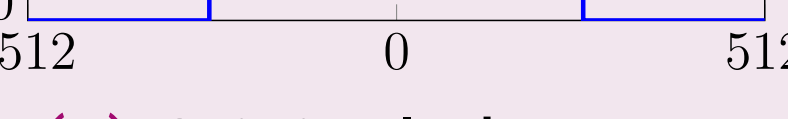
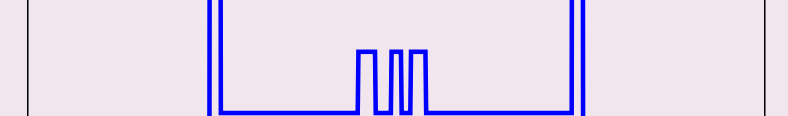
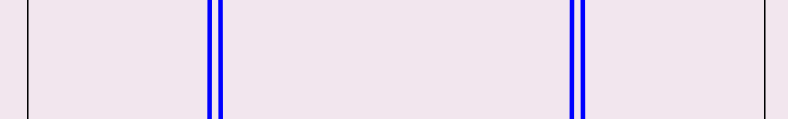
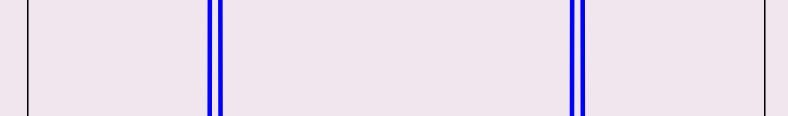
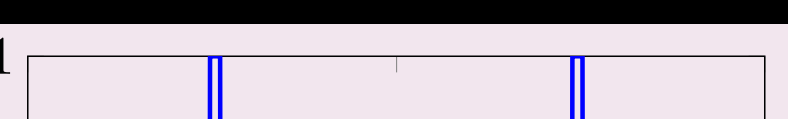
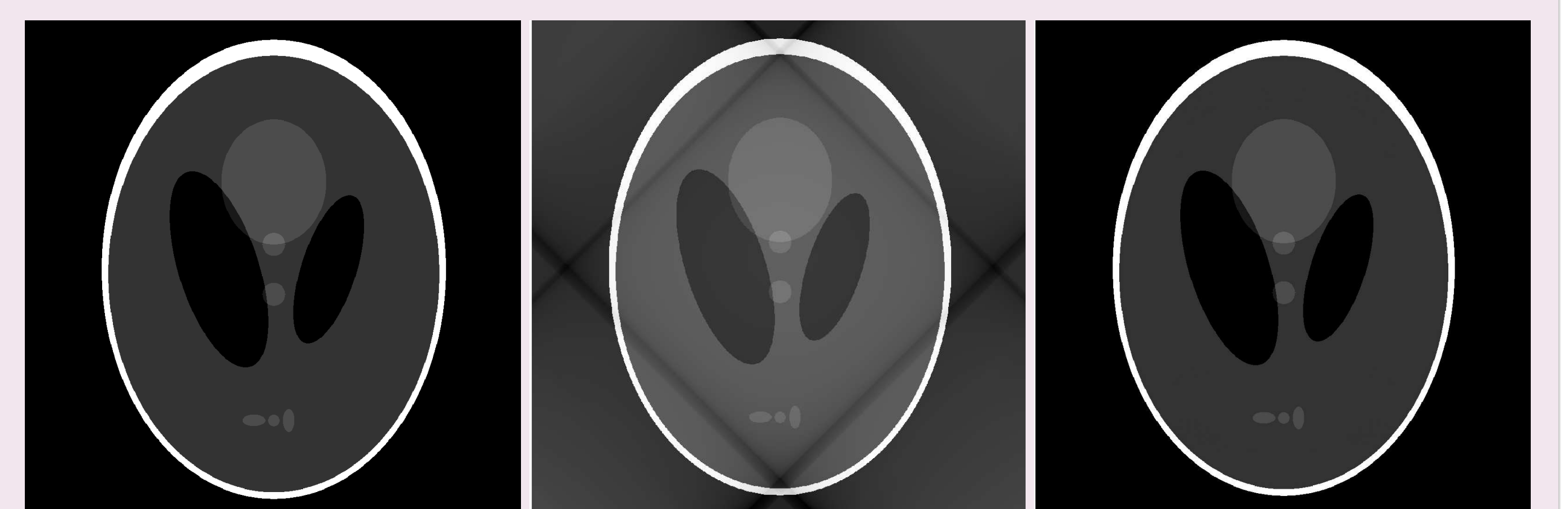
1. Compute optimized matrix \mathbf{B}_{opt} .
2. Compute the approximation $\tilde{\mathbf{h}} = \mathbf{D}^* \mathbf{F}^* \mathbf{B}_{\text{opt}}^* \mathbf{f}$ by means of a modified adjoint NFFT.

Output: $\tilde{\mathbf{h}} \approx \hat{\mathbf{f}}$

Example - Shepp-Logan phantom

- ▶ phantom size 1024×1024

- ▶ linogram grid of size $N = 2 \cdot 1024^2$



Results

Fundamentals

Dutt, Rokhlin. Fast Fourier transforms for nonequispaced data. *SIAM J. Sci. Stat. Comput.*, 1993.
Beylkin. On the fast Fourier transform of functions with singularities. *Appl. Comput. Harmon. Anal.*, 1995.
Plonka, Potts, Steidl, Tasche. Numerical Fourier Analysis. ANHA, Birkhäuser, 2018.

Corresponding paper

M. Kircheis and D. Potts. Efficient multivariate inversion of the nonequispaced fast Fourier transform. *PAMM*, 20(1):e202000120, 2021.

NFFT Software

<http://www.tu-chemnitz.de/~potts/nfft>



Refs