

# Using discretely subsampled quadrature points for function reconstruction

Felix Bartel

Lutz Kämmerer

Daniel Potts

Tino Ullrich

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Mathematik!  
TU Chemnitz

# Function recovery from samples

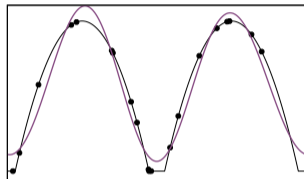
## Given:

- $L_2 = L_2(D, \nu)$
- reproducing kernel Hilbert space  $H(K) \hookrightarrow L_2$

## Goal:

- find good points  $\mathbf{X} = \{\mathbf{x}^1, \dots, \mathbf{x}^n\} \subset D$  and a sampling recovery operator  $S_{\mathbf{X}}: H(K) \rightarrow L_2$  with small **worst-case error**

$$\sup_{\|f\|_{H(K)} \leq 1} \|f - S_{\mathbf{X}} f\|_{L_2}$$



$$\begin{array}{ll} f(x) & \text{—} \\ (x^i, f(x^i)) & \bullet \\ (S_{\mathbf{X}} f)(x) & \text{—} \end{array}$$

# Least squares approximation

- let  $\eta_0, \dots, \eta_{m-1} \in L_2$  with  $n \geq m$
- ansatz  $f(\mathbf{x}) = \sum_{k=0}^{m-1} c_k \eta_k(\mathbf{x})$  with  $\mathbf{c} = (c_0, \dots, c_{m-1})^\top$  solving

$$\begin{pmatrix} \eta_0(\mathbf{x}^1) & \cdots & \eta_{m-1}(\mathbf{x}^1) \\ \vdots & \ddots & \vdots \\ \eta_0(\mathbf{x}^n) & \cdots & \eta_{m-1}(\mathbf{x}^n) \end{pmatrix} \begin{pmatrix} c_0 \\ \vdots \\ c_{m-1} \end{pmatrix} = \begin{pmatrix} f(\mathbf{x}^1) \\ \vdots \\ f(\mathbf{x}^n) \end{pmatrix}$$

# Least squares approximation

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$$\left\| \begin{pmatrix} \eta_0(\mathbf{x}^1) & \cdots & \eta_{m-1}(\mathbf{x}^1) \\ \vdots & \ddots & \vdots \\ \eta_0(\mathbf{x}^n) & \cdots & \eta_{m-1}(\mathbf{x}^n) \end{pmatrix} \begin{pmatrix} c_0 \\ \vdots \\ c_{m-1} \end{pmatrix} - \begin{pmatrix} f(\mathbf{x}^1) \\ \vdots \\ f(\mathbf{x}^n) \end{pmatrix} \right\|_{\mathbf{W}}^2 \rightarrow \min$$

for a weight matrix  $\mathbf{W} = \text{diag}(\omega_1, \dots, \omega_n)$

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$$\left\| \underbrace{\begin{pmatrix} \eta_0(\mathbf{x}^1) & \cdots & \eta_{m-1}(\mathbf{x}^1) \\ \vdots & \ddots & \vdots \\ \eta_0(\mathbf{x}^n) & \cdots & \eta_{m-1}(\mathbf{x}^n) \end{pmatrix}}_{=:L} \begin{pmatrix} c_0 \\ \vdots \\ c_{m-1} \end{pmatrix} - \underbrace{\begin{pmatrix} f(\mathbf{x}^1) \\ \vdots \\ f(\mathbf{x}^n) \end{pmatrix}}_{=:f} \right\|_W^2 \rightarrow \min$$

for a weight matrix  $W = \text{diag}(\omega_1, \dots, \omega_n)$

- the solution is given by

$$\mathbf{c} = (L^* W L)^{-1} L^* W \mathbf{f} \quad \text{and} \quad (S_X f)(\mathbf{x}) = \sum_{k=0}^{m-1} c_k \eta_k(\mathbf{x})$$

# Reproducing kernel Hilbert space

- assume finite trace  $\int_D K(\mathbf{x}, \mathbf{x}) \, d\nu(\mathbf{x}) < \infty$
- embedding  $\text{Id}_{K,\nu}: H(K) \hookrightarrow L_2$  has the representation

$$\text{Id}_{K,\nu}(f) = \sum_{k=0}^{\infty} \sigma_k \langle f, e_k \rangle_{H(K)} \eta_k$$

with  $e_k = \sigma_k \eta_k$  and

- **singular values**  $\sigma_0 \geq \sigma_1 \geq \dots \geq 0$
- **right singular functions**  $e_0, e_1, \dots$  forming an ONS in  $H(K)$
- **left singular functions**  $\eta_0, \eta_1, \dots$  forming an ONS in  $L_2$

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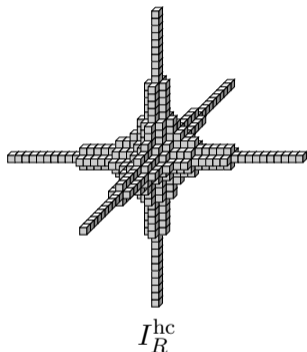
- **singular values**  $\sigma_0 \geq \sigma_1 \geq \dots \geq 0$
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  - **left singular functions**  $\eta_0, \eta_1, \dots$  forming an ONS in  $L_2$
- Kolmogorov width

$$\begin{aligned} d_m(H(K)) &:= \inf_{\substack{\ell_0, \dots, \ell_{m-1}: H(K) \rightarrow \mathbb{C} \\ \varphi_0, \dots, \varphi_{m-1} \in L_2}} \sup_{\|f\|_{H(K)} \leq 1} \left\| f - \sum_{k=0}^{m-1} \ell_k(f) \varphi_k \right\|_{L_2} \\ &= \sup_{\|f\|_{H(K)} \leq 1} \left\| f - \sum_{k=0}^{m-1} \langle f, \eta_k \rangle_{L_2} \eta_k \right\|_{L_2} = \sigma_m \end{aligned}$$

- $H_{\text{mix}}^s(\mathbb{T}^d) = \{f \in L_2 : \|f\|_{H_{\text{mix}}^s} < \infty\}$  ( $s > 1/2$ )  
with

$$\langle f, g \rangle_{H_{\text{mix}}^s} := \sum_{j \in \{0, s\}^d} \langle D^{(j)} f, D^{(j)} g \rangle_{L_2}$$

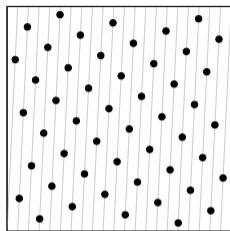
- singular functions  $\eta_{\mathbf{k}} = \exp(2\pi i \langle \mathbf{k}, \cdot \rangle)$
- singular values  $\sigma_{\mathbf{k}} = \prod_{j=1}^d (1 + (2\pi |k_j|)^{2s})^{-1/2}$
- Kolmogorov width  $d_m(H_{\text{mix}}^s) = m^{-s} (\log m)^{(d-1)s}$





# Structured points: rank-1 lattice

- $D = \mathbb{T}^d$ ,  $\{\eta_0, \dots, \eta_{m-1}\} = \{\exp(2\pi i \langle \mathbf{k}, \cdot \rangle)\}_{\mathbf{k} \in I}$
- $\mathbf{X} = \Lambda(\mathbf{z}, n) = \left\{ \frac{i}{n} \mathbf{z} \bmod \mathbf{1}, i = 0, \dots, n-1 \right\}$
- [Nuyens '07], [Kämmerer, Potts, Volkmer '15]:
  - given  $I$ , there exist algorithms that find  $\mathbf{z}$  and  $n$  such that  $\mathbf{L}^* \mathbf{W} \mathbf{L} = \mathbf{I}$
  - multiplication with  $\mathbf{L}$  can be carried out with the LFFT in  $\mathcal{O}(n \log n)$



$$\mathbf{z} = (1, 21)^\top, n = 55$$

## Theorem

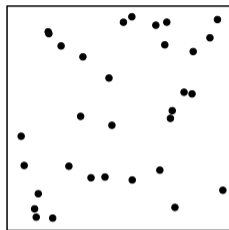
[Byrenheid, Kämmerer, T. Ullrich, Volkmer '17]

$$n^{-s} \lesssim \sup_{\|f\|_{H_{\text{mix}}^s} \leq 1} \|f - S_{\mathbf{X}} f\|_{L_2}^2 \lesssim n^{-s} (\log n)^{(d-2)s + (d-1)}$$

- $\mathbf{X} = \{\mathbf{x}^1, \dots, \mathbf{x}^n\}$  drawn randomly i.i.d. w.r.t. Lebesgue measure such that

$$n \sim m \log m$$

- no fast matrix-vector multiplication due to lack of structure



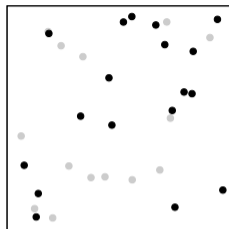
## Theorem

[M. Ullrich, Krieg '19]

$$n^{-2s} (\log n)^{2(d-1)s} \lesssim \sup_{\|f\|_{H_{\text{mix}}^s} \leq 1} \|f - S_{\mathbf{X}} f\|_{L_2}^2 \lesssim n^{-2s} (\log n)^{2(d-1)s+2s}$$

# Subsampled random points

- $\mathbf{X} = \{\mathbf{x}^1, \dots, \mathbf{x}^n\}$  drawn randomly i.i.d. w.r.t. Lebesgue measure such that  $n \sim m \log m$
- there exist  $\mathbf{X}' = \{\mathbf{x}^{i_1}, \dots, \mathbf{x}^{i_{n'}}\} \subset \mathbf{X}$  subsampled points suitable for reconstruction with
$$n' \in \mathcal{O}(|I|)$$
- no fast matrix-vector multiplication



## Theorem

[Nagel, Schäfer, T. Ullrich '21], [B., Schäfer, T. Ullrich '22]

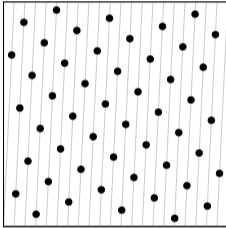
$$n^{-s}(\log n)^{2(d-1)s} \lesssim \sup_{\|f\|_{H(K)} \leq 1} \|f - S_{\mathbf{X}} f\|_{L_2}^2 \lesssim n^{-2s}(\log n)^{2(d-1)s+1}$$

## Theorem

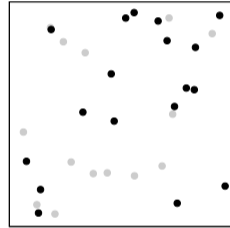
[Dolbeaut, Krieg, M. Ullrich '22]

$$\sup_{\|f\|_{H(K)} \leq 1} \|f - S_{\mathbf{X}} f\|_{L_2}^2 \sim n^{-2s}(\log n)^{2(d-1)s}$$

**Goal:** combine both advantages



structure for **fast algorithms**

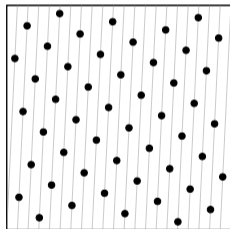


**optimal rates**

# Approach (motivated by [Kunis, Rauhut '08])

1

$\mathbf{X}_{\text{MZ}}$ : good  
deterministic set of  
 $M$  points

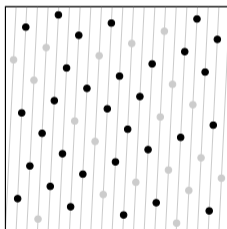


$$\mathbf{X}_{\text{MZ}} = \{\mathbf{x}^1, \dots, \mathbf{x}^M\}$$
$$\mathbf{W}_{\text{MZ}} \in [0, \infty)^{M \times M}$$

random  
→  
subsampling

2

$\mathbf{X}$ : randomly  
subsample  $n$  points  
with  $n \sim m \log m$

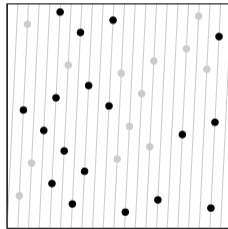


$$\mathbf{X} = \{\mathbf{x}^i\}_{i \in J}$$
$$\mathbf{W} \in [0, \infty)^{|\mathbf{X}| \times |\mathbf{X}|}$$

BSS  
→  
subsampling

3

$\mathbf{X}'$ : BSS subsample  
 $n'$  points with  
 $n' \sim m$



$$\mathbf{X}' = \{\mathbf{x}^i\}_{i \in J'}$$
$$\mathbf{W}' \in [0, \infty)^{|\mathbf{X}'| \times |\mathbf{X}'|}$$

- points  $\mathbf{X}_{\text{MZ}}$  and weights  $\omega_1, \dots, \omega_M$  fulfill a  $L_2$  **Marcinkiewicz-Zygmund inequality** for  $V \subset L_2$ , iff

$$A\|g\|_{L_2}^2 \leq \sum_{i=1}^M \omega_i |g(\mathbf{x}^i)|^2 \leq B\|g\|_{L_2}^2 \quad \text{for all } g \in V$$

- the full system matrix is well-conditioned

$$A \leq \sigma_{\min}^2(\mathbf{W}^{1/2}\mathbf{L}) \leq \sigma_{\max}^2(\mathbf{W}^{1/2}\mathbf{L}) \leq B$$

- **MZ inequalities are widely available:** [Mhaskar, Narcowich, Ward '01], [Keiner, Kunis, Potts '07], [Filbir, Mhaskar '11], [Müller-Gronbach, Novak, Ritter '12], [Temlyakov '18], [Gröchenig '20], [Filbir, Hielscher, Jahn, T. Ullrich to appear], ...

### Lemma

The MZ inequality on  $V$  with  $A = B$  is equivalent to the exact integration

$$\sum_{i=1}^M \omega_i g(\mathbf{x}^i) \overline{h(\mathbf{x}^i)} = \int_D g(\mathbf{x}) \overline{h(\mathbf{x})} d\nu(\mathbf{x}) \quad \text{for all } g, h \in V.$$

- **exact integration is widely available:** [Nuyens, Cools '06], [Kämmerer, Potts, Volkmer '15], [Trefethen '19], ...

# Incorporate fast algorithms

- big point set  $\mathbf{X}_{MZ} = \{\mathbf{x}^1, \dots, \mathbf{x}^M\}$  with  $\mathcal{O}(M \log M)$  algorithm
- small point set  $\mathbf{X} = \{\mathbf{x}^{j_1}, \dots, \mathbf{x}^{j_n}\} \subset \mathbf{X}_{MZ}$
- we may use the algorithm for the big point set

$$\mathbf{L}_{\mathbf{X}} = \mathbf{P}\mathbf{L}_{\mathbf{X}_{MZ}} \quad \text{where} \quad \mathbf{P} = \begin{pmatrix} 1 & 0 & 0 & \dots & 0 & 0 & 0 \\ 0 & 0 & 1 & \dots & 0 & 0 & 0 \\ \vdots & & & & & \vdots & \\ 0 & 0 & 0 & \dots & 0 & 1 & 0 \end{pmatrix} \begin{matrix} 1 \\ 2 \\ \vdots \\ n \end{matrix}$$

$j_1 \qquad j_2 \qquad \qquad j_n$

- with  $n \sim m \log m$  and  $M \sim m^2$  we obtain the same complexity as for naive matrix-vector multiplication

$$\mathcal{O}(mn) = \mathcal{O}(M \log M)$$



## Assumptions 1 of 2

Let

- $H(K)$  separable RKHS with  $\sup_{\mathbf{x} \in D} K(\mathbf{x}, \mathbf{x}) < \infty$  and  $\int_D K(\mathbf{x}, \mathbf{x}) d\mathbf{x} < \infty$ ,
- $\sigma_k$  and  $\eta_k$  singular values and functions of  $\text{Id}_{K,\nu}: H(K) \hookrightarrow L_2(\nu)$
- $I = \{0, \dots, m-1\} \subset I_M$
- $\mathbf{X}_{\text{MZ}}$  points and  $\mathbf{W}_{\text{MZ}}$  weights fulfilling an  $L_2$ -Marcinkiewicz-Zygmund inequality with constants  $A$  and  $B$  for  $V = \text{span}\{\eta_k\}_{k \in I_M}$ ,

## Assumptions 2 of 2

- $n := \left\lceil \frac{72B}{A} |I| \log |I| \right\rceil$ ,
- $\mathbf{X} = \{\mathbf{x}^i\}_{i \in J}$ ,  $|J| = n$  points drawn i.i.d. from  $\mathbf{X}_{MZ}$  w.r.t. the **discrete density weights**  $\varrho_i = \omega_i \varrho(\mathbf{x}^i)$  with

$$\varrho(\mathbf{x}^i) = \frac{\omega_i \sum_{k \in I} |\eta_k(\mathbf{x}^i)|^2}{3 \sum_{j=1}^M \omega_j \sum_{k \in I} |\eta_k(\mathbf{x}^j)|^2} + \frac{\omega_i \sum_{k \in I_{MZ} \setminus I} |e_k(\mathbf{x}^i)|^2}{3 \sum_{j=1}^M \omega_j \sum_{k \in I_{MZ} \setminus I} |e_k(\mathbf{x}^j)|^2} + \frac{\omega_i}{3}$$

- $b > 1 + \frac{1}{|I|}$ .

- uses spectral properties of the embedding
- discrete version of [M. Ullrich, Krieg '19], [Kämmerer, T. Ullrich, Volkmer '19]
- $\varrho(\mathbf{x}^i)$  neglectable for BOS, i.e.,  $\sup_{x \in D} |\eta_k(x)| \leq B$  for all  $k$

## Theorem

[B, Kämmerer, Potts, T. Ullrich '22]

Given the assumptions we construct  $\mathbf{X}' \subset \mathbf{X}_{\text{MZ}}$  with  $|\mathbf{X}'| \leq \lceil b|I| \rceil$  such that

$$\sup_{\|f\|_{H(K)} \leq 1} \|f - S_I^{\mathbf{X}'} f\|_{L_2}^2 \leq C_{A,B,b} \log |I| \left( \sup_{k \notin I} \sigma_k^2 + \frac{1}{|I|} \sum_{k \in I_{\text{MZ}} \setminus I} \sigma_k^2 \right) + \sup_{\|f\|_{H(K)} \leq 1} \|f - P_{I_{\text{MZ}}} f\|_{\ell_\infty(D)}^2$$

with probability larger than  $1 - 4/|I|$  and  $C_{A,B,b} = 79\,616 \left(\frac{B}{A}\right)^2 \frac{(b+1)^2}{(b-1)^3}$ .

## Theorem for rank-1 lattices

[B, Kämmerer, Potts, T. Ullrich '22]

Let •  $s > 1/2$ ,

- $I \subset I_{\text{MZ}} \subset \mathbb{Z}^d$  hyperbolic cross frequency index sets,  $|I| \geq 3$ ,
- $\mathbf{X}_{\text{MZ}}$  reconstructing rank-1 lattice for  $I_{\text{MZ}}$  with  $M$  points,
- $b > 1 + \frac{1}{|I|}$ .

Then we construct  $\mathbf{X}' \subset \mathbf{X}_{\text{MZ}}$  with  $|\mathbf{X}'| = n \leq \lceil b|I| \rceil$  such that

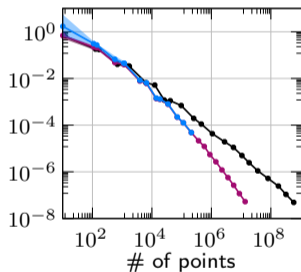
$$\sup_{\|f\|_{H_{\text{mix}}^s} \leq 1} \|f - S_I^{\mathbf{X}'} f\|_{L_2}^2 \leq C_{d,s,b} \left( n^{-2s} (\log n)^{2(d-1)s+1} \right. \\ \left. + M^{-s+1} (\log M)^{(d-1)(s+1)-s} \right)$$

with probability  $1 - 4/|I|$ .

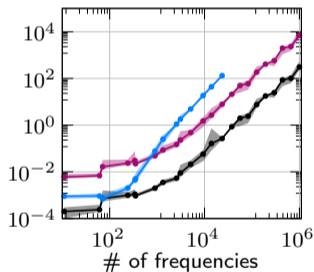
# Numerical experiment for random subsampling

- $d = 5$ ,  $f$  tensorized bump function in  $H_{\text{mix}}^s(\mathbb{T}^5)$  for  $s < 3/2$
- initial  $\mathbf{X}_{\text{MZ}}$ : fast probabilistic CBC rank-1 lattice construction [Kämmerer '20]

$L_2$ -error



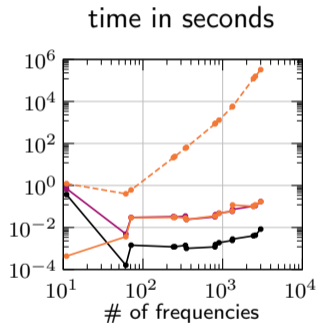
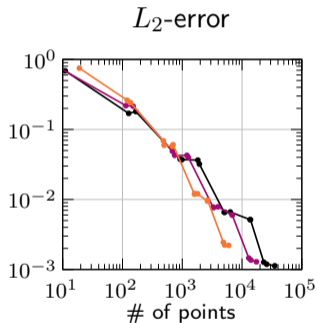
time in seconds



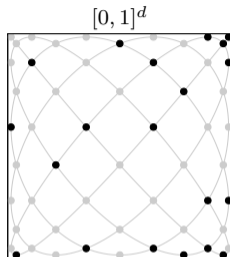
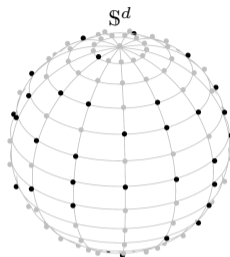
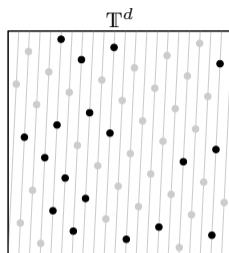
black: full rank-1 lattice  
azure: continuously random  
magenta: randomly subsampled

# Numerical experiment for random + BSS subsampling

- $d = 5$ ,  $f$  tensorized bump function in  $H_{\text{mix}}^s(\mathbb{T}^5)$  for  $s < 3/2$
- initial  $\mathbf{X}_{\text{MZ}}$ : fast probabilistic CBC rank-1 lattice construction [Kämmerer '20]



black: full rank-1 lattice  
magenta: randomly subsampled  
orange: BSS subsampled



- we proposed **random + BSS** subsampling procedure for **MZ inequalities**
  - achieving the near optimal convergence rates ✓
  - making it possible to utilize fast algorithms ✓
- [B., Kämmerer, Potts, T. Ullrich '22] on [arXiv:2208.13597](https://arxiv.org/abs/2208.13597) “On the reconstruction of functions from values at subsampled quadrature points”