

Probabilistic Subproblem Selection in Branch-and-Bound Algorithms

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General B&B Framework (1)

Problem:
$$\begin{aligned} & \max f(x) \\ & \text{s.t. } x \in \mathcal{M}. \end{aligned}$$

Step 0. Compute a compact set $X \supset \mathcal{M}$ of simple structure. Compute an upper bound $u(X) \geq \max_{x \in \mathcal{M}} f(x)$, as well as a lower bound $\ell(X) \leq \max_{x \in \mathcal{M}} f(x)$. Let \mathcal{L} be an empty list, and set $k := 1$.

Step 1. Partition X into ν compact subsets X_1, \dots, X_ν , and add them to the list \mathcal{L} .

Step 2. Calculate upper bounds $u(X_i) \geq \max_{x \in X_i \cap \mathcal{M}} f(x)$ and lower bounds $\ell(X_i) \leq \max_{x \in X_i \cap \mathcal{M}} f(x)$ for all new sets X_i .

General B&B Framework (2)

Step 3. Update the current best upper and lower bounds: Put

$$u_k := \max_{X \in \mathcal{L}} u(X) \quad \text{and} \quad \ell_k := \max_{X \in \mathcal{L}} \ell(X).$$

Step 4. Discard elements from \mathcal{L} which can not contain a global optimizer, i.e. discard all elements X with the property (i) $X \cap \mathcal{M} = \emptyset$, or (ii) $u(X) < \ell_k$.

Step 5. **Select a new $X \in \mathcal{L}$ which is to be subdivided in the next iteration**, and remove it from \mathcal{L} .

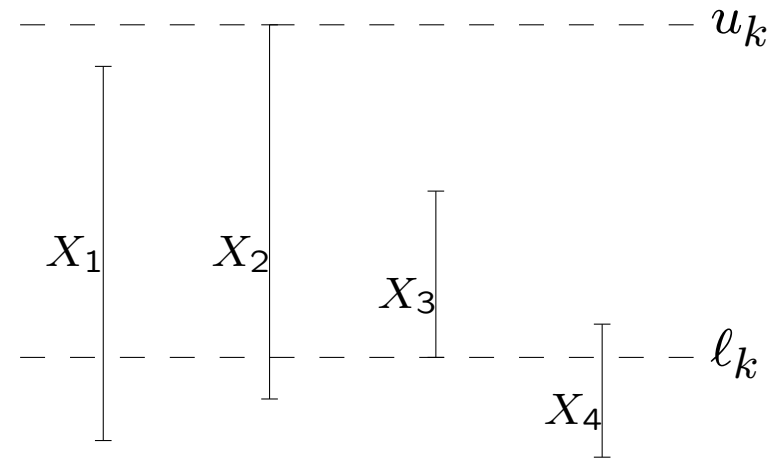
Step 6. While stopping criteria are not fulfilled, increment $k := k + 1$, and go to Step 1.

Possible Selection Rules (1)

Traditional Rule:

Select the dominating set \bar{X}
(i.e. the set with the highest
upper bound $u(X)$).

Convergence: guaranteed.



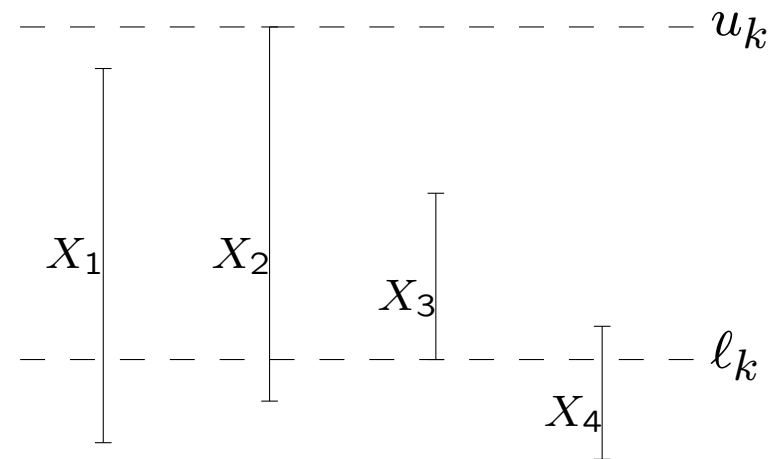
Possible Selection Rules (2)

Stix' Rule:

Select all sets whose upper bound exceeds a certain target value: $u(X) \geq \tau$.

Convergence: guaranteed.

Performance: improved.



Possible Selection Rules (3)

Csendes' Rule:

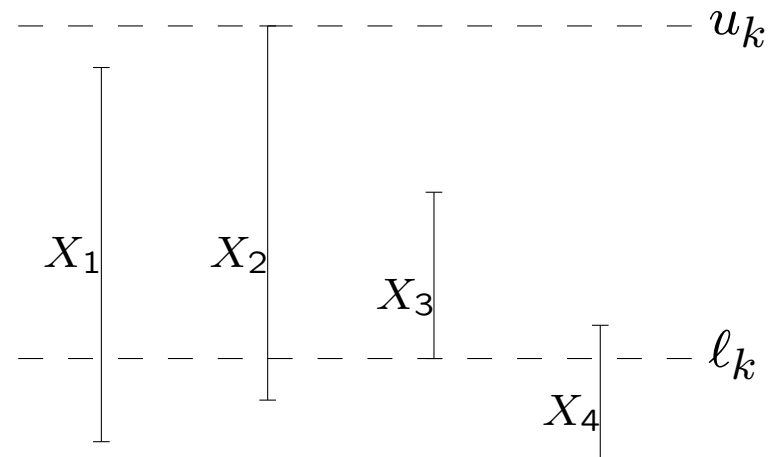
Select the set for which

$$I(X) = \frac{u(X) - \max_{x \in \mathcal{M}} f(x)}{u(X) - \ell(X)}$$

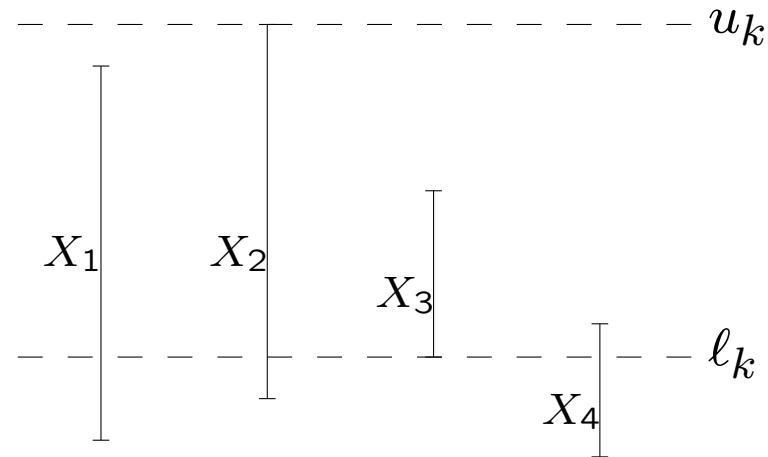
is maximized.

Convergence:
not guaranteed.

Performance: very good.



Idea: Probabilistic Selection Rule



Some sets “more likely” to contain global solution than others.

Therefore: select according to a probability distribution!

$$P_k(X_i \text{ is selected in iteration } k) \propto \frac{u(X_i) - \ell_k}{u_k - \ell_k}.$$

Some Definitions

Definition 1:

A partitioning procedure is called **exhaustive**, if every nested sequence $\{X_i\}_{i \in \mathbb{N}}$ of partition sets eventually shrinks to a singleton:

$$\bigcap_{i \in \mathbb{N}} X_i = \{x^*\}.$$

Definition 2:

We say that a bounding procedure has the **zero convergence property**, if for every exhaustive sequence $\{X_i\}_{i \in \mathbb{N}}$ of subsets of X_0 we have

$$\lim_{i \rightarrow \infty} u(X_i) = \lim_{i \rightarrow \infty} \ell(X_i).$$

Convergence of B&B

Proposition:

Assume that in a Branch-and-Bound procedure

- (a) a selection rule is used which ensures that the dominating set is chosen infinitely often,
- (b) an exhaustive partitioning procedure is used, and
- (c) the bounding procedure has the zero convergence property.

Then the B&B procedure is convergent, i.e.

$$\lim_{k \rightarrow \infty} u_k = \lim_{k \rightarrow \infty} \ell_k.$$

Convergence of B&B with Probabilistic Selection Rule

Theorem:

Assume that in a Branch-and-Bound procedure

- (a) a probabilistic selection rule is used which fulfills

$$\sum_{k=0}^{\infty} P_k(\bar{X}_k \text{ is selected in iteration } k) = +\infty,$$

- (b) an exhaustive partitioning procedure is used, and
(c) the bounding procedure has the zero convergence property.

Then the B&B procedure converges with probability 1.

Which probabilistic selection rules fulfill convergence conditions? (1)

$$P_k(X_i \text{ is selected in iteration } k) = \frac{u(X_i) - \ell_k}{\sum_{X \in \mathcal{L}_k} (u(X) - \ell_k)}.$$

The dominating set is assigned the highest probability among all $X \in \mathcal{L}_k$, and $|\mathcal{L}_k| \leq k\nu$.

$$P_k(\bar{X}_k \text{ is selected in iteration } k) \geq \frac{1}{k\nu}.$$

Therefore:

$$\sum_{k=0}^{\infty} P_k(\bar{X}_k \text{ is selected in iteration } k) \geq \frac{1}{\nu} \sum_{k=0}^{\infty} \frac{1}{k} = +\infty.$$

Which probabilistic selection rules fulfill convergence conditions? (2)

Theorem:

Assume that in a Branch-and-Bound procedure

- (a) a selection rule is used which ensures that in each iteration the dominating set is assigned the highest probability among all elements of the list,
- (b) the number of elements in the list increases at most linearly in the iterations,
- (c) an exhaustive partitioning procedure is used, and
- (d) the bounding procedure has the zero convergence property.

Then the B&B procedure is convergent with probability 1.

Which probabilistic selection rules fulfill convergence conditions? (3)

Other possibility:

$$P'_k(X_i \text{ is selected in iteration } k) = \frac{F\left(\frac{u(X_i) - \ell_k}{u_k - \ell_k}\right)}{\sum_{X \in \mathcal{L}_k} F\left(\frac{u(X) - \ell_k}{u_k - \ell_k}\right)},$$

where F is any probability distribution function on $[0, 1]$.

Order of probabilities of the sets does not change:

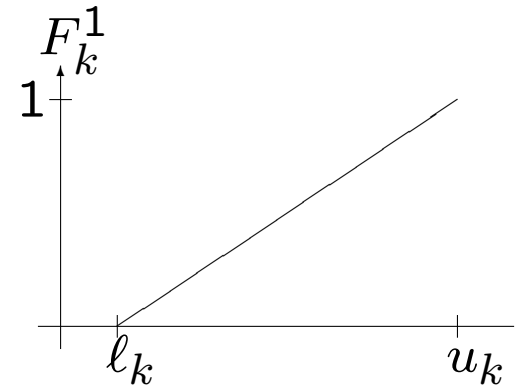
$$P_k(X_i) \leq P_k(X_j) \iff P'_k(X_i) \leq P'_k(X_j).$$

\Rightarrow Convergence conditions fulfilled.

Special Instances of F (1)

Uniform distribution:

yields our original probability
measure:

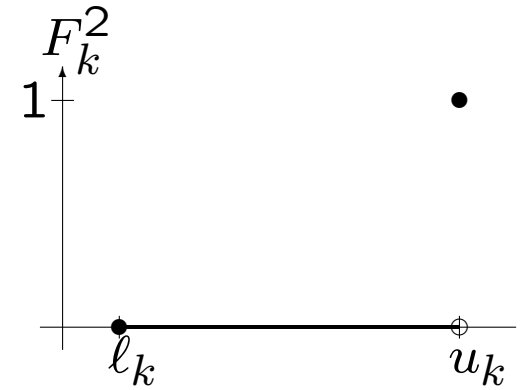


$$P_k(X_i \text{ is selected in iteration } k) = \frac{u(X_i) - l_k}{\sum_{X \in \mathcal{L}_k} (u(X) - l_k)}.$$

Special Instances of F (2)

Distribution F^2 :

$$F^2(x) = \begin{cases} 0 & \dots & 0 \leq x < 1 \\ 1 & \dots & x = 1 \end{cases}$$



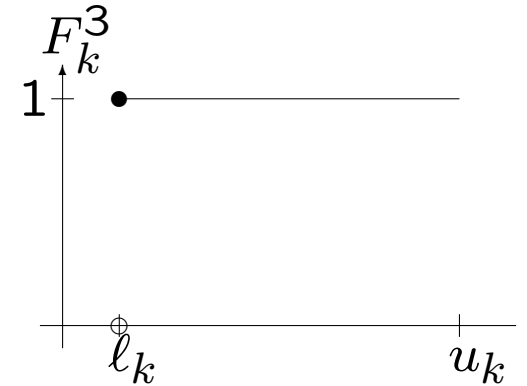
yields the purely deterministic selection rule:

Always select the set with the highest upper bound.

Special Instances of F (3)

Distribution F^3 :

$$F^3(x) = 1 \quad \text{for} \quad 0 \leq x \leq 1.$$



yields the purely random selection rule:

$$P_k^3(X_i \text{ is selected in iteration } k) = \frac{1}{\sum_{X \in \mathcal{L}} 1} = \frac{1}{|\mathcal{L}|},$$

Numerical Tests on Maximum Clique Problem

Reformulation:

$$\begin{aligned} \max x^T A_G x \\ \text{s.t. } x \in \Delta \end{aligned} \quad (\star)$$

where: A_G is regularization of adjacency matrix – indefinite!!
 Δ is the standard simplex.

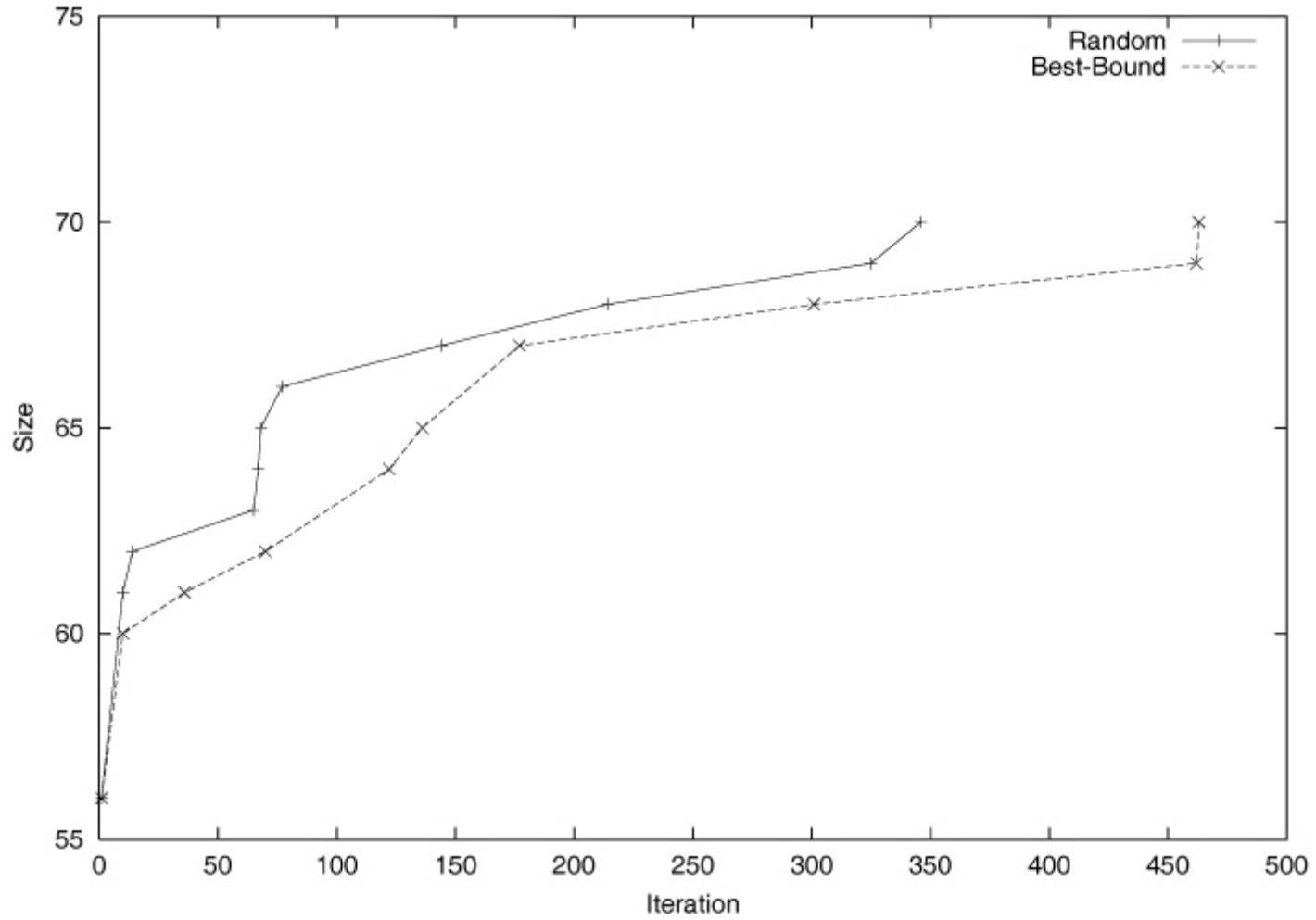
Theorem (I.M.Bomze):

There is a one to one correspondance between:

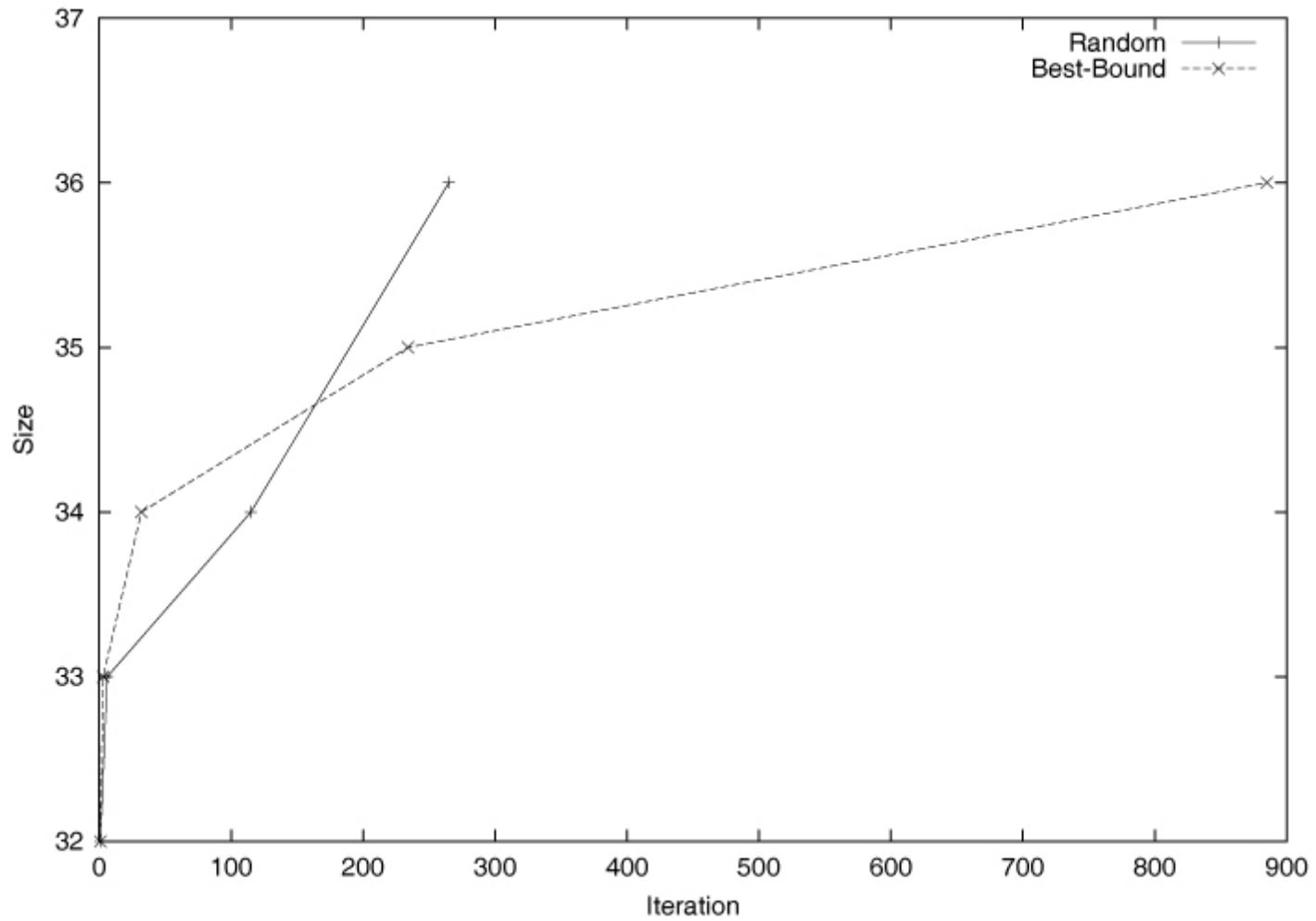
- 1.) local solutions of (\star) and maximal cliques,
- 2.) global solutions of (\star) and maximum cliques.

We used DIMACS benchmark graphs as testproblems.

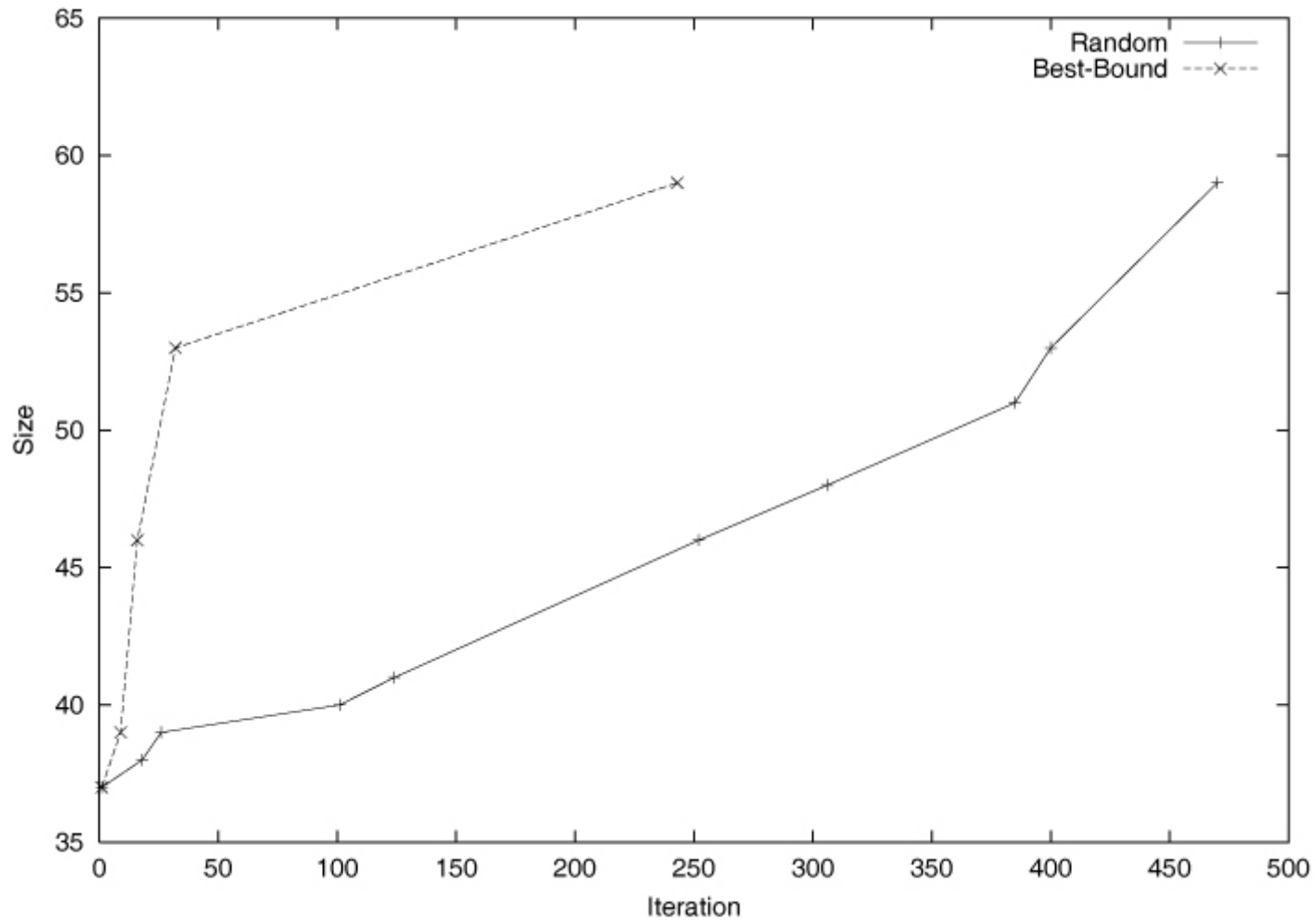
Performance: san200_0.9_1



Performance: san200_0.9_3



Performance: san200_0.9_2



Experiments on DIMACS Graphs:

name	dimension	random			best bound	performance
		min.	avg.	max.		
brock200_1	200	1497	1935.45	2103	2802	69%
brock400_1	400	25	72.85	102	105	69%
brock400_2	400	17	67.20	107	90	71%
brock400_3	400	41	163.45	259	99	165%
brock400_4	400	8	254.45	560	330	77%
MANN_a27	378	9	16.25	30	123	13%
p_hat300-2	300	2833	4603.75	5079	3628	127%
p_hat300-3	300	53	145.10	288	156	93%
san200_0.9_2	200	136	253.90	348	243	104%
san200_0.9_3	200	92	853.2	1732	885	96%
san400_0.7_1	400	23	142.30	207	88	162%
san400_0.7_2	400	688	1046.60	1865	2777	38%
san400_0.7_3	400	38	833.70	1291	462	181%
san400_0.9_1	400	18	192.10	511	272	71%
sanr200_0.7	200	12	77.25	132	144	53%
sanr200_0.9	200	694	3468.05	8613	3512	98%
sanr400_0.7	400	44	216.90	372	257	84%