

Spectra and Finite Sections of Band Operators

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What is this course about?

- **spectral theory**
- **Fredholm theory**
- **stable approximation**

of **infinite matrices** (a_{ij}) , understood as bounded linear operators on a sequence space E .

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of **infinite matrices** (a_{ij}) , understood as bounded linear operators on a sequence space E .

Simplest example: $E = \ell^2(\mathbb{Z}, \mathbb{C})$

$$\begin{pmatrix} \ddots & \vdots & \ddots \\ \cdots & a_{ij} & \cdots \\ \ddots & \vdots & \ddots \end{pmatrix} : \begin{pmatrix} \vdots \\ x_j \\ \vdots \end{pmatrix} \mapsto \begin{pmatrix} \vdots \\ b_i \\ \vdots \end{pmatrix}$$

with indices $i, j \in \mathbb{Z}$ and entries $a_{ij}, x_j, b_i \in \mathbb{C}$

1 Classes of Infinite Matrices

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$$E = \ell^p(\mathbb{Z}^N, X)$$

with

- $p \in [1, \infty]$
- $N \in \mathbb{N}$
- X ... complex **Banach space**

Clearly, the entries of a matrix (a_{ij}) acting on E also need to be indexed by $i, j \in \mathbb{Z}^N$, and the entries a_{ij} are themselves bounded linear **operators** $X \rightarrow X$.

Vector-valued ℓ^p -spaces

$x \in E = \ell^p(\mathbb{Z}^N, X)$ iff $x = (x_k)_{k \in \mathbb{Z}^N}$ with $x_k \in X$ for $k \in \mathbb{Z}^N$ and

$$\|x\|_E := \sqrt[p]{\sum_{k \in \mathbb{Z}^N} \|x_k\|_X^p} < \infty, \quad p < \infty,$$

$$\|x\|_E := \sup_{k \in \mathbb{Z}^N} \|x_k\|_X < \infty, \quad p = \infty.$$

Vector-valued ℓ^p -spaces

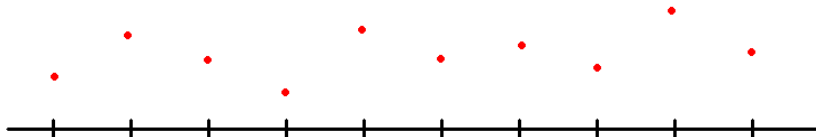
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Simplest case

$$E = \ell^p := \ell^p(\mathbb{Z}, \mathbb{C}), \quad N = 1, \quad X = \mathbb{C}$$



Vector-valued ℓ^p -spaces

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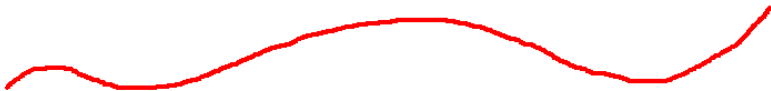
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Somewhat more complicated case

$$E = L^p(\mathbb{R}^N) \cong \ell^p(\mathbb{Z}^N, X), \quad X = L^p([0, 1]^N)$$

via identification of $f \in L^p(\mathbb{R}^N)$ with $(f|_{\alpha + [0, 1]^N})_{\alpha \in \mathbb{Z}^N}$



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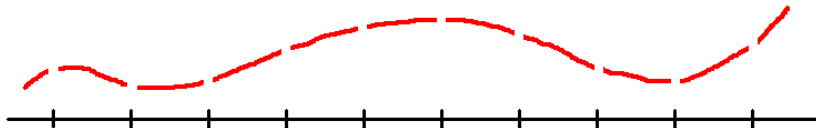
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Operators: Some first notations

Let $E = \ell^p(\mathbb{Z}^N, X)$ be one of our sequence spaces.

Then we denote by

$L(E)$... space of all **bounded linear** operators $E \rightarrow E$,

$K(E)$... space of all **compact** operators $E \rightarrow E$.

Important:

$L(E)$ is a Banach algebra and

$K(E)$ is a closed two-sided ideal in $L(E)$.

Two basic types of operators

- **shift operators**, V_k
- **multiplication operators**, M_b

$$\begin{pmatrix} \ddots & & & & & & \\ & b_{-2} & & & & & \\ & & b_{-1} & & & & \\ & & & b_0 & & & \\ & & & & b_1 & & \\ & & & & & b_2 & \\ & & & & & & \ddots \end{pmatrix} : \begin{pmatrix} \vdots \\ x_{-2} \\ x_{-1} \\ x_0 \\ x_1 \\ x_2 \\ \vdots \end{pmatrix} \mapsto \begin{pmatrix} \vdots \\ b_{-2}x_{-2} \\ b_{-1}x_{-1} \\ b_0x_0 \\ b_1x_1 \\ b_2x_2 \\ \vdots \end{pmatrix}$$

M_b is bounded if $b = (b_k)$ is so; $\|M_b\| = \|b\|_\infty$

Two basic types of operators

- **shift operators**, V_k
- **multiplication operators**, M_b

Now let them mingle: Take

- scalar multiples
- sums
- products

or combinations of those.

\implies an **operator algebra**

An algebra of shifts and multiplications

Typical elements of the algebra look like this:

$$A = V^2 M_b + 3M_b + M_a V^{-1}$$

i.e.

$$A = \begin{pmatrix} \ddots & \ddots & & & & & & & \\ \ddots & 3b_{-2} & a_{-2} & & & & & & \\ \ddots & 0 & 3b_{-1} & a_{-1} & & & & & \\ & b_{-2} & 0 & 3b_0 & a_0 & & & & \\ & & b_{-1} & 0 & 3b_1 & a_1 & & & \\ & & & b_0 & 0 & 3b_2 & \ddots & & \\ & & & & \ddots & \ddots & \ddots & & \end{pmatrix}$$

Note: $M_a M_b = M_{a \cdot b}$, $VM_b = M_{Vb}V$

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Algebras of shifts and multiplications

It is easy to see that

A is a finite sum-product of shifts and multiplications



A acts as a band matrix

The set of these operators is denoted by $BO(E)$, where BO is short for **band operator**.

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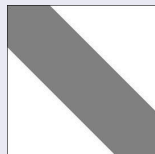


A acts as a band matrix

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Band operators

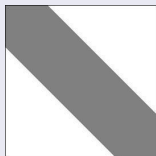
$$\underline{A \in BO(E)} : \quad A = \sum_{k=-w}^w M_{b^{(k)}} V^k$$



The number w is called the **band-width** of A .

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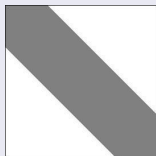


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It is nice to have an algebra of operators (i.e. a set that is closed under addition, multiplication and taking scalar multiples) but it is even nicer to have a **Banach algebra** (also closed w.r.t. $\| \cdot \|$).

Band operators

$$\underline{A \in BO(E)} : \quad A = \sum_{k=-w}^w M_{b(k)} V^k$$



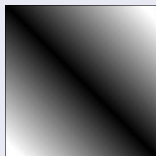
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Band-dominated operators

$$BDO(E) := \text{clos}_{L(E)} BO(E)$$

The matrices have a certain off-diagonal decay.



Algebras of shifts and multiplications

The norm under which $BDO(E)$ is closed is the usual **operator norm**

$$\|A\| := \sup_{\|x\|=1} \|Ax\|.$$

Here is **another norm**: For

$$A = \sum_{k=-w}^w M_{b^{(k)}} V^k \in BO(E),$$

we have

$$\|A\| \leq \sum_{k=-w}^w \|M_{b^{(k)}}\| \|V^k\| = \sum_{k=-w}^w \|b^{(k)}\|_{\infty} =: [A]$$

Wiener norm

$$\|A\| := \sum_{k=-w}^w \|b^{(k)}\|_{\infty}$$

Now let $\mathcal{W}(E)$ denote the completion of $BO(E)$ w.r.t. $\|\cdot\|$.

Algebras of shifts and multiplications

Wiener norm

$$\|A\| := \sum_{k=-w}^w \|b^{(k)}\|_\infty$$

Now let $\mathcal{W}(E)$ denote the completion of $BO(E)$ w.r.t. $\|\cdot\|$.
This also gives a Banach algebra (w.r.t. $\|\cdot\|$):

Wiener algebra

$$\mathcal{W}(E) = \left\{ \underbrace{\sum_{k=-\infty}^{+\infty} M_{b^{(k)}} V^k}_A : \underbrace{\sum_{k=-\infty}^{+\infty} \|b^{(k)}\|_\infty}_{\|A\|} < \infty \right\}$$

Clearly: $BO(E) \subset \mathcal{W}(E) \subset BDO(E) \subset L(E)$

Example: Laurent operator

Let \mathbb{T} be the unit circle in \mathbb{C} and fix a function $a \in L^\infty(\mathbb{T})$.

Fourier series of a :
$$\sum_{k=-\infty}^{+\infty} a_k t^k, \quad t \in \mathbb{T}.$$

This function is closely related to the associated operator

$$L(a) := \sum_{k=-\infty}^{+\infty} a_k V^k,$$

which is a so-called **Laurent operator** (constant matrix diagonals):

$$L(a) = \begin{pmatrix} \ddots & \ddots & \ddots & & \ddots \\ \ddots & a_0 & a_{-1} & a_{-2} & \\ \ddots & a_1 & a_0 & a_{-1} & \ddots \\ & a_2 & a_1 & a_0 & \ddots \\ \ddots & & \ddots & \ddots & \ddots \end{pmatrix}$$

Example: Laurent operator

The function a is called the **symbol** of the operator $L(a)$.

$L(a)$ = discrete convolution by $(a_k)_{k \in \mathbb{Z}}$ $\stackrel{\text{Fourier}}{\cong}$ multiplication by a

For simplicity, suppose $E = \ell^2$. Then $E \stackrel{\text{Fourier}}{\cong} L^2(\mathbb{T})$.

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Correspondence between $L(a)$ and its symbol a on \mathbb{T} :

Laurent operator $L(a)$	symbol a
bounded operator	bounded function
$\ L(a)\ $	$= \ a\ _\infty$

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in $BO(E)$ in $\mathcal{W}(E)$ in $BDO(E)$	trig. polynomial Wiener function continuous function

Example: Laurent operator

Here we call $a : \mathbb{T} \rightarrow \mathbb{C}$ a Wiener function and write $a \in W(\mathbb{T})$ if its Fourier coefficients are summable, i.e. $(a_k) \in \ell^1$. Note that

$$\|a\|_W := \sum_{k=-\infty}^{+\infty} |a_k| = \llbracket L(a) \rrbracket.$$

Wiener's theorem

If $a \in W(\mathbb{T})$ has no zeros then also $a^{-1} \in W(\mathbb{T})$.

Wiener's theorem in Laurent operator language

If $L(a) \in \mathcal{W}(E)$ is invertible then $L(a)^{-1} = L(a^{-1}) \in \mathcal{W}(E)$.

Operator algebras: inverse closedness

This theorem

Wiener's theorem in Laurent operator language

If $L(a) \in \mathcal{W}(E)$ is invertible then $L(a)^{-1} = L(a^{-1}) \in \mathcal{W}(E)$.

has an amazing generalisation:

The Wiener algebra is inverse closed

If $A \in \mathcal{W}(E)$ is an invertible operator then $A^{-1} \in \mathcal{W}(E)$.

And now that we're at it:

Also $BDO(E)$ is inverse closed

If $A \in BDO(E)$ is an invertible operator then $A^{-1} \in BDO(E)$.

The last two theorems hold in the general case $E = \ell^p(\mathbb{Z}^N, X)$.

The Finite Section Method, Part I

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We look at linear equations $Ax = b$ in infinitely many variables:

$$\underbrace{\begin{pmatrix} \ddots & \vdots & \vdots & \vdots & \ddots \\ \cdots & a_{-1,-1} & a_{-1,0} & a_{-1,1} & \cdots \\ \cdots & a_{0,-1} & a_{0,0} & a_{0,1} & \cdots \\ \cdots & a_{1,-1} & a_{1,0} & a_{1,1} & \cdots \\ \ddots & \vdots & \vdots & \vdots & \ddots \end{pmatrix}}_A \underbrace{\begin{pmatrix} \vdots \\ x_{-1} \\ x_0 \\ x_1 \\ \vdots \end{pmatrix}}_x = \underbrace{\begin{pmatrix} \vdots \\ b_{-1} \\ b_0 \\ b_1 \\ \vdots \end{pmatrix}}_b$$

Assumption: $A \in BDO(E)$ with $E = \ell^p$, i.e. $N = 1$ and $X = \mathbb{C}$.

Task: Given such an A and a RHS $b \in E$, find $x \in E$.

The Finite Section Method

Let A be **invertible** (bijective) as a map $E \rightarrow E$, so that $Ax = b$ is **uniquely solvable** for **every** RHS b .

How do we compute this unique solution x of $Ax = b$, i.e.

$$\sum_{j=-\infty}^{+\infty} a_{ij}x_j = b_i, \quad i \in \mathbb{Z} \quad ? \quad (1)$$

Replace the infinite system (1) by the sequence of finite systems

$$\sum_{j=-n}^n a_{ij}x_j = b_i, \quad i = -n, \dots, n$$

for $n = 1, 2, \dots$

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Or, more flexible: Take two monotonous sequences of integers

$$-\infty \leftarrow \dots < l_2 < l_1 \quad < \quad r_1 < r_2 < \dots \rightarrow +\infty$$

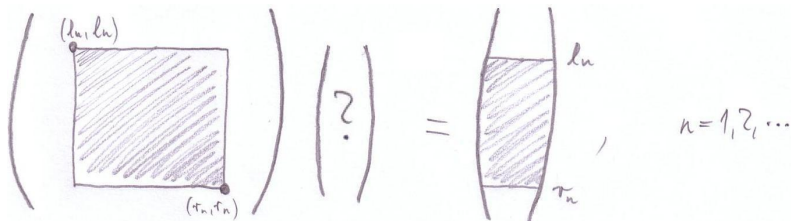
and replace the infinite system (1) by the sequence of finite systems

$$\sum_{j=l_n}^{r_n} a_{ij}x_j = b_i, \quad i = l_n, \dots, r_n \quad (2)$$

for $n = 1, 2, \dots$

The Finite Section Method

Graphically, (2) means



We say the **finite section method** is **applicable** to A if the truncated equations (2) are uniquely solvable for all $n > n_0$ and their solutions converge componentwise to the unique solution x of (1).

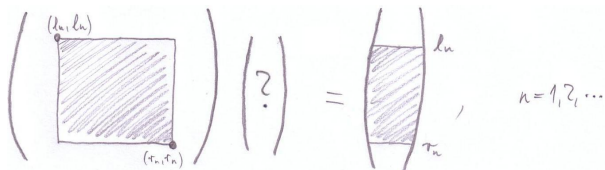
The Finite Section Method

$$\left(\begin{array}{c} (l_n, l_n) \\ \text{[shaded square]} \\ (r_n, r_n) \end{array} \right) \left(\begin{array}{c} ? \\ \cdot \end{array} \right) = \left(\begin{array}{c} \text{[shaded column]} \\ l_n \\ r_n \end{array} \right), \quad n=1, 2, \dots$$

We say the **finite section method** is **applicable** to A if the truncated equations (2) are uniquely solvable for all $n > n_0$ and their solutions converge componentwise to the unique solution x of (1).

All of this should happen **independently** of the right-hand side b . So applicability of the method only depends on A .

The Finite Section Method



Precisely: The finite section method (FSM) is applicable to A
iff
 A is invertible and its so-called **finite sections**

$$A_n := (a_{ij})_{i,j=1}^{r_n}, \quad n = 1, 2, \dots$$

form a stable sequence.

Here we call a sequence (A_n) **stable** if there exists n_0 such that

$$\sup_{n > n_0} \|A_n^{-1}\| < \infty.$$

When the Finite Section Method goes wrong

There are very simple examples where the FSM fails to apply.

Example 1: a block-flip

$$A = \begin{pmatrix} \ddots & & & & & & & & \\ & 0 & 1 & & & & & & \\ & 1 & 0 & & & & & & \\ & & & 1 & & & & & \\ & & & & 0 & 1 & & & \\ & & & & 1 & 0 & & & \\ & & & & & & \ddots & & \end{pmatrix} : \begin{pmatrix} \vdots \\ x_{-2} \\ x_{-1} \\ x_0 \\ x_1 \\ x_2 \\ \vdots \end{pmatrix} \mapsto \begin{pmatrix} \vdots \\ x_{-1} \\ x_{-2} \\ x_0 \\ x_2 \\ x_1 \\ \vdots \end{pmatrix}$$

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Only 50% of the finite systems are uniquely solvable.

Choosing good cut-off intervals $[l_n, r_n]$ will solve the problem!

When the Finite Section Method goes wrong

Example 2: the shift

$$A = \begin{pmatrix} \ddots & & & & & & \\ \ddots & & & & & & \\ & 0 & & & & & \\ & 1 & 0 & & & & \\ & & 1 & 0 & & & \\ & & & 1 & 0 & & \\ & & & & 1 & 0 & \\ & & & & & \ddots & \ddots \end{pmatrix} : \begin{pmatrix} \vdots \\ \hline x_{-2} \\ x_{-1} \\ x_0 \\ x_1 \\ \hline x_2 \\ \vdots \end{pmatrix} \mapsto \begin{pmatrix} \vdots \\ \hline 0 \\ x_{-2} \\ x_{-1} \\ x_0 \\ x_1 \\ \hline \vdots \end{pmatrix}$$

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No one of the finite systems is uniquely solvable.

When the Finite Section Method goes wrong

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$$A = \begin{pmatrix} \ddots & & & & & & \\ & \ddots & & & & & \\ & & \boxed{\begin{matrix} 0 & & & & \\ 1 & 0 & & & \\ & 1 & 0 & & \\ & & 1 & 0 & \\ & & & 1 & 0 \end{matrix}} & & & & \\ & & & \ddots & & & \end{pmatrix} : \begin{pmatrix} \vdots \\ \hline x_{-2} \\ x_{-1} \\ x_0 \\ x_1 \\ \hline x_2 \\ \vdots \end{pmatrix} \mapsto \begin{pmatrix} \vdots \\ \hline 0 \\ x_{-2} \\ x_{-1} \\ x_0 \\ x_1 \\ \hline \vdots \end{pmatrix}$$

No one of the finite systems is uniquely solvable.

Adapting the cut-off points $[l_n, r_n]$ will **not** help here!
Instead, place the corners of A_n **along another diagonal!**

The Finite Section Method

Clearly, there is a lot of room in a bi-infinite matrix and therefore a lot of **freedom** to place the finite sections.

The previous examples have shown that sometimes one **needs** to make use of that freedom by

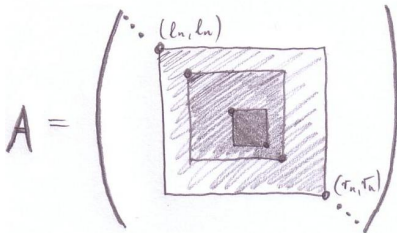
- picking the appropriate “main” diagonal (which one is it?),
- choosing good cut-off sequences (l_n) and (r_n)

We will learn how to do both of that.

The Finite Section Method

We need theorems that tell us when the FSM works and when not.

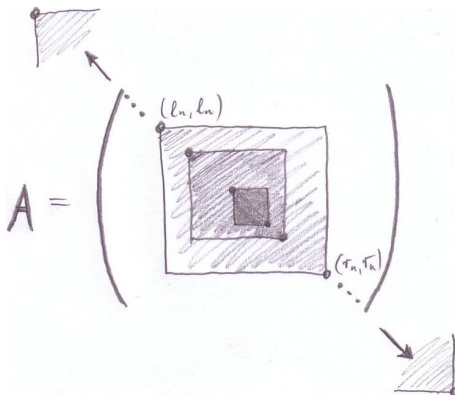
One can show that applicability of the finite section method



The Finite Section Method

We need theorems that tell us when the FSM works and when not.

One can show that applicability of the finite section method



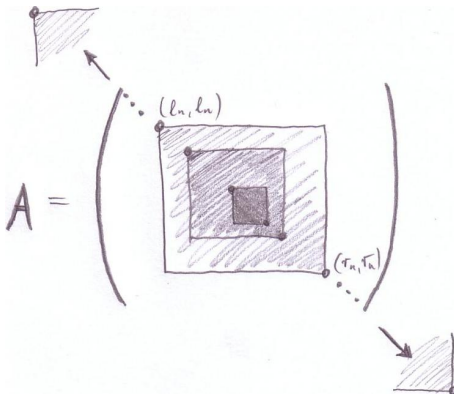
is controlled by certain **limits** of the upper left and lower right **corners** of the finite sections A_n as $n \rightarrow \infty$.

Following the corners as they move out to infinity

So we have to follow the two “corners” (semi-infinite matrices)

$$\begin{pmatrix} a_{l_n, l_n} & \cdots \\ \vdots & \ddots \end{pmatrix} \quad \text{and} \quad \begin{pmatrix} \ddots & \vdots \\ \cdots & a_{r_n, r_n} \end{pmatrix}$$

of A_n as $n \rightarrow \infty$ and find (partial) limits of these matrix sequences:

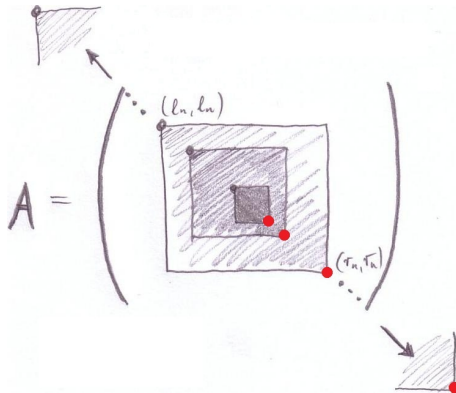


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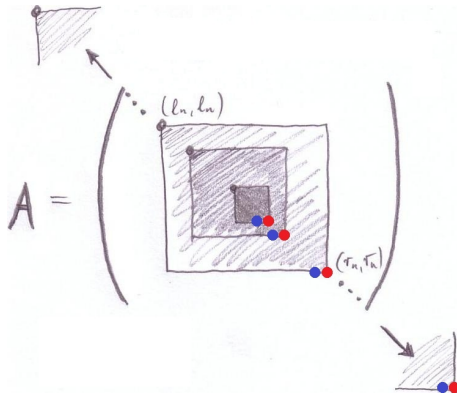


Following the corners as they move out to infinity

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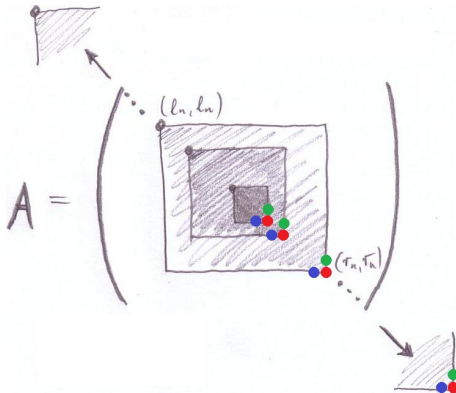


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of A_n as $n \rightarrow \infty$ and find (partial) limits of these matrix sequences:



Limit Operators: Definition

This leads to the study of so-called limit operators.

Definition: Limit Operator

For a given sequence $h_1, h_2, \dots \in \mathbb{Z}$ with $|h_n| \rightarrow \infty$ and a matrix $A = (a_{ij})_{i,j \in \mathbb{Z}}$, we call $B = (b_{ij})_{i,j \in \mathbb{Z}}$ a **limit operator** of A with respect to that sequence $h = (h_1, h_2, \dots)$ if for all $i, j \in \mathbb{Z}$,

$$a_{i+h_n, j+h_n} \rightarrow b_{ij} \quad \text{as} \quad n \rightarrow \infty.$$

We write A_h instead of B , where $h = (h_1, h_2, \dots)$.

For our FSM, we will use $l = (l_1, l_2, \dots)$ and $r = (r_1, r_2, \dots)$ (or subsequences of those) in place of $h = (h_1, h_2, \dots)$.

One more notation: A_+ and A_-

Think of a bi-infinite band matrix A as 2×2 block matrix:

A hand-drawn diagram illustrating a bi-infinite band matrix A as a 2×2 block matrix. On the left, a large letter A is written. To its right is an equals sign followed by a large pair of parentheses. Inside the parentheses, a horizontal line and a vertical line intersect at the center, dividing the space into four quadrants. The top-left quadrant contains the letter A with a minus sign ($-$) next to it. The bottom-right quadrant contains the letter A with a plus sign ($+$) next to it. A diagonal band of small, shaded rectangular blocks runs from the top-left to the bottom-right, passing through the intersection of the lines. The top-right and bottom-left quadrants are empty.

So A_+ and A_- are one-sided infinite submatrices of A .

The Finite Section Method

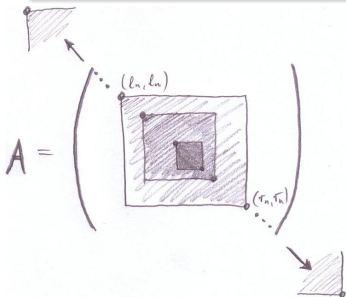
Theorem

ML, Roch 2010; Seidel, Silbermann 2011

The finite section method (2) is applicable to A **iff** the following operators are invertible:

$$A, \quad B_-, \quad C_+$$

for **all** limit operators B of A w.r.t. a subsequence of r
and **all** limit operators C of A w.r.t. a subsequence of l .



Strategy:

Choose the sequences

$r = (r_1, r_2, \dots)$ and $l = (l_1, l_2, \dots)$

so that B_- and C_+ are **invertible!**

The FSM for one-sided infinite matrices

For a banded and one-sided infinite matrix $A = A_+ = (a_{ij})_{i,j \in \mathbb{N}}$, acting boundedly on $\ell^p(\mathbb{N})$, the situation is similar:

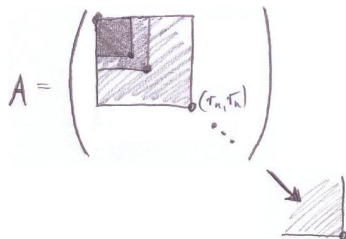
Now $l_n \equiv 1$ is fixed and $r_1 < r_2 < \dots \rightarrow +\infty$.

Theorem

The FSM (2) is applicable **iff**

$$A \quad \text{and} \quad B_-$$

are invertible for **all** limops B of A w.r.t. a subsequence of r .



Strategy: Again, make sure B_- is/are invertible – by choosing the sequence $r = (r_1, r_2, \dots)$.

Ok, limit operators seem to be useful.

Time to learn more about them!

(We come back to the FSM at some later point.)

- 1 Classes of Infinite Matrices
- 2 The Finite Section Method, Part I
- 3 Limit Operators**
- 4 The Spectrum: Formulas and Bounds
- 5 Spectral Bounds: An Example
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Limit Operators: Definition

Let $A \in BDO(E)$. Recall:

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$$a_{i+h_n, j+h_n} \rightarrow b_{ij} \quad \text{as} \quad n \rightarrow \infty.$$

We write A_h instead of B , where $h = (h_1, h_2, \dots)$.

In short: A_h is the entrywise limit of $V_{-h_n} A V_{h_n}$ as $n \rightarrow \infty$.

By $\sigma^{\text{op}}(A)$ we denote the set of all limit operators of A .

Example: Discrete Schrödinger operator

The Schrödinger operator $-\Delta + M_b$ with a bounded potential $b \in L^\infty(\mathbb{R})$ is usually discretized as

$$A = V_{-1} + M_c + V_1 = \begin{pmatrix} \ddots & \ddots & & & & & & & \\ & \ddots & c_{-2} & 1 & & & & & \\ & & 1 & c_{-1} & 1 & & & & \\ & & & 1 & c_0 & 1 & & & \\ & & & & 1 & c_1 & 1 & & \\ & & & & & 1 & c_2 & \ddots & \\ & & & & & & \ddots & \ddots & \end{pmatrix}$$

where $c = (\dots, c_{-1}, c_0, c_1, \dots) \in \ell^\infty(\mathbb{Z})$.

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where $c = (\dots, c_{-1}, c_0, c_1, \dots) \in \ell^\infty(\mathbb{Z})$. Clearly,

$$A_h = (V_{-1})_h + (M_c)_h + (V_1)_h = V_{-1} + (M_c)_h + V_1,$$

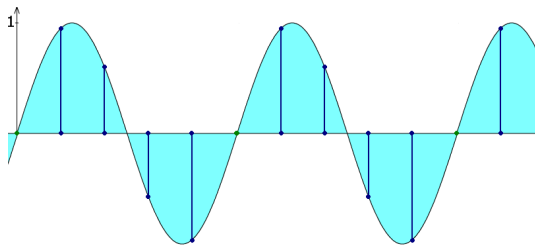
so that everything depends on the limit operators of M_c only.

Example 1: Periodic potential

If there is a $P \in \mathbb{N}$ such that

$$c_{k+P} = c_k \quad \text{for every } k \in \mathbb{Z},$$

then $\sigma^{\text{op}}(M_c) = \left\{ M_{V_k c} : k \in \{0, 1, \dots, P-1\} \right\}$.

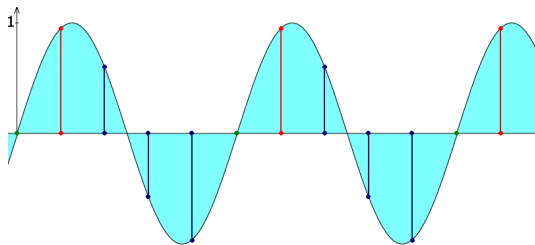


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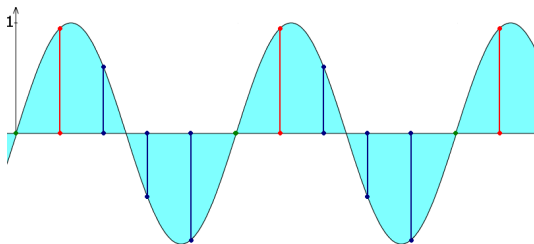
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$$c_{k+P} = c_k \quad \text{for every} \quad k \in \mathbb{Z},$$

then $\sigma^{\text{op}}(M_C) = \left\{ M_{V_{kC}} : k \in \{0, 1, \dots, P-1\} \right\}$.



But then $\sigma^{\text{op}}(A) = \left\{ V_{-k} A V_k : k \in \{0, 1, \dots, P-1\} \right\}$.

Example 2: Almost-periodic potential

Note: $c \in \ell^\infty$ is **periodic** iff the set $\{V_k c : k \in \mathbb{Z}\}$ of all its translates is **finite**.

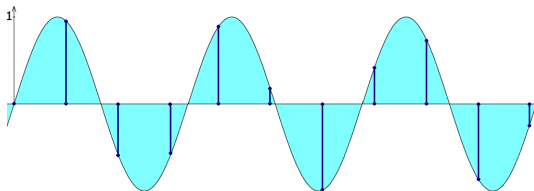
E.g. if $c = (\dots, c_1, c_2, c_3, c_1, c_2, c_3, \dots)$ then

$$\{V_k c : k \in \mathbb{Z}\} = \left\{ \begin{array}{l} V_0 c = (\dots, c_1, c_2, c_3, c_1, c_2, c_3, \dots), \\ V_1 c = (\dots, c_3, c_1, c_2, c_3, c_1, c_2, \dots), \\ V_2 c = (\dots, c_2, c_3, c_1, c_2, c_3, c_1, \dots) \end{array} \right\}$$

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Note: $c \in \ell^\infty$ is **periodic** iff the set $\{V_k c : k \in \mathbb{Z}\}$ of all its translates is **finite**.

Definition: $c \in \ell^\infty$ is **almost-periodic** iff the set $\{V_k c : k \in \mathbb{Z}\}$ of all its translates is **relatively compact** in ℓ^∞ .



The set $h(c) := \text{clos}\{V_k c : k \in \mathbb{Z}\} \subset \ell^\infty$ is called the **hull** of c .

In that case $\sigma^{\text{op}}(M_c) = \{M_d : d \in h(c)\} = \text{clos}\{M_{V_k c} : k \in \mathbb{Z}\}$,
 $\sigma^{\text{op}}(A) = \{V_{-1} + M_d + V_1 : d \in h(c)\} = \text{clos}\{V_{-k} A V_k : k \in \mathbb{Z}\}$.

Example: Discrete Schrödinger operator

Example 2: Almost-periodic potential (continued)

Fix $\lambda \in \mathbb{R}$ and look at $c = (c_k)_{k \in \mathbb{Z}} \in \ell^\infty$ with entries

$$c_k := e^{i\lambda k}, \quad k \in \mathbb{Z}.$$

If λ/π is rational then c is periodic;
otherwise it is almost-periodic and its hull is

$$h(c) = \left\{ d = \left(e^{i(\lambda k + \gamma)} \right)_{k \in \mathbb{Z}} : \gamma \in [0, 2\pi) \right\}.$$

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Slightly more advanced: If

$$c_k := \alpha e^{i\lambda k} + \beta e^{i\mu k}, \quad k \in \mathbb{Z}$$

with λ/π , μ/π and λ/μ all irrational then the hull is

$$h(c) = \left\{ d = \left(\alpha e^{i(\lambda k + \gamma)} + \beta e^{i(\mu k + \delta)} \right)_{k \in \mathbb{Z}} : \gamma, \delta \in [0, 2\pi) \right\}.$$



Example 2: Almost-periodic potential (continued 2)

For the so-called Almost-Mathieu operator, the potential c is the real part of what we studied earlier:

Fix $\alpha, \lambda \in \mathbb{R}$ and take $c = (c_k)_{k \in \mathbb{Z}} \in \ell^\infty$ with entries

$$c_k := \alpha \cos(\lambda k), \quad k \in \mathbb{Z}.$$

If λ/π is irrational then c is almost-periodic (but non-periodic) and its hull is

$$h(c) = \left\{ d = (\alpha \cos(\lambda k + \gamma))_{k \in \mathbb{Z}} : \gamma \in [0, 2\pi) \right\}.$$

Ten-Martini problem: Show that $\text{spec}(V_{-1} + M_c + V_1)$ is a Cantor set. (Puig 2003, Avila&Jitomiskaya 2005)

Example 3: Slowly oscillating potential

If

$$c_{k+1} - c_k \rightarrow 0 \quad \text{as} \quad k \rightarrow \pm\infty,$$

then $\sigma^{\text{op}}(M_c) = \{aI : a \in c(\infty)\}$.

Now all limit operators A_h are Laurent operators (i.e., they have constant diagonals).

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Now all limit operators A_h are Laurent operators (i.e., they have constant diagonals).

For example, let $c = (c_k)_{k \in \mathbb{Z}}$ with

$$c_k = \sin \sqrt{|k|}, \quad k \in \mathbb{Z}.$$

Then c is slowly oscillating and $c(\infty) = [-1, 1]$.

Example 4: Pseudo-ergodic potential

Let Σ be a compact subset of \mathbb{C} .

Definition

Davies 2001

A sequence $c = (c_k)_{k \in \mathbb{Z}}$ over Σ is called **pseudoergodic** over Σ if every finite vector $f = (f_i)_{i \in F}$ with values $f_i \in \Sigma$ can be found, up to arbitrary precision $\varepsilon > 0$, somewhere inside the infinite sequence c , i.e.

$$\forall \varepsilon > 0 \quad \exists m \in \mathbb{Z} : \max_{i \in F} |c_{i+m} - f_i| < \varepsilon.$$

Idea: Model random behaviour by a purely deterministic concept.

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Idea: Model random behaviour by a purely deterministic concept.

Example: $dec(\pi) = 3, 1, 4, 1, 5, 9, 2, 6, 5, 3, 5, \dots$ is pseudo-ergodic over $\Sigma = \{0, 1, 2, \dots, 9\}$ (conjecture)

Example: $bin(1), bin(2), bin(3), \dots$ is pseudo-erg. over $\Sigma = \{0, 1\}$.

Example 4: Pseudo-ergodic potential (continued)

If c is pseudo-ergodic over Σ then **every** multiplication operator M_d with a sequence $d = (\dots, d_{-1}, d_0, d_1, \dots)$ over Σ is a limit operator of M_c :

$$\sigma^{\text{op}}(M_c) = \{ M_d : d : \mathbb{Z} \rightarrow \Sigma \} \quad (3)$$

The statement also holds the other way round!

So c is pseudo-ergodic **iff** (3) holds.

Example 5: A locally constant potential

$$c = (\dots, \underbrace{\beta, \beta, \beta, \beta}_4, \underbrace{\alpha, \alpha, \alpha}_3, \underbrace{\beta, \beta}_2, \underbrace{\alpha}_1, \underbrace{\beta, \beta}_2, \underbrace{\alpha, \alpha, \alpha}_3, \underbrace{\beta, \beta, \beta, \beta}_4, \dots).$$

Example 5: A locally constant potential

$$c = (\dots, \underbrace{\beta, \beta, \beta, \beta}_4, \underbrace{\alpha, \alpha, \alpha}_3, \underbrace{\beta, \beta}_2, \underbrace{\alpha}_1, \underbrace{\beta, \beta}_2, \underbrace{\alpha, \alpha, \alpha}_3, \underbrace{\beta, \beta, \beta, \beta}_4, \dots).$$

Then all limit operators of A are of the form

$$\left(\begin{array}{cccc} \ddots & \ddots & & \\ \ddots & \beta & 1 & \\ & 1 & \beta & \ddots \\ & & \ddots & \ddots \end{array} \right), \quad \left(\begin{array}{cccc} \ddots & \ddots & & \\ \ddots & \alpha & 1 & \\ & 1 & \alpha & \ddots \\ & & \ddots & \ddots \end{array} \right),$$

Limit Operators: Some rules

Let $A, B \in BDO(E)$ and $h = (h_1, h_2, \dots)$ be a sequence of integers going to $\pm\infty$.

Recall: limit op A_h is the entrywise limit of $V_{-h_n} A V_{h_n}$ as $n \rightarrow \infty$

Basic rules

If the right-hand side exists then also the left-hand side exists and equality holds:

$$(A + B)_h = A_h + B_h$$

$$(AB)_h = A_h B_h$$

$$(\alpha A)_h = \alpha A_h$$

$$(\lim A_n)_h = \lim (A_n)_h \quad (\text{w.r.t. op-}\|\cdot\|)$$

\Rightarrow Compute limit operators of elements of an operator algebra in terms of limit operators of generators of the algebra.

Moreover, $\|A_h\| \leq \|A\|$.

The Spectrum: Formulas and Bounds

- 1 Classes of Infinite Matrices
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Besides studying the **FSM**, limit operators are also important for determining the **essential spectrum** of A .

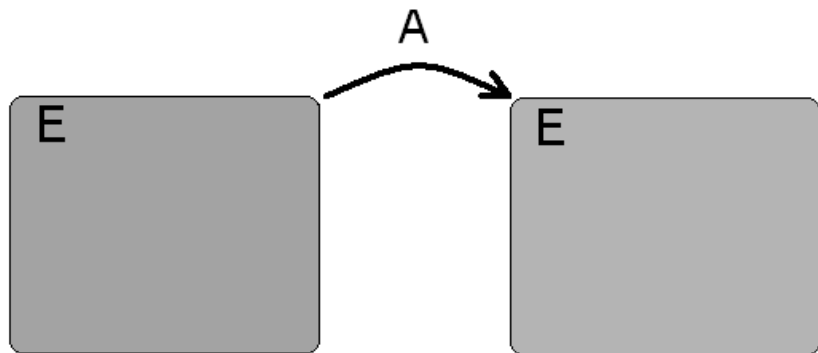
Here we define

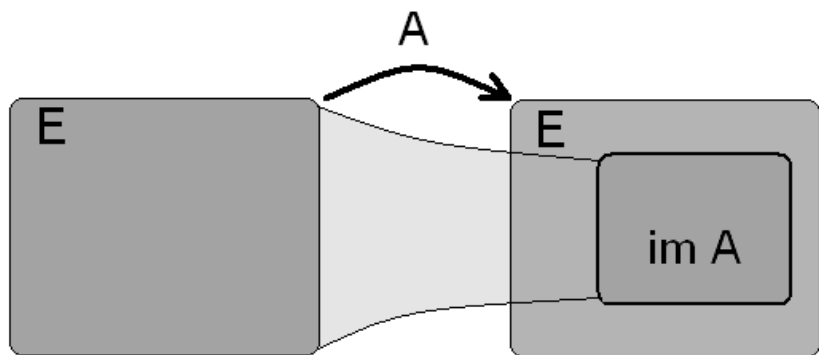
Definition: Essential Spectrum

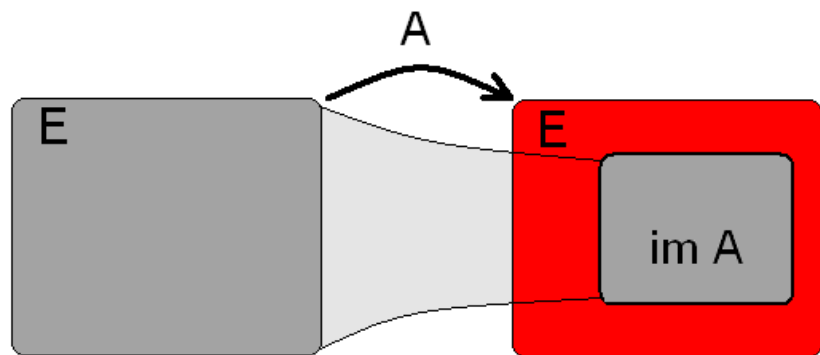
For a (not necessarily self-adjoint) operator A , we denote by

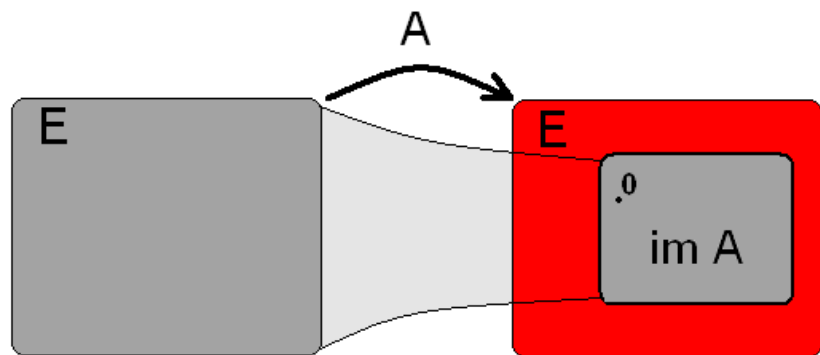
$$\text{spec}_{\text{ess}} A := \{ \lambda \in \mathbb{C} : A - \lambda I \text{ is not Fredholm} \}$$

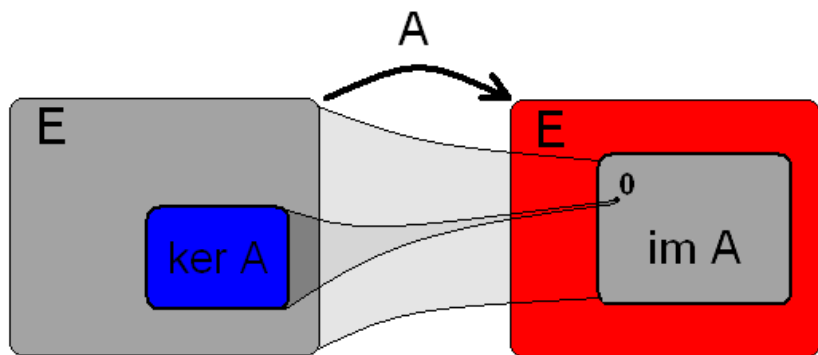
the **essential spectrum** of A .

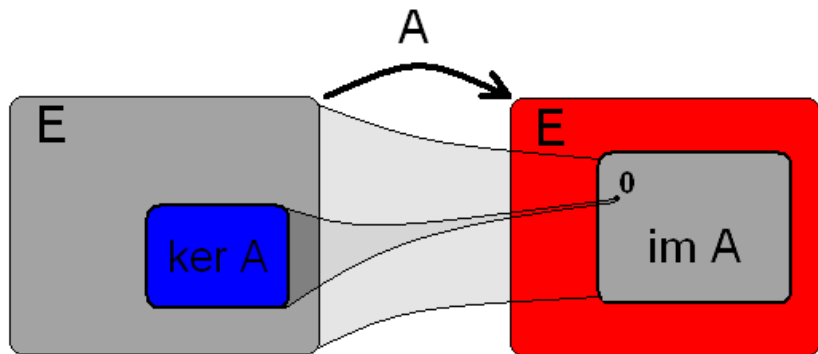










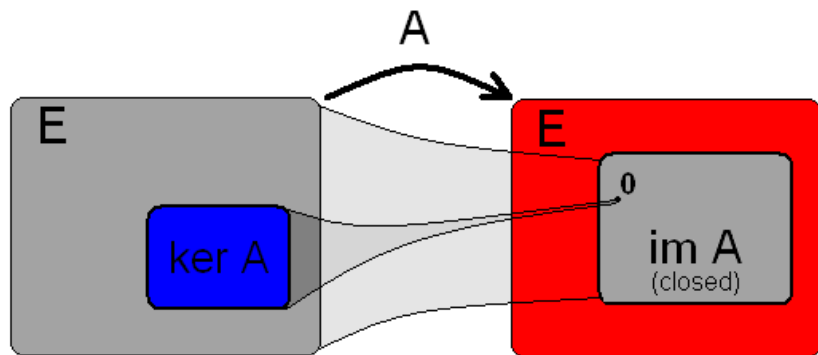


Definition

$A : E \rightarrow E$ is a **Fredholm operator**

if $\alpha := \dim(\ker A)$ and $\beta := \text{codim}(\text{im } A)$ are both finite.

The difference $\alpha - \beta$ is then called the **index of A** .

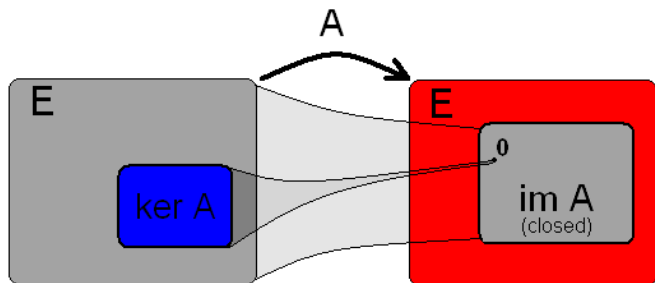


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if $\alpha := \dim(\ker A)$ and $\beta := \operatorname{codim}(\operatorname{im} A)$ are both finite.

The difference $\alpha - \beta$ is then called the **index of A** .

$A \in L(E)$ is Fredholm iff $A + K(E)$ is invertible in $L(E)/K(E)$.

Limit Operators vs. Fredholmness

Take $A \in \mathcal{W}(E)$.

Then it is not hard to see that

A **Fredholm** \implies all limit operators of A are **invertible**.

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If we repeat the same argument with $A - \lambda I$ in place of A , we get:

Essential Spectrum

Rabinovich, Roch, Silbermann 1998; ML 2003; Chandler-Wilde, ML 2007

$$\text{spec}_{\text{ess}}^p(A) = \bigcup_h \text{spec}^p(A_h) = \bigcup_h \text{spec}_{\text{point}}^\infty(A_h), \quad p \in [1, \infty]$$

Limit operators vs. Fredholm index

Think of a bi-infinite band matrix A as 2×2 block matrix:

$$A = \begin{pmatrix} A_- & \\ & A_+ \end{pmatrix}$$

Then

$$\begin{aligned} A \text{ is Fredholm} & \text{ iff } A_+ \text{ and } A_- \text{ are Fredholm} \\ \text{spec}_{\text{ess}} A &= \text{spec}_{\text{ess}} A_+ \cup \text{spec}_{\text{ess}} A_- \\ \text{ind } A &= \text{ind } A_+ + \text{ind } A_- \end{aligned}$$

Limit operators vs. Fredholm index

$$C = \left(\begin{array}{c} \\ \\ \end{array} \right) \text{invertible}$$

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$$C = \begin{pmatrix} + \\ + \end{pmatrix} \text{ invertible}$$

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\mathcal{K}_+

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$$C = \begin{pmatrix} \color{red}{\kappa_-} & \\ & \color{blue}{\kappa_+} \end{pmatrix} \text{ invertible}$$

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$$C = \begin{pmatrix} k_- & \\ & -k_- \end{pmatrix} \text{ invertible}$$

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Theorem

Rabinovich, Roch, Roe 2004

$\text{ind } A_+ = \text{ind } B_+$ for all limops B of A at $+\infty$

$\text{ind } A_- = \text{ind } C_-$ for all limops C of A at $-\infty$

Example: Random Jacobi Operator

So, if the random (meaning pseudoergodic) operator A is Fredholm then, for **all** choices $u \in U$, $v \in V$ and $w \in W$:

$$\text{ind } A_- = \text{ind} \begin{pmatrix} \ddots & \ddots & & & \\ \ddots & \ddots & & & \\ & & & & w \\ & & & u & v \\ & & & & \end{pmatrix}$$

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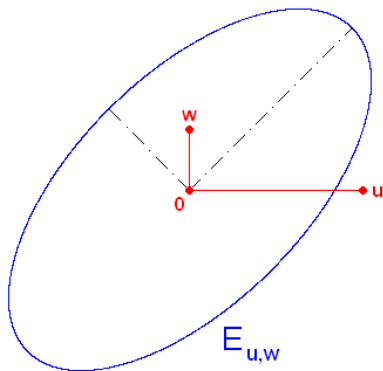
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The previous slide has shown how 'hard' it is for a pseudoergodic Jacobi operator to be Fredholm. Let us underline this. Put

$$\begin{aligned} J(U, V, W) &:= \{ \text{Jacobi ops (4)} : u_i \in U, v_i \in V, w_i \in W \}, \\ \Psi E(U, V, W) &:= \{ A \in J(U, V, W) : A \text{ pseudoergodic} \}. \end{aligned}$$

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In particular,

$$\text{spec} A \supseteq \bigcup_{\text{Laurent}} \text{spec} L = \bigcup_{u,v,w} (v + E_{u,w}).$$

Example 1: Periodic potential

If

$$c_{k+P} = c_k \quad \text{for every } k \in \mathbb{Z},$$

then $\sigma^{\text{op}}(M_C) = \left\{ M_{V_k c} : k \in \{0, 1, \dots, P-1\} \right\}$.

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$$\text{But then } \sigma^{\text{op}}(A) = \left\{ V_{-k} A V_k : k \in \{0, 1, \dots, P-1\} \right\}.$$

Consequently, A is invertible iff any/all of its limit operators are invertible. So in this case, A is **Fredholm** iff it is **invertible**, and

$$\text{spec}_{\text{ess}} A = \text{spec } A = \text{spec}_{\text{point}}^{\infty} A$$

Example 2: Almost-periodic potential

Let c be almost-periodic with hull $h(c) = \text{clos}\{V_k c : k \in \mathbb{Z}\}$.

$$\sigma^{\text{op}}(M_c) = \{M_d : d \in h(c)\} = \text{clos}\{M_{V_k c} : k \in \mathbb{Z}\},$$

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If any $A_h = \lim V_{-h_n} A V_{h_n}$ is invertible then $V_{-h_n} A V_{h_n}$ is invertible for large n , so that A itself is invertible!

Hence

$$A \text{ Fredholm} \iff A \text{ invertible} \iff \text{any } A_h \text{ invertible}$$

$$\text{spec}_{\text{ess}} A = \text{spec } A = \text{spec } A_h = \bigcup_h \text{spec}_{\text{point}}^{\infty} A_h$$

Example 3: Slowly oscillating potential

If

$$c_{i+1} - c_i \rightarrow 0 \quad \text{as} \quad i \rightarrow \pm\infty,$$

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But then all limit operators A_h are Laurent operators (i.e., they have constant diagonals), for which invertibility & spectrum are well-understood.

$$\text{spec}_{\text{ess}} A = \bigcup \text{spec} A_h = c(\infty) + [-2, 2]$$

Summary: Limit operators help us to determine the essential spectrum. So they give us **lower bounds** on the spectrum.

It would be good to also have **upper bounds**!

Classical upper bounds

- Gershgorin circles
- numerical range

We will discuss another approach.

One more notion: Pseudospectrum

For $A \in L(E)$ and $\varepsilon > 0$, we put

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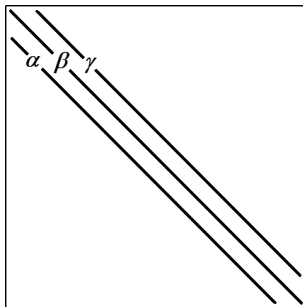
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The sets $\operatorname{spec}_{\varepsilon} A$, $\varepsilon > 0$, are the so-called ε -**pseudospectra** of A . It holds that

$$\operatorname{spec} A =: \operatorname{spec}_0 A \subset \operatorname{spec}_{\varepsilon_1} A \subset \operatorname{spec}_{\varepsilon_2} A, \quad 0 < \varepsilon_1 < \varepsilon_2.$$

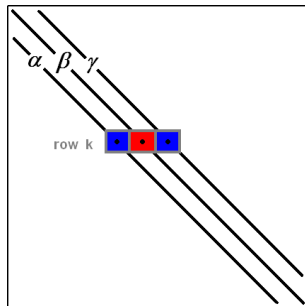
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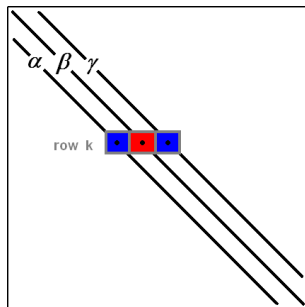


For every row k , consider the **disk** with

center at $a_{k,k}$ and radius $|a_{k,k-1}| + |a_{k,k+1}|$

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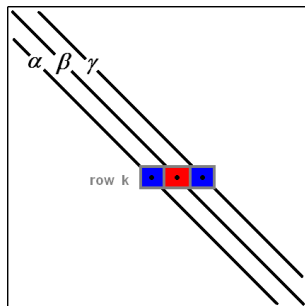
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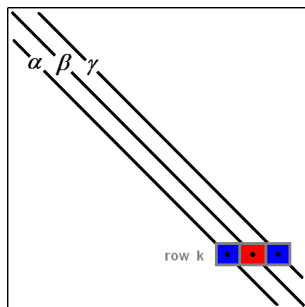
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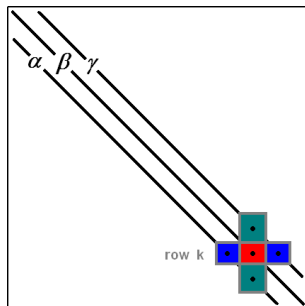
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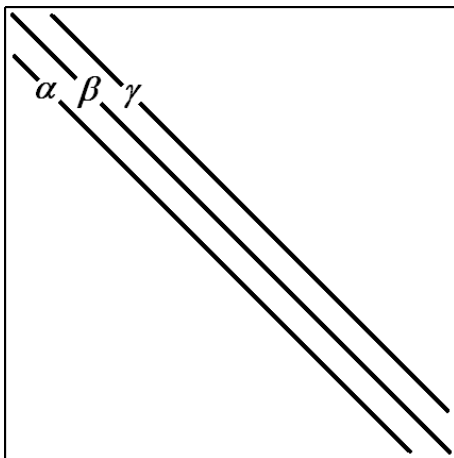
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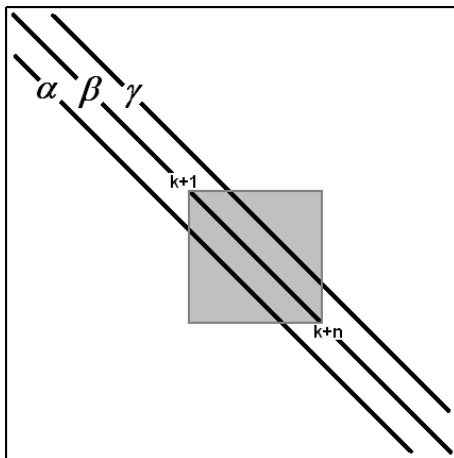
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Look at (pseudo)spectra of the **finite principal submatrices** of A :



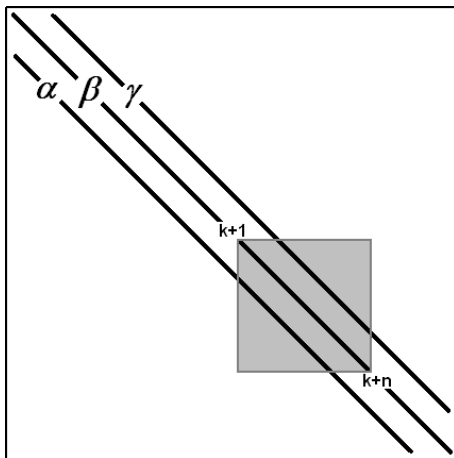
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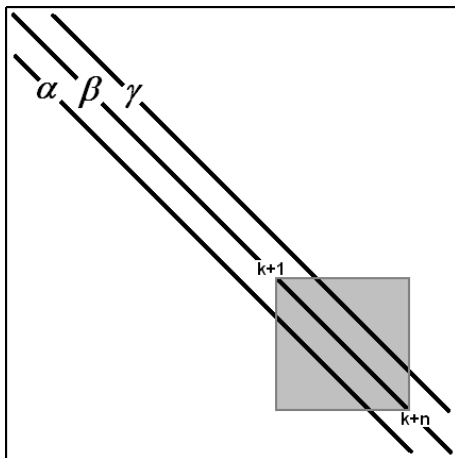
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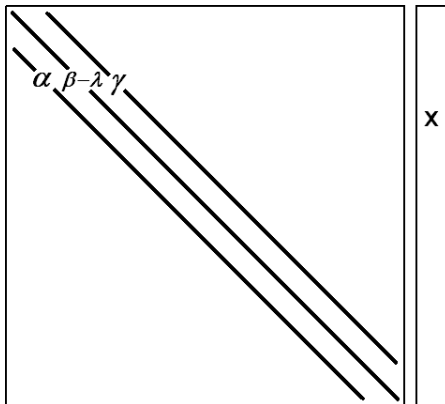
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Method 1: Finite principal submatrices

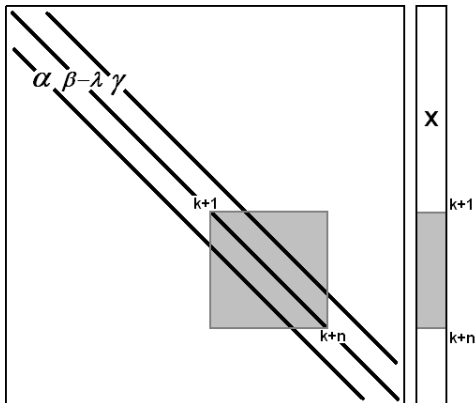
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$$\|(A - \lambda I)x\| < \varepsilon \|x\|$$

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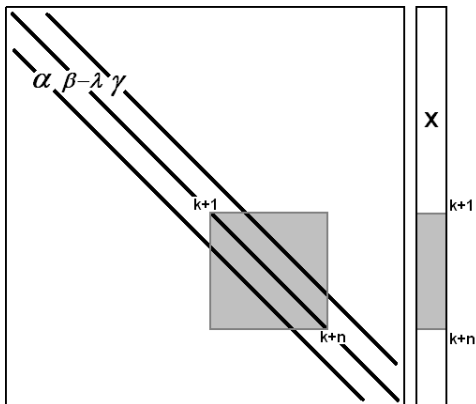
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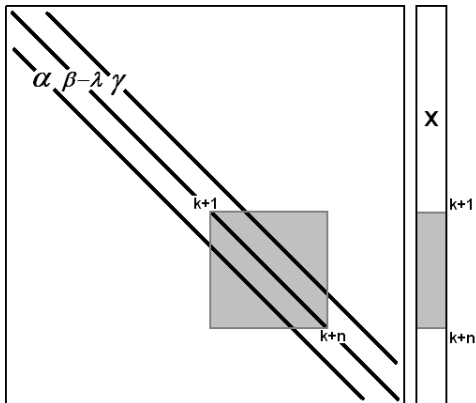
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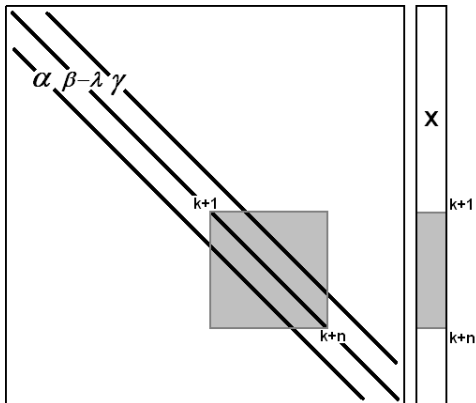
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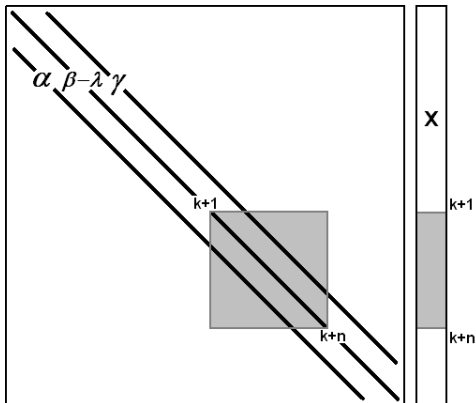
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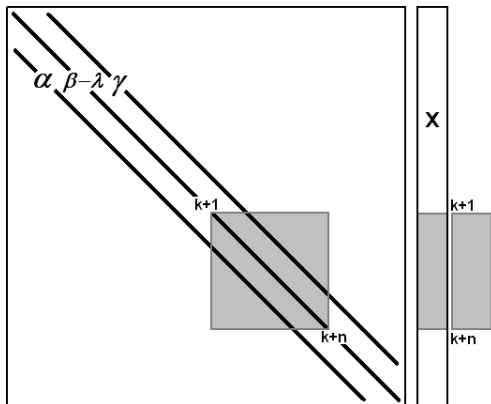


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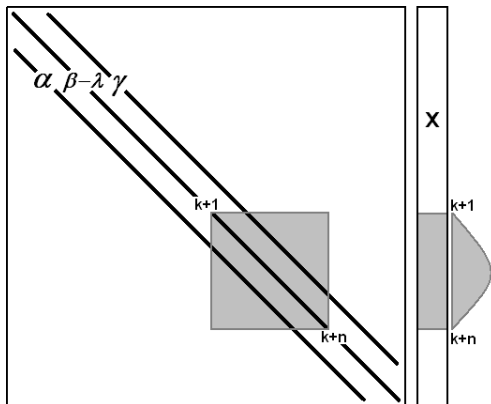
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$$\varepsilon_n < \frac{1}{\sqrt{n}} (\|\alpha\|_\infty + \|\gamma\|_\infty)$$

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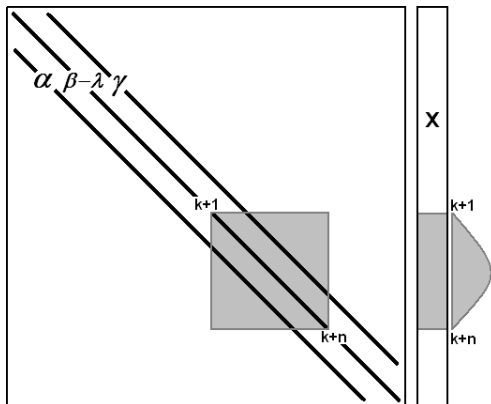
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Method 1: Finite principal submatrices

Let $\lambda \in \text{spec}_\varepsilon(A)$ and let $x \in \ell^2$ be a corresponding pseudomode.



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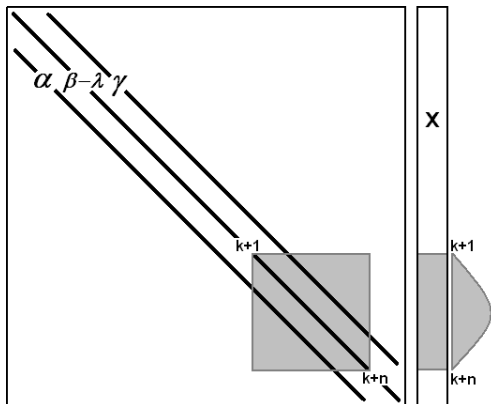
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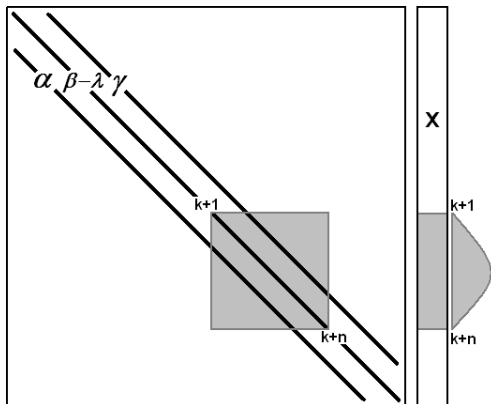
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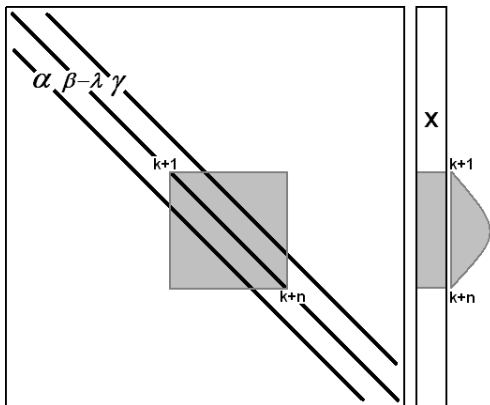
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So one gets

Upper Bound

$$\text{spec}_\varepsilon(A) \subset \bigcup_{k \in \mathbb{Z}} \text{spec}_{\varepsilon + \varepsilon_n}(A_{n,k}),$$

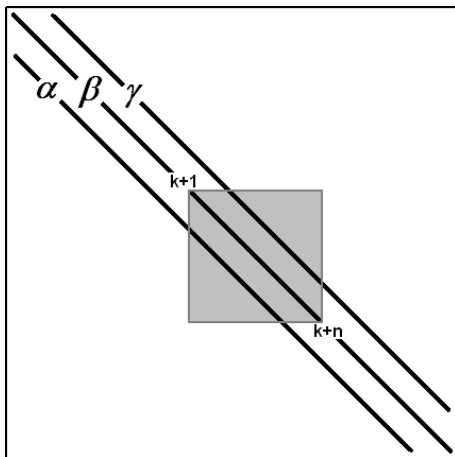
where

$$\varepsilon_n < \frac{\pi}{n} (\|\alpha\|_\infty + \|\gamma\|_\infty).$$

In particular, $\varepsilon_n \rightarrow 0$ as $n \rightarrow \infty$.

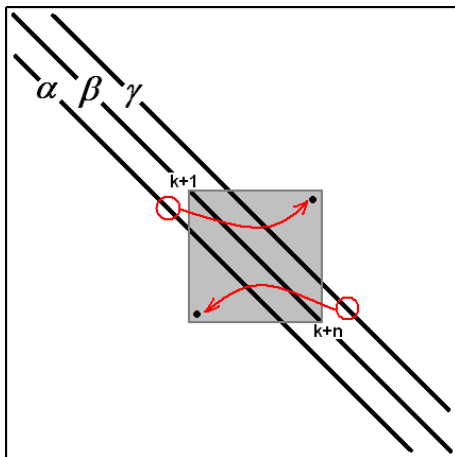
Method 2: **Periodised** finite principal submatrices

If the finite submatrices $A_{n,k}$ are “periodised”,



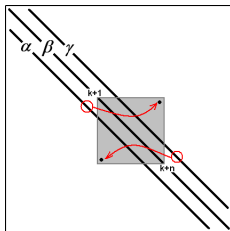
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very similar computations show that, again,

$$\text{spec}_\varepsilon(A) \subset \bigcup_{k \in \mathbb{Z}} \text{spec}_{\varepsilon + \varepsilon_n}(A_{n,k}^{\text{per}})$$

$$\text{with } \varepsilon_n < \frac{\pi}{n} (\|\alpha\|_\infty + \|\gamma\|_\infty)$$

but this upper bound on $\text{spec}_\varepsilon(A)$ generally seems sharper than in method 1.

Method 1 vs. Method 2: An Example

Look at the shift operator

$$(Ax)(i) = x(i+1), \quad \text{i.e.} \quad A = \begin{pmatrix} \cdot & \cdot & \cdot & \cdot & \cdot \\ & 0 & 1 & & \\ & & 0 & 1 & \\ & & & 0 & 1 \\ & & & & \ddots \end{pmatrix}.$$

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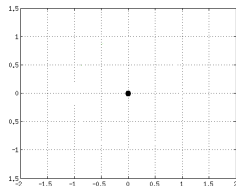
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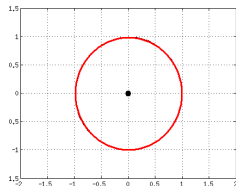
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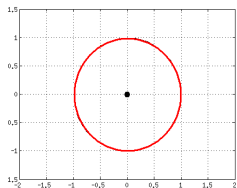
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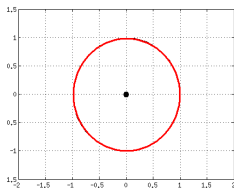
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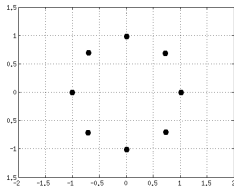
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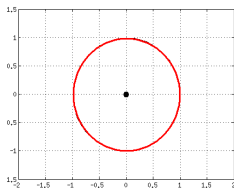
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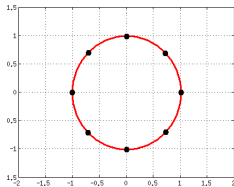
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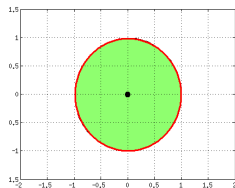
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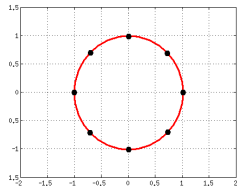
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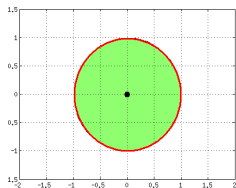
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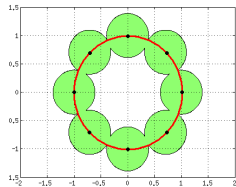
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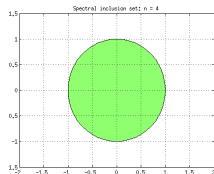
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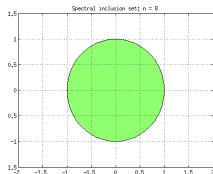
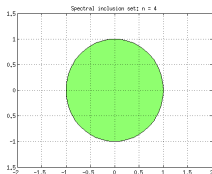
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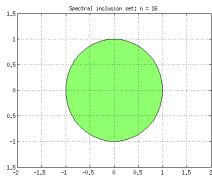
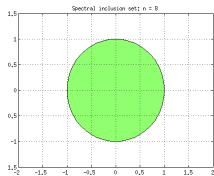
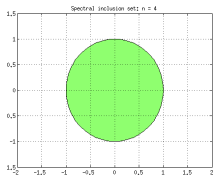
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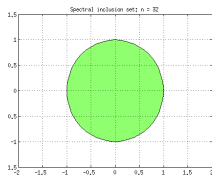
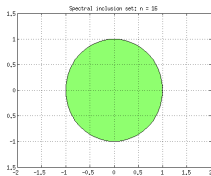
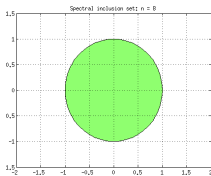
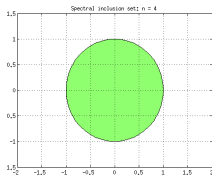
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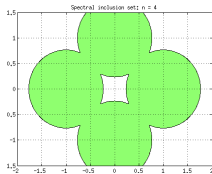
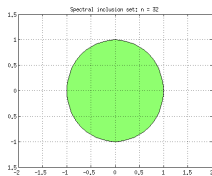
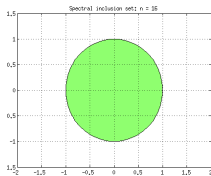
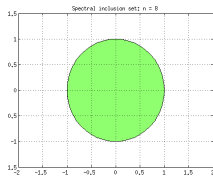
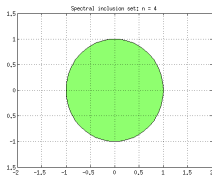
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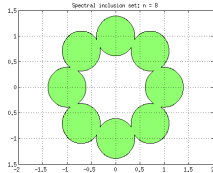
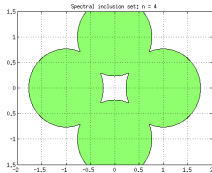
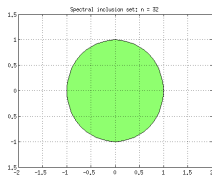
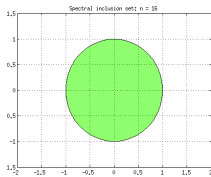
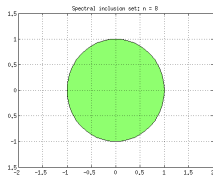
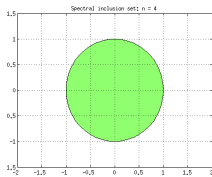
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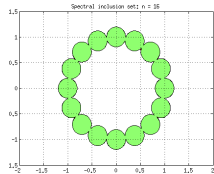
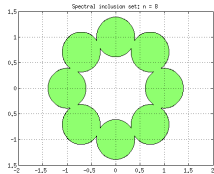
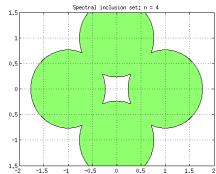
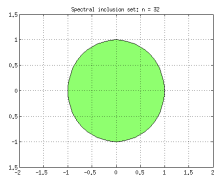
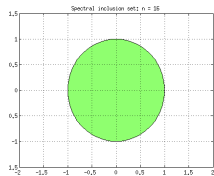
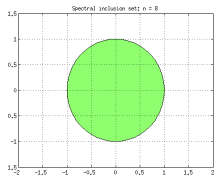
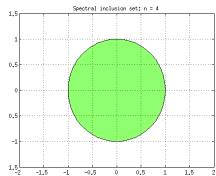
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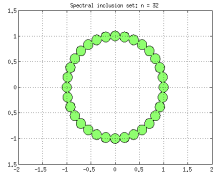
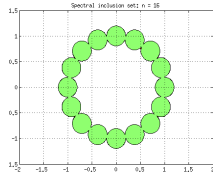
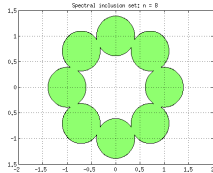
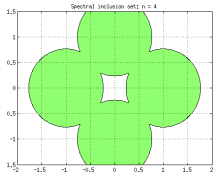
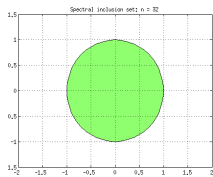
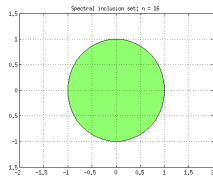
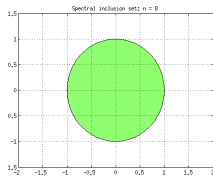
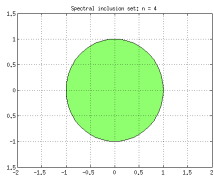
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Method 1 vs. Method 2: An Example

Look at the shift operator

$$(Ax)(i) = x(i+1), \quad \text{i.e.} \quad A = \begin{pmatrix} \cdot & \cdot & \cdot & \cdot & \cdot \\ & 0 & 1 & & \\ & & 0 & 1 & \\ & & & 0 & 1 \\ & & & & \ddots \end{pmatrix}.$$



Summary on Methods 1 & 2

- Both methods give **upper bounds** on $\text{spec } A$ and $\text{spec}_\varepsilon A$.

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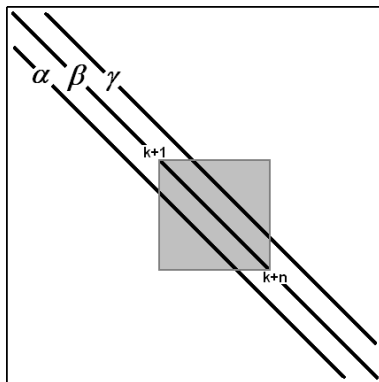
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Summary on Methods 1 & 2

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- The bound from Method 2 always appears to be **sharper**.
- Conjecture: Method 2 **converges** to $\text{spec}_\varepsilon A$ as $n \rightarrow \infty$.
- Method 1 also works for **semi-infinite** and **finite** matrices A !

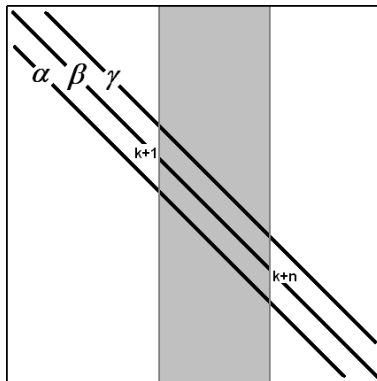
Here is another idea: Method 3

Instead of



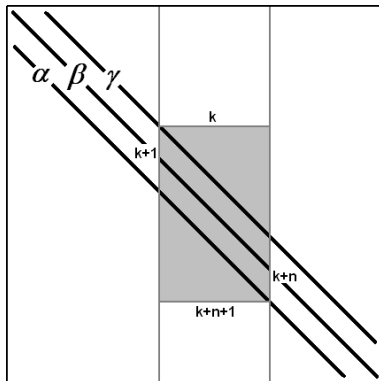
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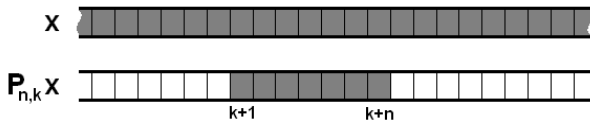
In a sense, we work with **rectangular** finite submatrices.

This is motivated by work of Davies 1998 and Hansen 2008.
(Also see Heinemeyer/ML/Potthast [SIAM Num. Anal. 2007].)

Method 3: Projection Operator

For $n \in \mathbb{N}$ and $k \in \mathbb{Z}$, let $P_{n,k} : \ell^2 \rightarrow \ell^2$ denote the projection

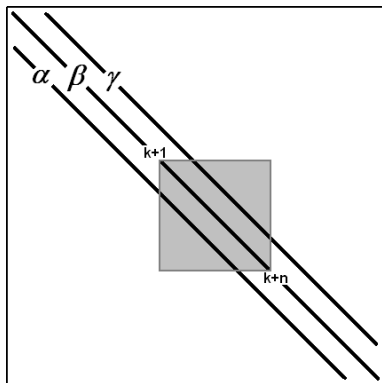
$$(P_{n,k}x)(i) := \begin{cases} x(i), & i \in \{k+1, \dots, k+n\}, \\ 0 & \text{otherwise.} \end{cases}$$



Further, we put $X_{n,k} := \text{im } P_{n,k}$ and identify it with \mathbb{C}^n in the obvious way.

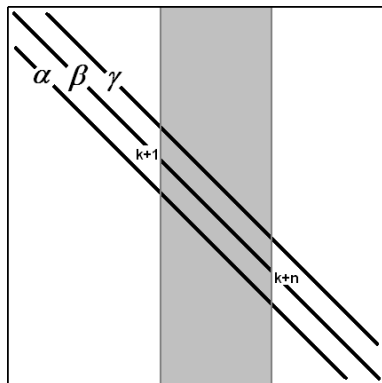
Method 3: Truncations

Method 1:



$$P_{n,k}(A - \lambda I)P_{n,k}|_{X_{n,k}}$$

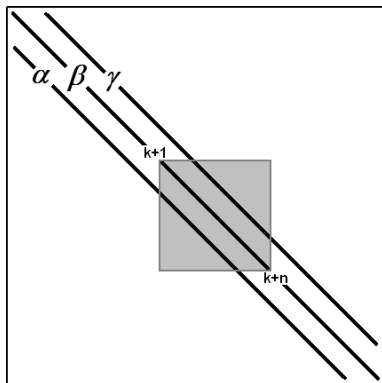
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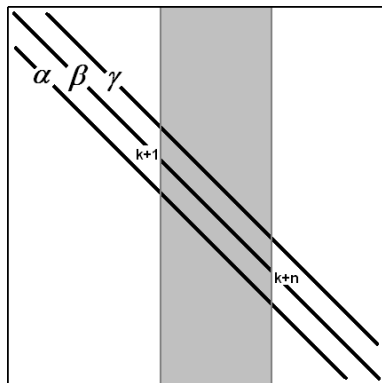
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$\lambda \in \text{spec}_\varepsilon(A) \implies$ For some $k \in \mathbb{Z}$:

$$\lambda \in \text{spec}_{\varepsilon+\varepsilon_n}(P_{n,k}AP_{n,k}|_{X_{n,k}})$$

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and $\min \operatorname{spec} \left(P_{n,k}(A - \lambda I)(A - \lambda I)^*P_{n,k} \right) < (\varepsilon + \varepsilon_n)^2.$

Then put

$$\Gamma_\varepsilon^n(A) := \bigcup_{k \in \mathbb{Z}} \gamma_\varepsilon^{n,k}(A).$$

Method 3: Spectral bounds

Again we get (as in Methods 1 & 2)

Upper Bound

$$\text{spec}_\varepsilon(A) \subset \bigcup_{k \in \mathbb{Z}} \gamma_{\varepsilon + \varepsilon_n}^{n,k}(A) = \Gamma_{\varepsilon + \varepsilon_n}^n(A)$$

$$\text{with } \varepsilon_n < \frac{\pi}{n} (\|\alpha\|_\infty + \|\gamma\|_\infty)$$

and this time the upper bound looks even sharper than before.

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and this time the upper bound looks even sharper than before.
But now we also have

Lower Bound

$$\Gamma_\varepsilon^n(A) \subset \text{spec}_\varepsilon(A).$$

Method 3: Spectral bounds

From the lower and upper bound

$$\Gamma_{\varepsilon}^n(A) \subset \text{spec}_{\varepsilon}(A) \quad \text{and} \quad \text{spec}_{\varepsilon}(A) \subset \Gamma_{\varepsilon+\varepsilon_n}^n(A)$$

we get

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Sandwich 2

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Sandwich 2

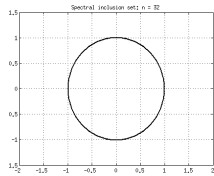
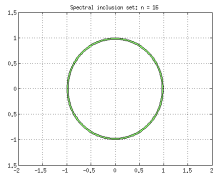
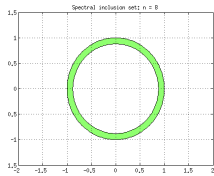
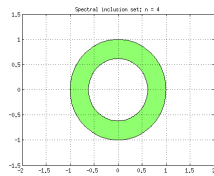
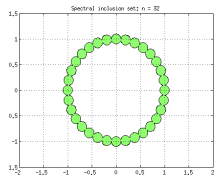
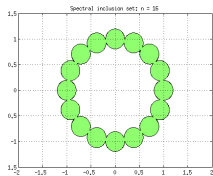
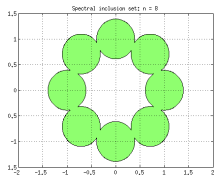
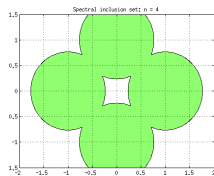
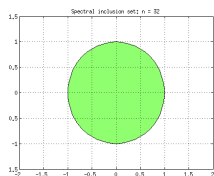
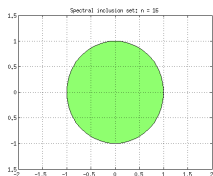
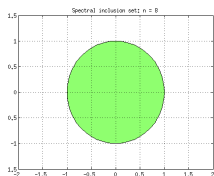
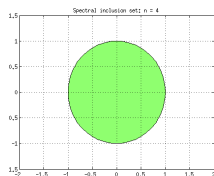
$$\text{spec}_{\varepsilon}(A) \subset \Gamma_{\varepsilon+\varepsilon_n}^n(A) \subset \text{spec}_{\varepsilon+\varepsilon_n}(A).$$

In particular, it follows that

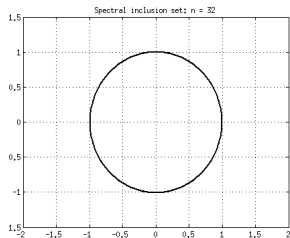
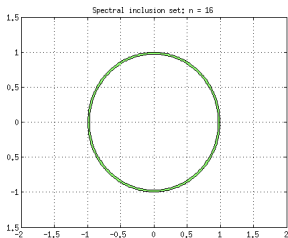
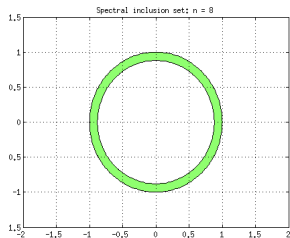
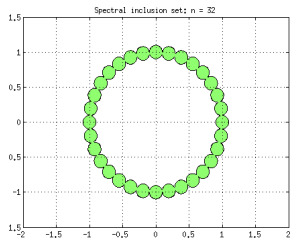
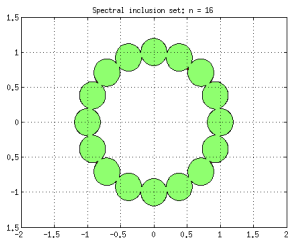
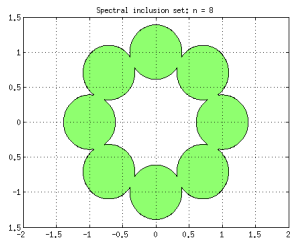
$$\Gamma_{\varepsilon+\varepsilon_n}^n(A) \rightarrow \text{spec}_{\varepsilon}(A), \quad n \rightarrow \infty$$

in the Hausdorff metric.

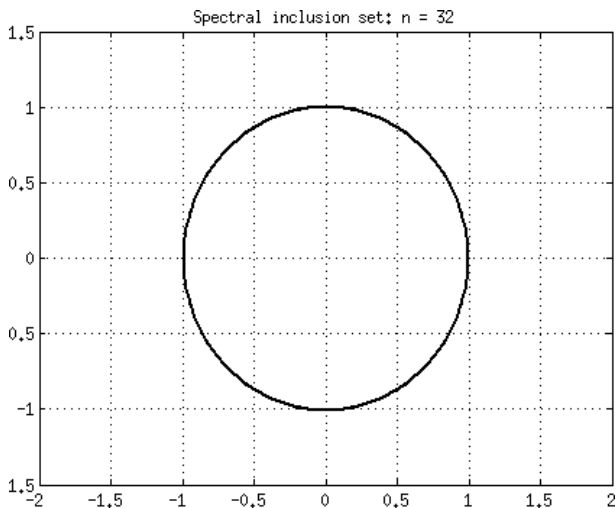
Methods 1, 2 & 3: The Shift Operator



Methods 2 & 3: The Shift Operator



Method 3: The Shift Operator

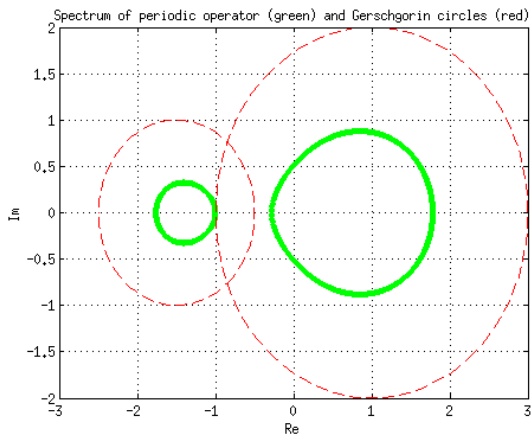


Methods 1, 2 & 3: Second Example

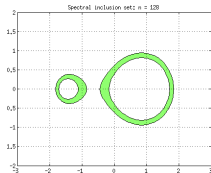
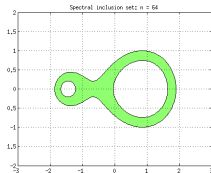
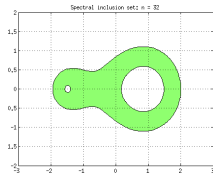
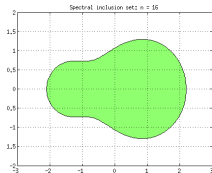
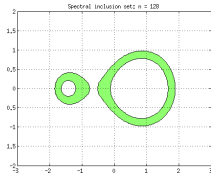
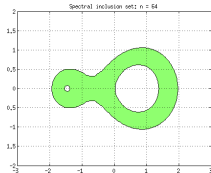
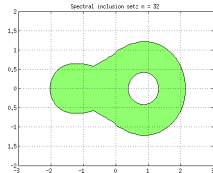
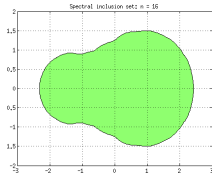
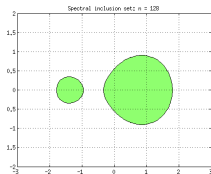
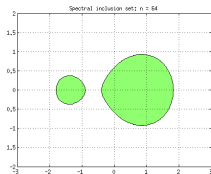
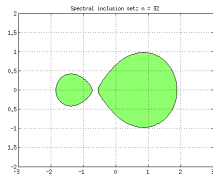
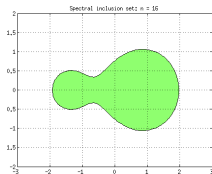
We now look at a matrix A with 3-periodic diagonals:

main diagonal: $\dots, -\frac{3}{2}, 1, 1, \dots$

super-diagonal: $\dots, 1, 2, 1, \dots$



Methods 1, 2 & 3: Second Example

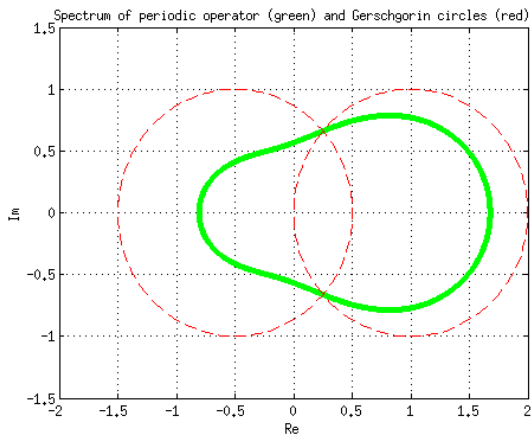


Methods 1, 2 & 3: Third Example

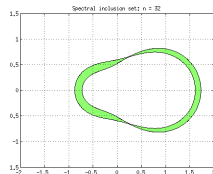
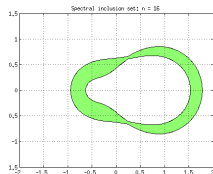
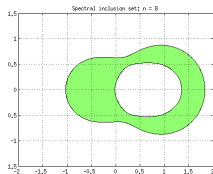
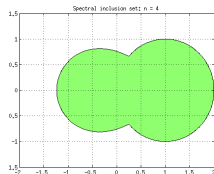
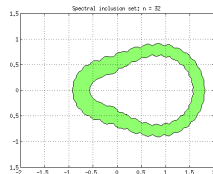
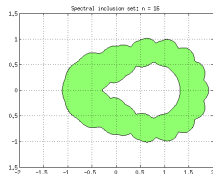
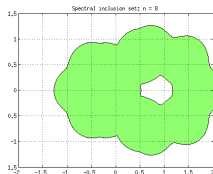
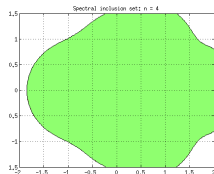
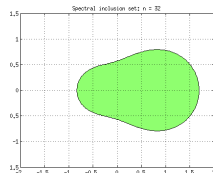
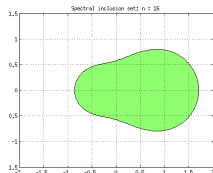
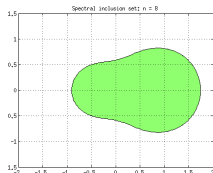
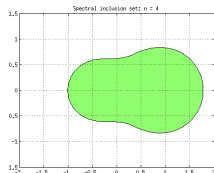
We now look at a matrix A with 3-periodic diagonals:

main diagonal: $\dots, -\frac{1}{2}, 1, 1, \dots$

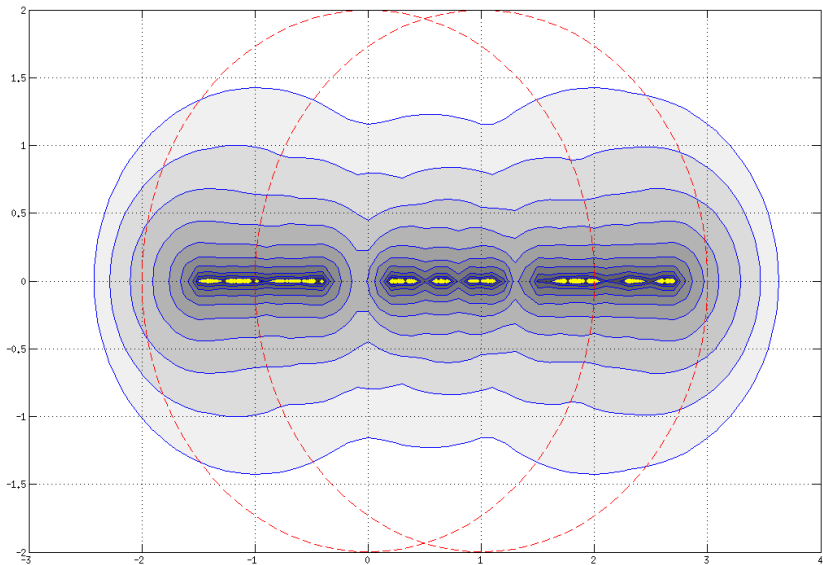
super-diagonal: $\dots, 1, 1, 1, \dots$



Methods 1, 2 & 3: Third Example



Method 3: Schrödinger operator with Cantor spectrum



Spectral Bounds: An Example

- 1 Classes of Infinite Matrices
- 2 The Finite Section Method, Part I
- 3 Limit Operators
- 4 The Spectrum: Formulas and Bounds
- 5 Spectral Bounds: An Example**
- 6 The Finite Section Method, Part II

Spectral Formula

Chandler-Wilde, ML 2007

If b is pseudoergodic then

$$\operatorname{spec} A^b = \operatorname{spec}_{\text{ess}} A^b = \bigcup_{c \in \{\pm 1\}^{\mathbb{Z}}} \operatorname{spec}_{\text{point}}^{\infty} A^c.$$

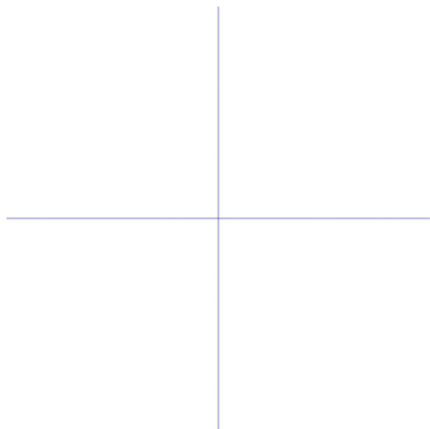
Spectral Formula

Chandler-Wilde, ML 2007

If b is pseudoergodic then

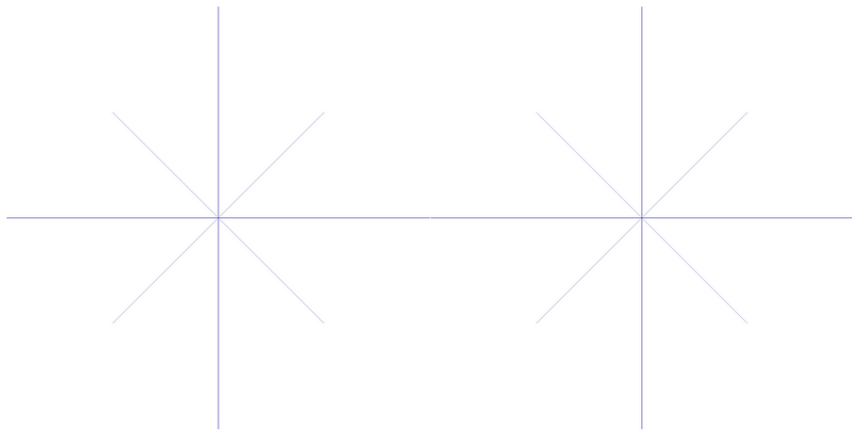
$$\operatorname{spec} A^b = \operatorname{spec}_{\text{ess}} A^b = \bigcup_{c \in \{\pm 1\}^{\mathbb{Z}}} \operatorname{spec}_{\text{point}}^{\infty} A^c.$$

Idea: Try to “exhaust” the RHS by running through all **periodic** ± 1 sequences c .



Period 1

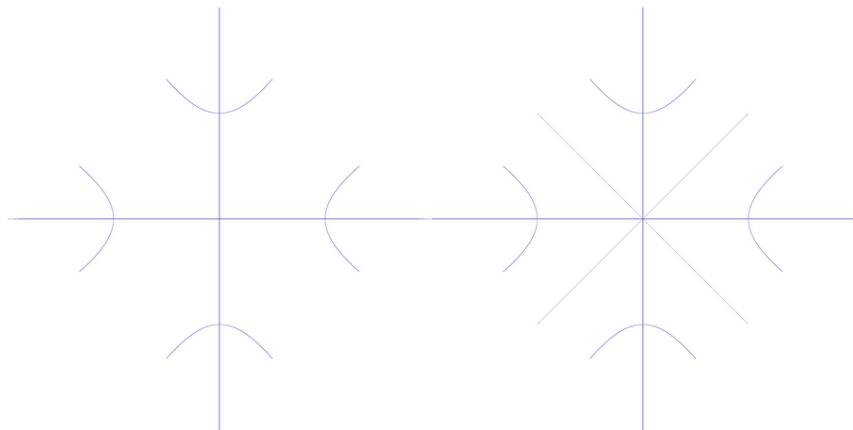
Example [Feinberg/Zee 1999]



Period 2

Periods 1, 2

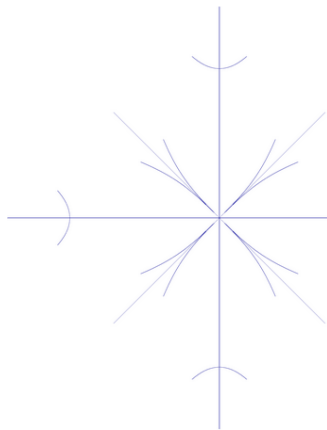
Example [Feinberg/Zee 1999]



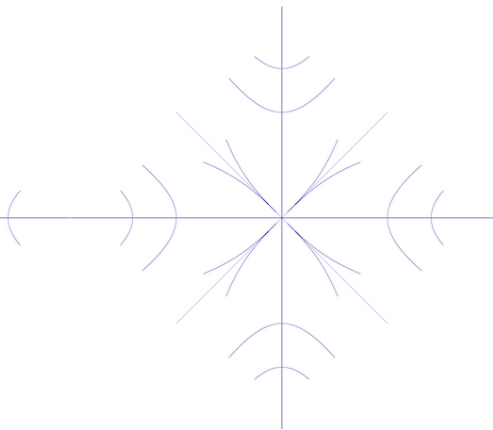
Period 3

Periods 1, ..., 3

Example [Feinberg/Zee 1999]

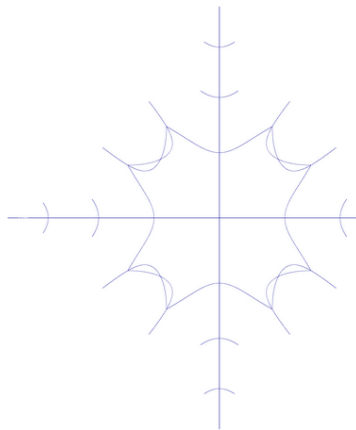


Period 4

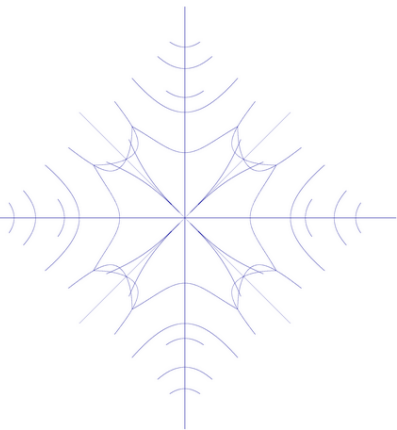


Periods 1, ..., 4

Example [Feinberg/Zee 1999]

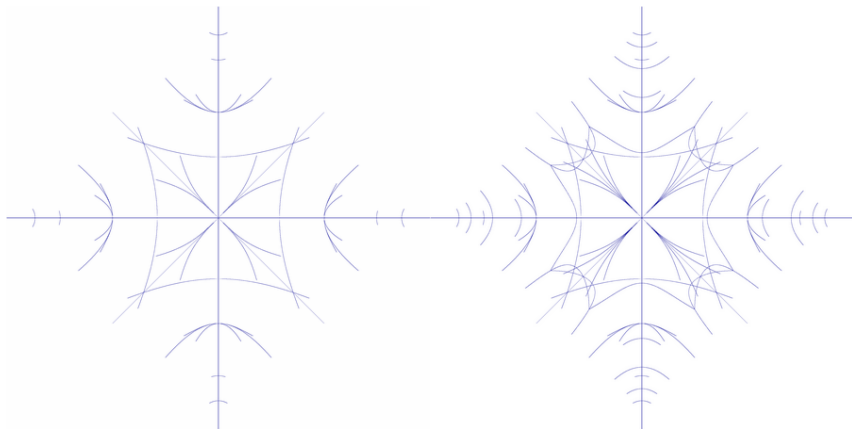


Period 5



Periods 1, ..., 5

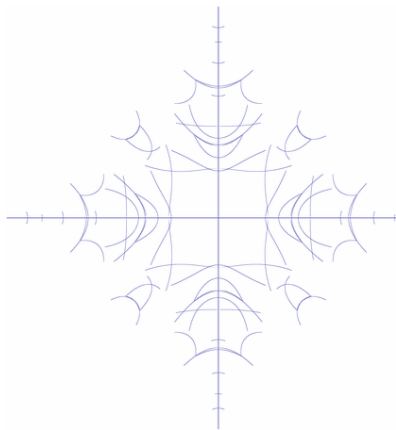
Example [Feinberg/Zee 1999]



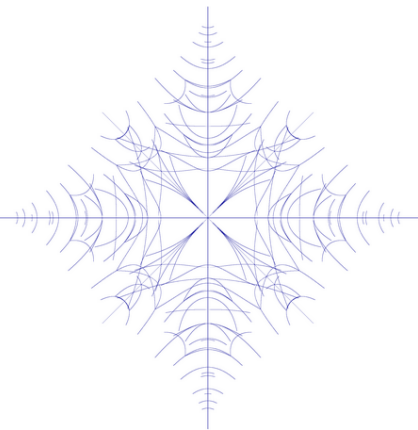
Period 6

Periods 1, ..., 6

Example [Feinberg/Zee 1999]

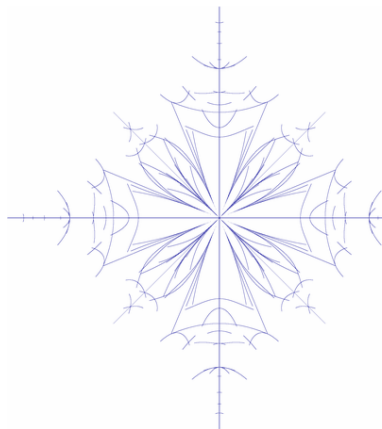


Period 7

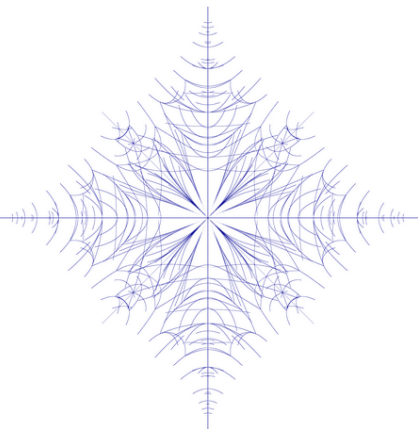


Periods 1, ..., 7

Example [Feinberg/Zee 1999]

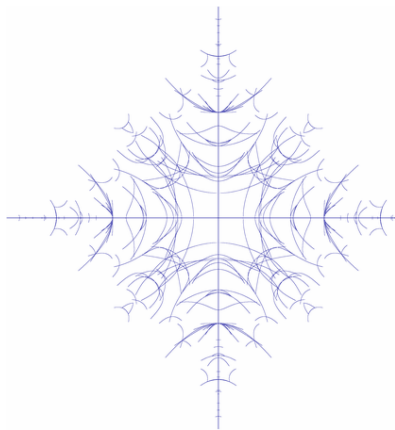


Period 8

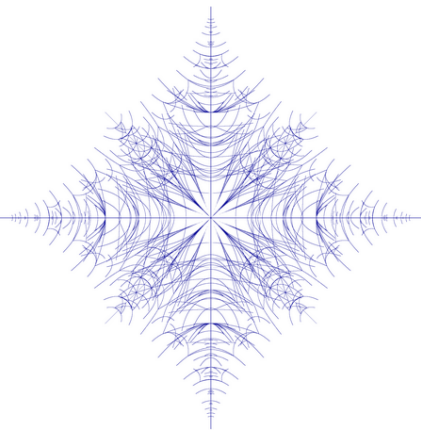


Periods 1, ..., 8

Example [Feinberg/Zee 1999]

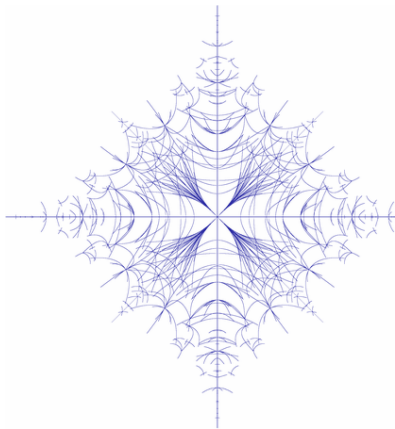


Period 9

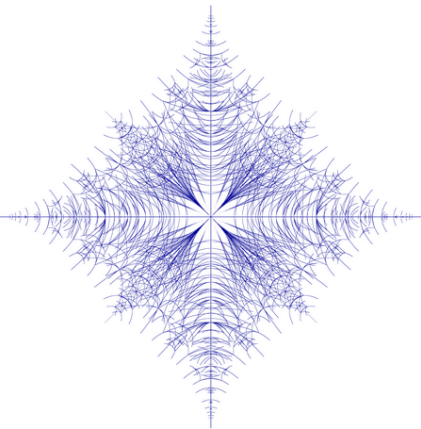


Periods 1, ..., 9

Example [Feinberg/Zee 1999]

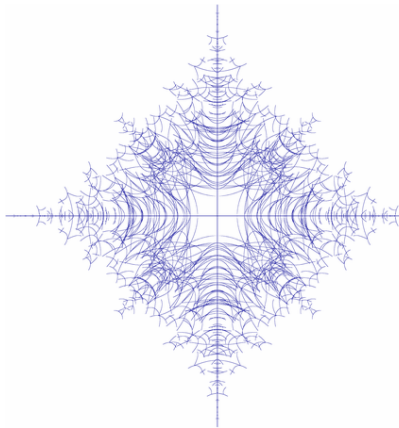


Period 10

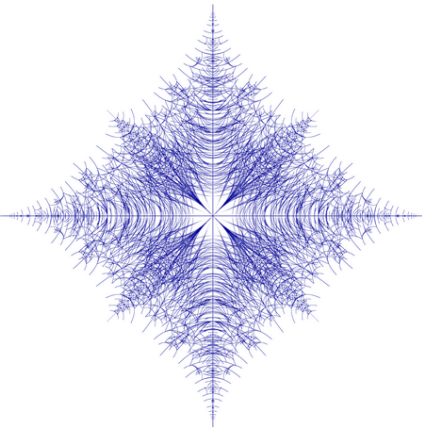


Periods 1, ..., 10

Example [Feinberg/Zee 1999]

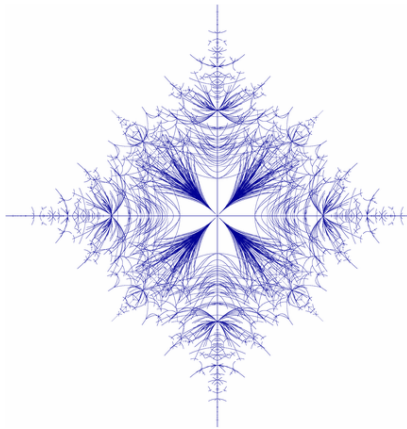


Period 11

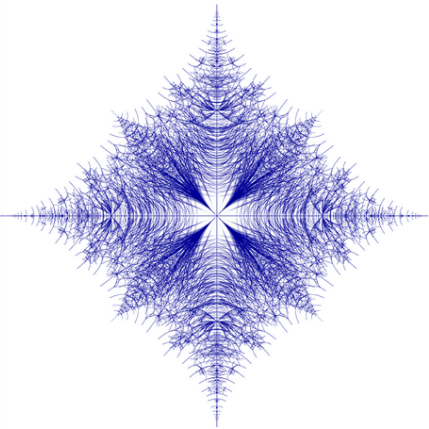


Periods 1, ..., 11

Example [Feinberg/Zee 1999]

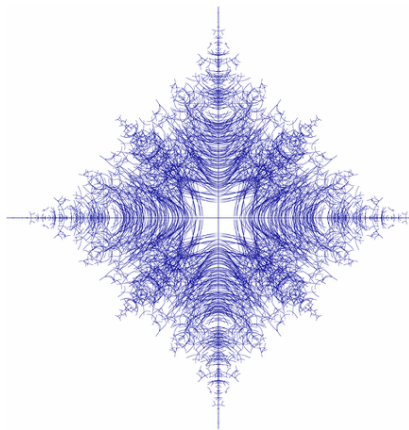


Period 12

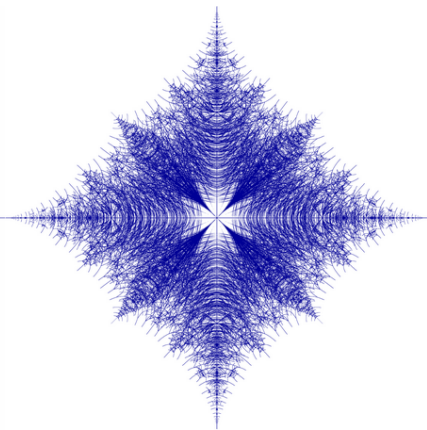


Periods 1, ..., 12

Example [Feinberg/Zee 1999]

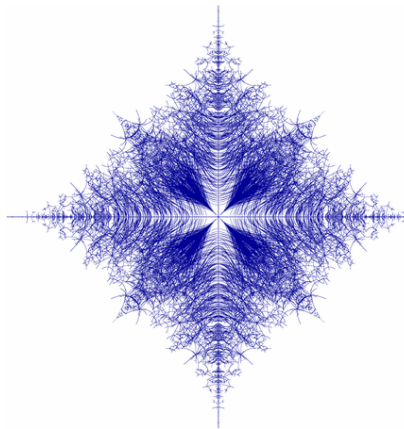


Period 13

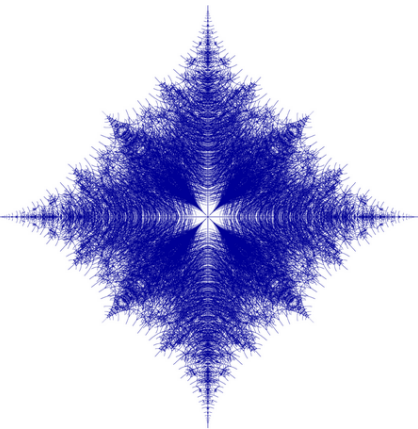


Periods 1, ..., 13

Example [Feinberg/Zee 1999]

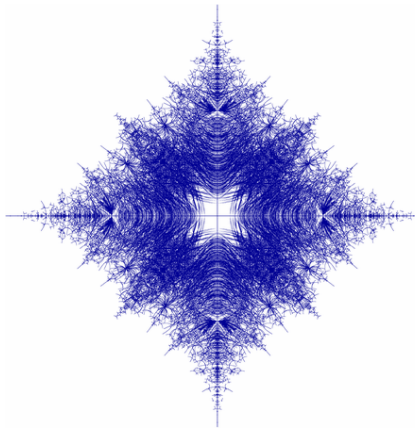


Period 14

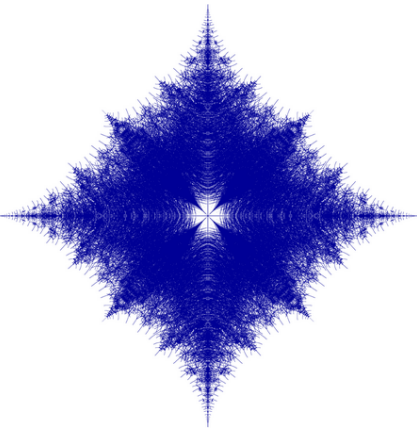


Periods 1, ..., 14

Example [Feinberg/Zee 1999]

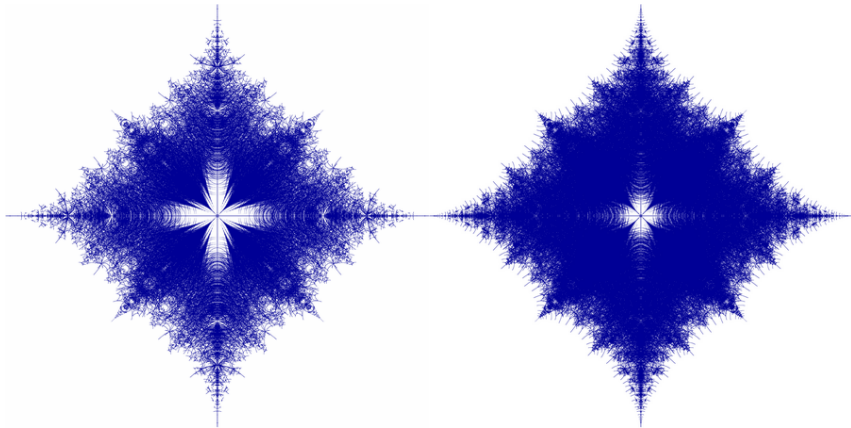


Period 15



Periods 1, ..., 15

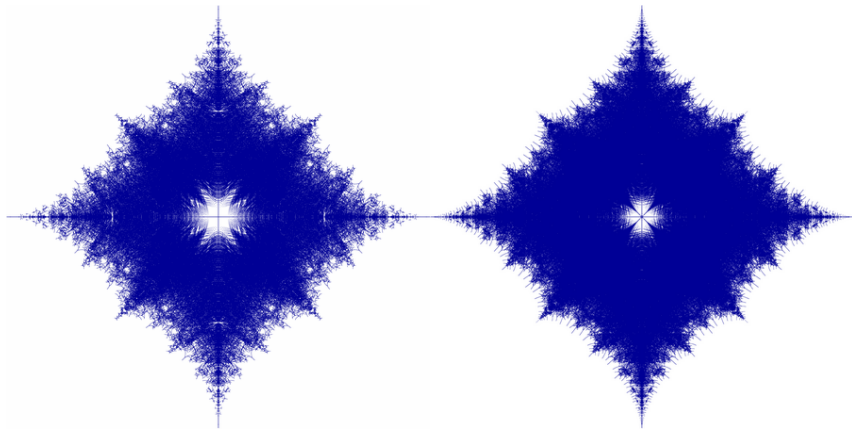
Example [Feinberg/Zee 1999]



Period 16

Periods 1, ..., 16

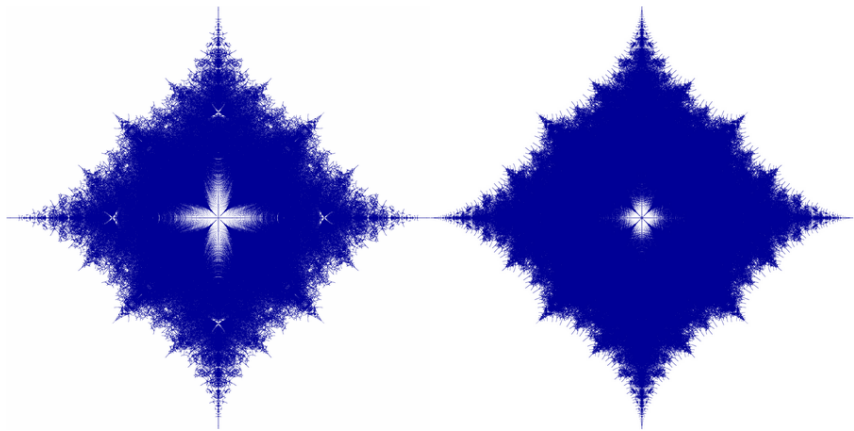
Example [Feinberg/Zee 1999]



Period 17

Periods 1, ..., 17

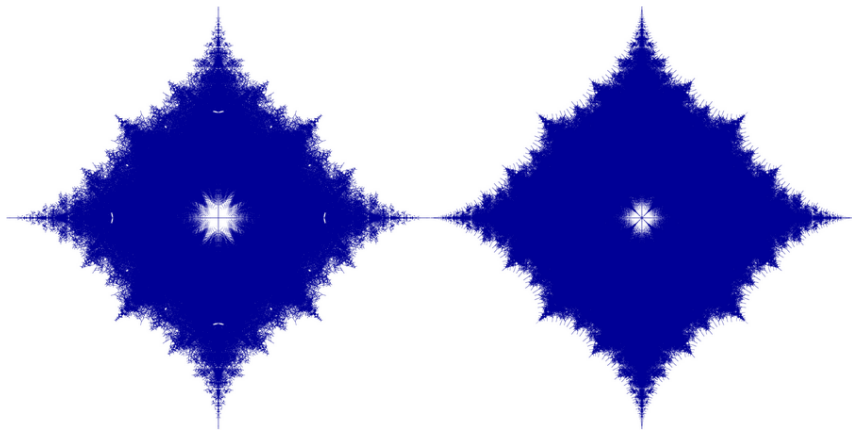
Example [Feinberg/Zee 1999]



Period 18

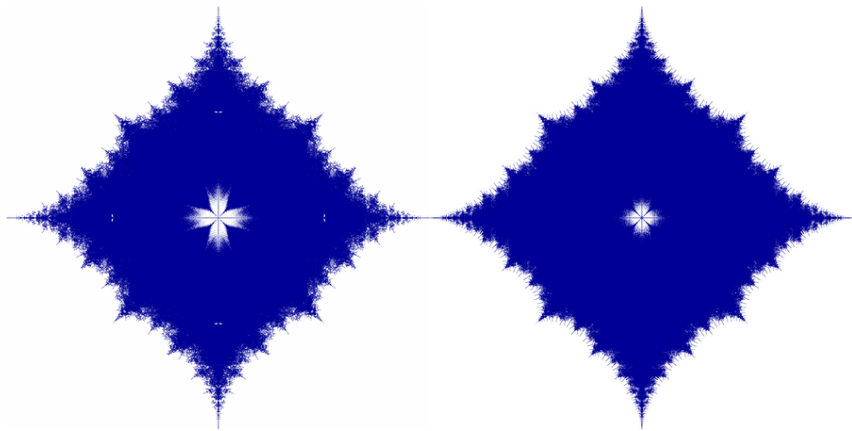
Periods 1, ..., 18

Example [Feinberg/Zee 1999]



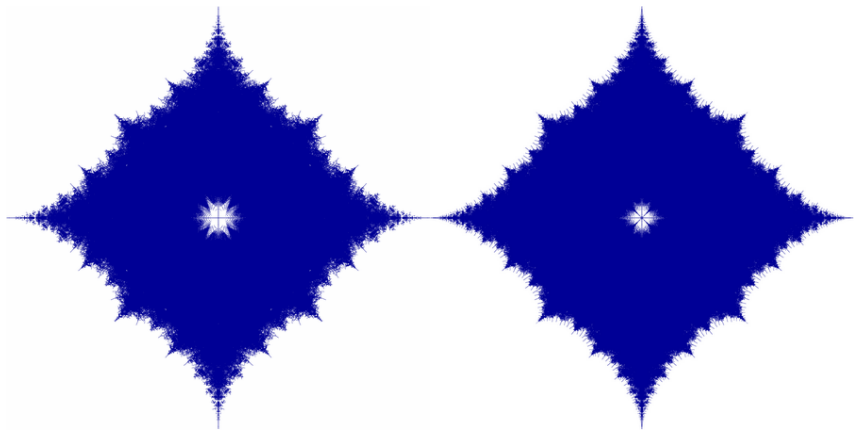
Period 19

Periods 1, ..., 19



Period 20

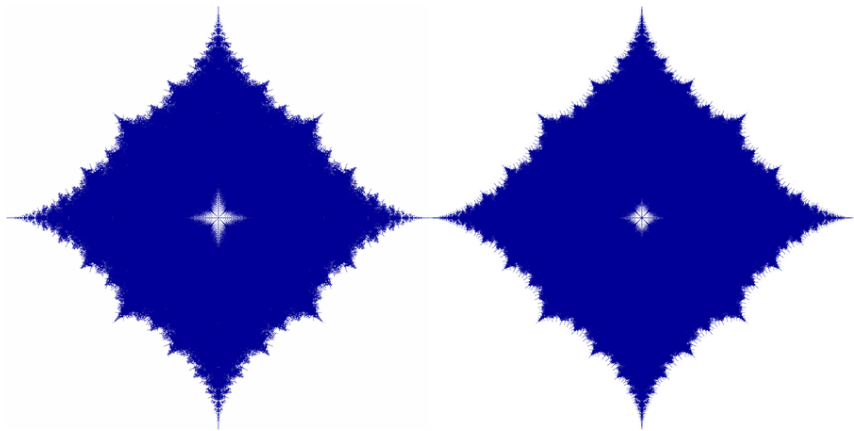
Periods 1, ..., 20



Period 21

Periods 1, ..., 21

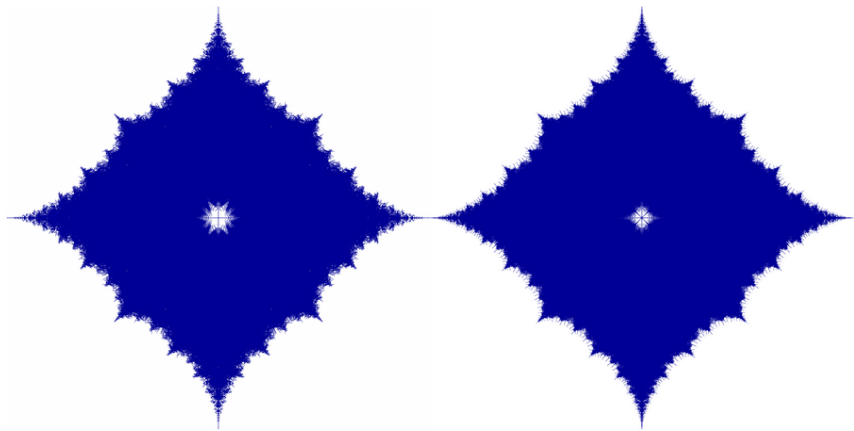
Example [Feinberg/Zee 1999]



Period 22

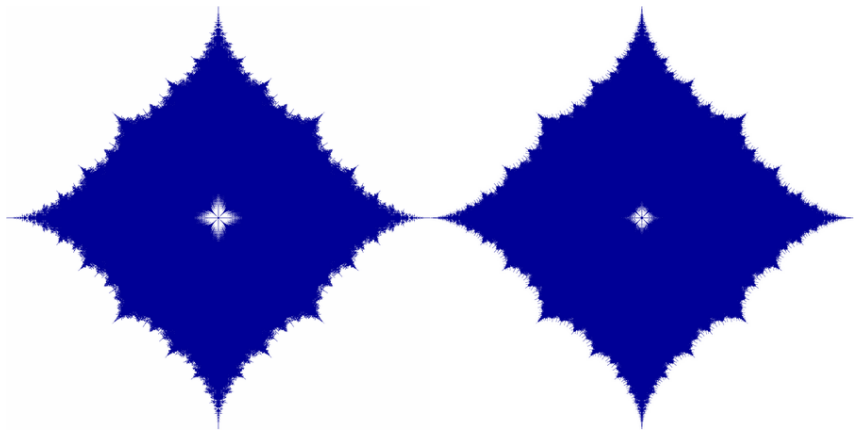
Periods 1, ..., 22

Example [Feinberg/Zee 1999]



Period 23

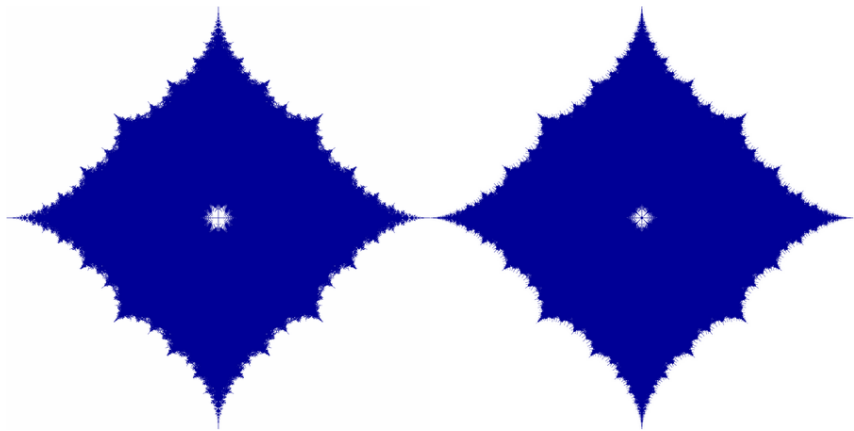
Periods 1, ..., 23



Period 24

Periods 1, ..., 24

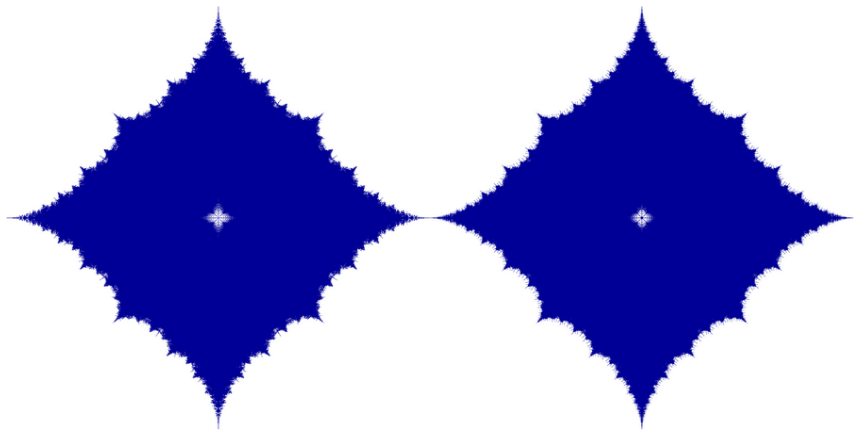
Example [Feinberg/Zee 1999]



Period 25

Periods 1, ..., 25

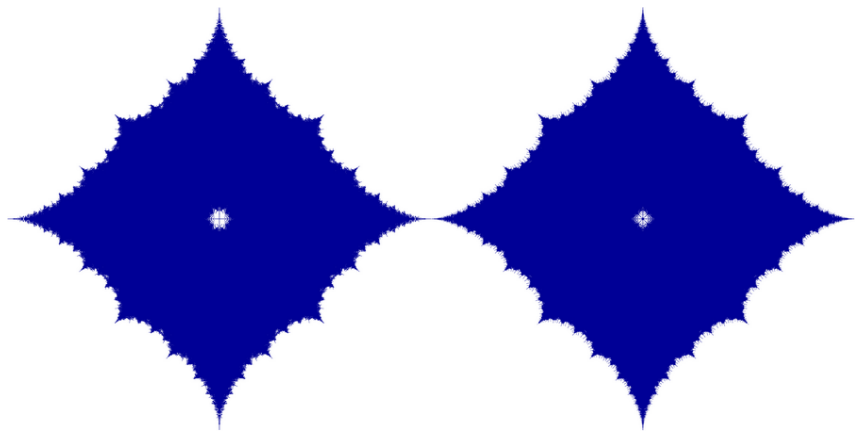
Example [Feinberg/Zee 1999]



Period 26

Periods 1, ..., 26

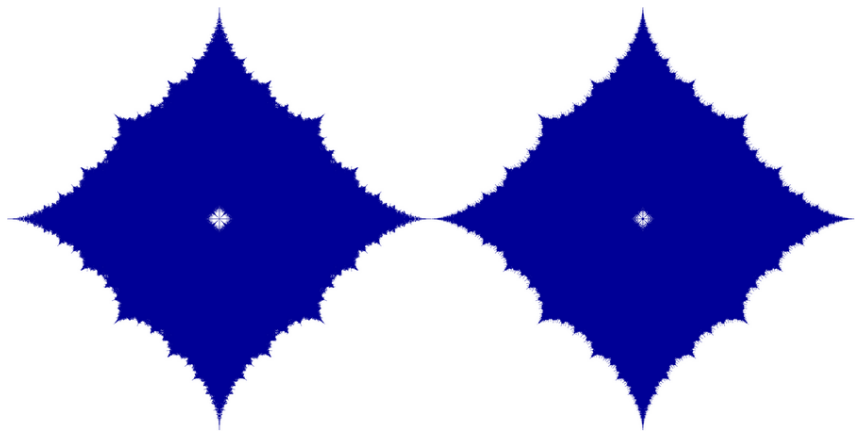
Example [Feinberg/Zee 1999]



Period 27

Periods 1, ..., 27

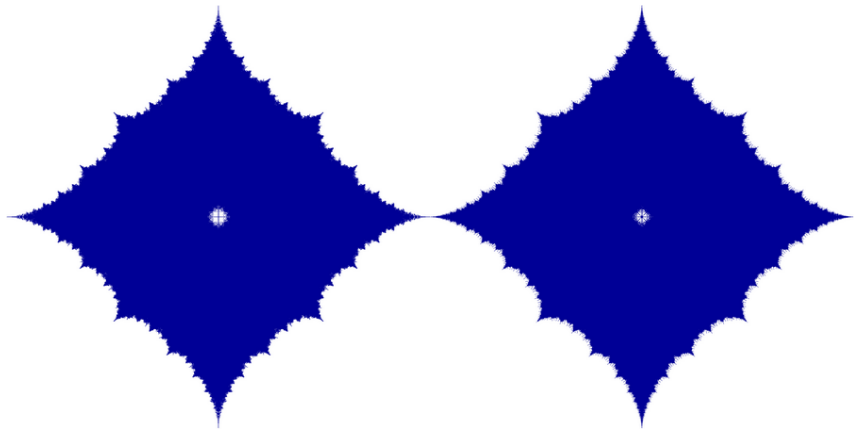
Example [Feinberg/Zee 1999]



Period 28

Periods 1, ..., 28

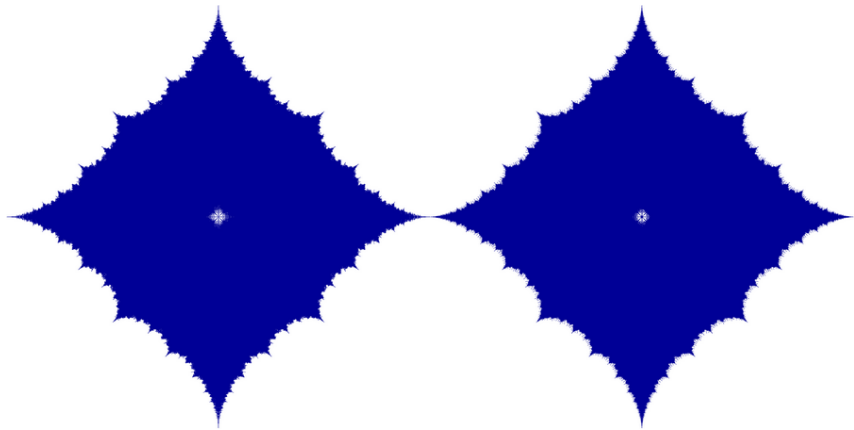
Example [Feinberg/Zee 1999]



Period 29

Periods 1, ..., 29

Example [Feinberg/Zee 1999]



Period 30

Periods 1, ..., 30

Example [Feinberg/Zee 1999]

Recall our “Sandwich 1”: In this example, one has

$$\underbrace{\bigcup_{k \in \mathbb{Z}} \text{spec}_{\varepsilon}(P_{n,k} A^b P_{n,k})}_{=: \sigma_n^{\varepsilon}} \subset \text{spec}_{\varepsilon}(A^b) \subset \underbrace{\bigcup_{k \in \mathbb{Z}} \text{spec}_{\varepsilon + \varepsilon_n}(P_{n,k} A^b P_{n,k})}_{\Sigma_n^{\varepsilon} := \sigma_n^{\varepsilon + \varepsilon_n}}$$

for all $n \in \mathbb{N}$, so let's look at σ_n^{ε} for $\varepsilon = 0$.

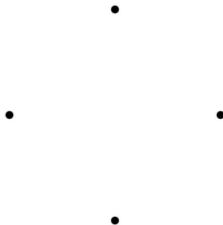
Here are the $n \times n$ matrix eigenvalues

$$\sigma_n^0 = \bigcup_{k \in \mathbb{Z}} \text{spec}(P_{n,k} A^b P_{n,k})$$

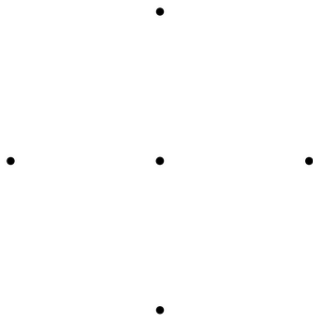
for $n = 1, \dots, 30$:



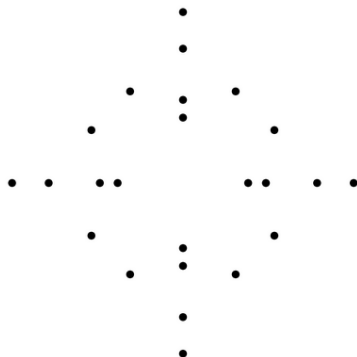
Size 1



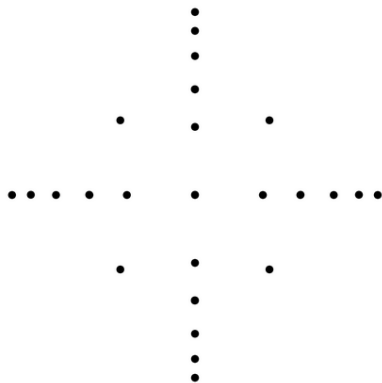
Size 2



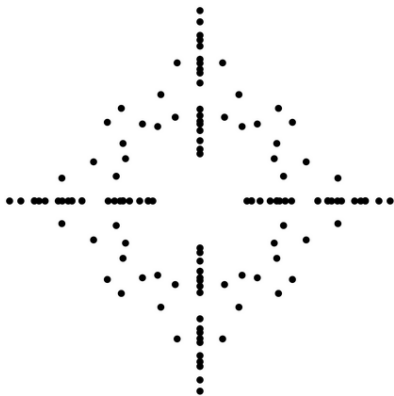
Size 3



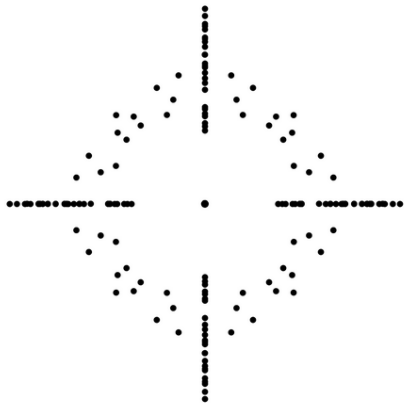
Size 4



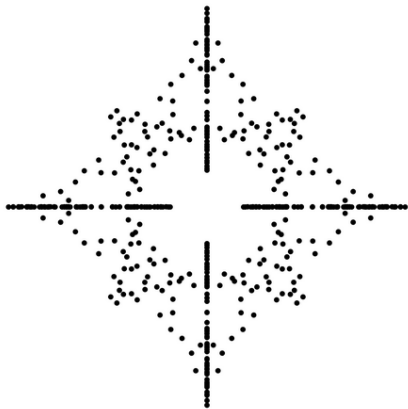
Size 5



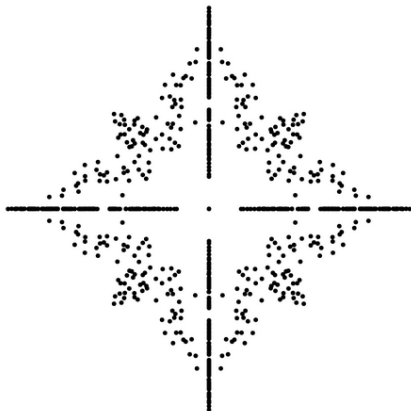
Size 6



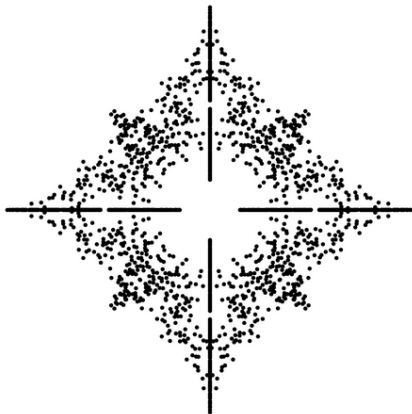
Size 7



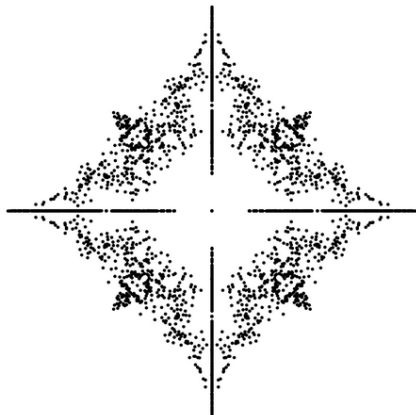
Size 8



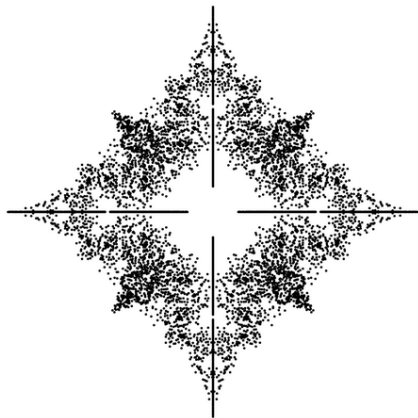
Size 9



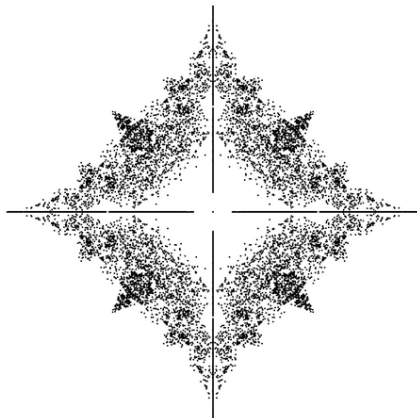
Size 10



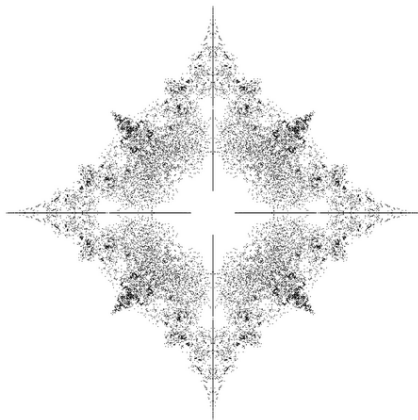
Size 11



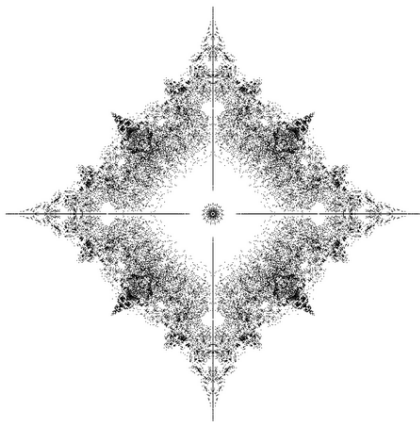
Size 12



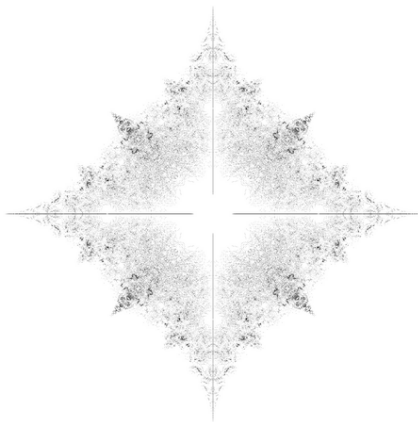
Size 13



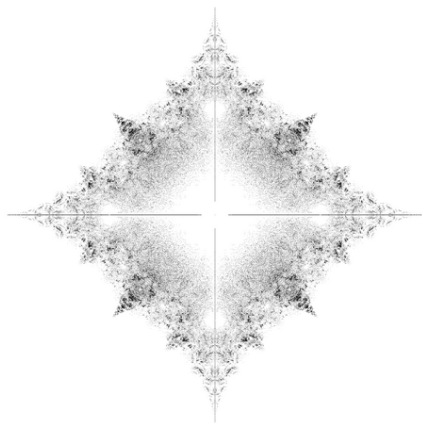
Size 14



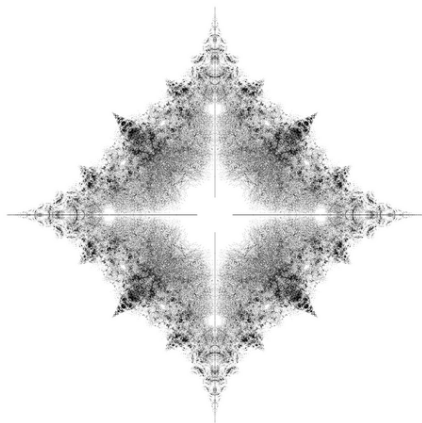
Size 15



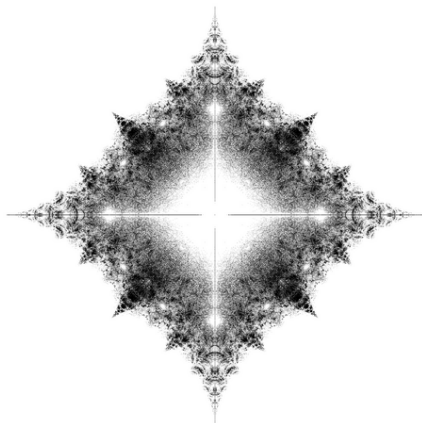
Size 16



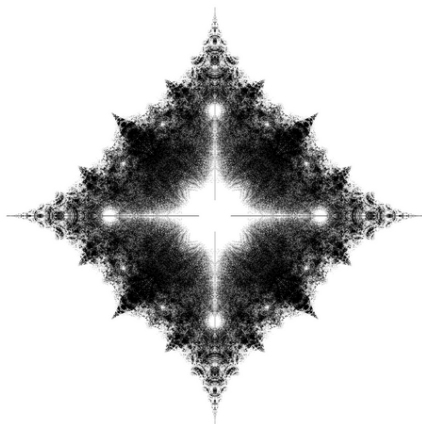
Size 17



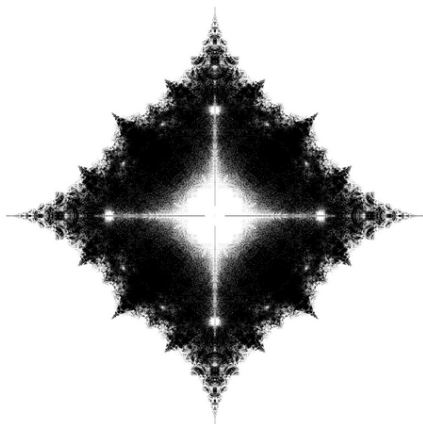
Size 18



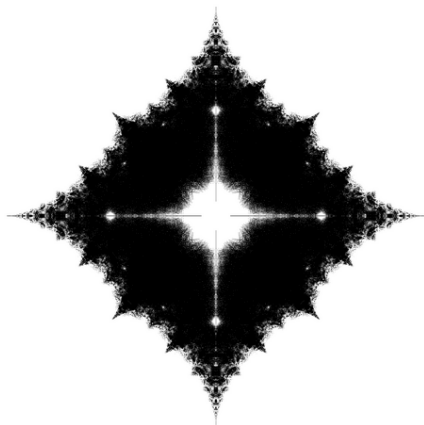
Size 19



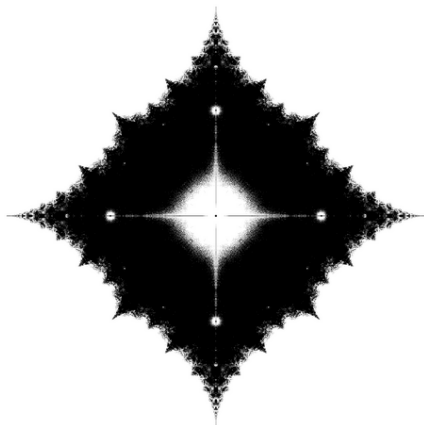
Size 20



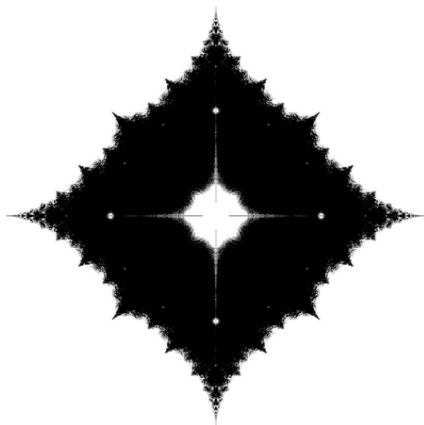
Size 21



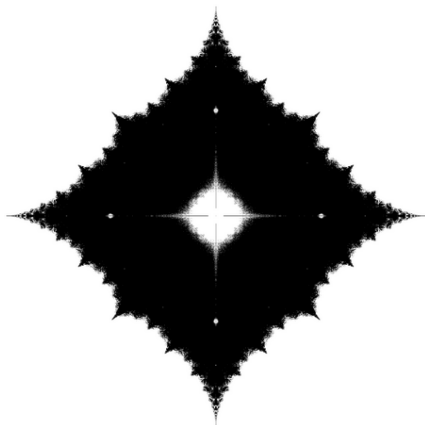
Size 22



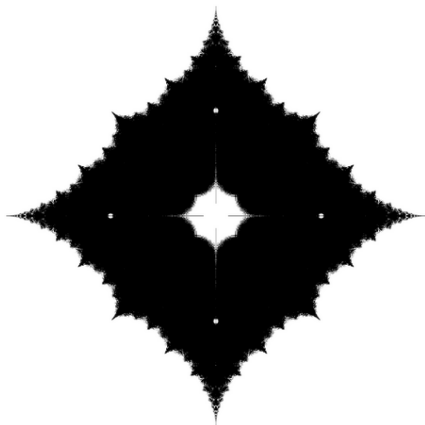
Size 23



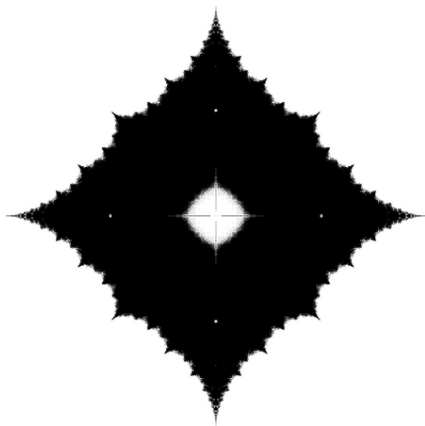
Size 24



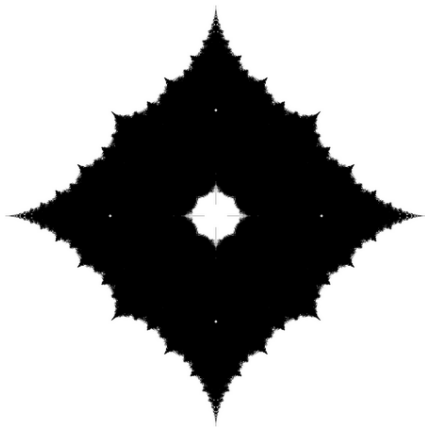
Size 25



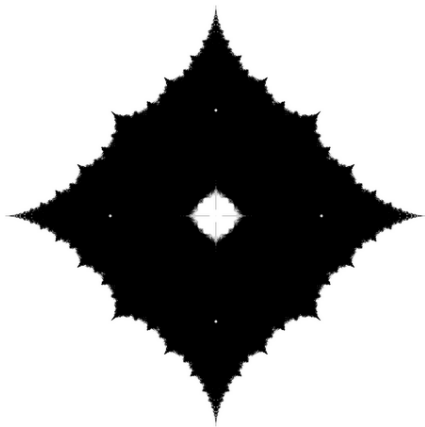
Size 26



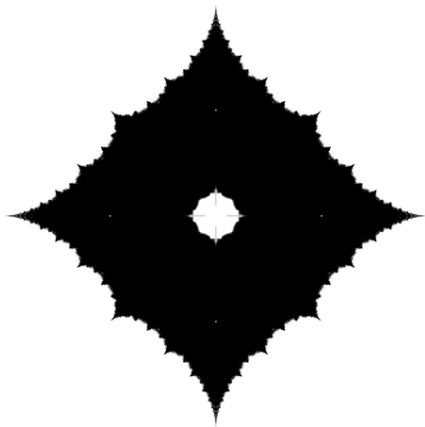
Size 27



Size 28

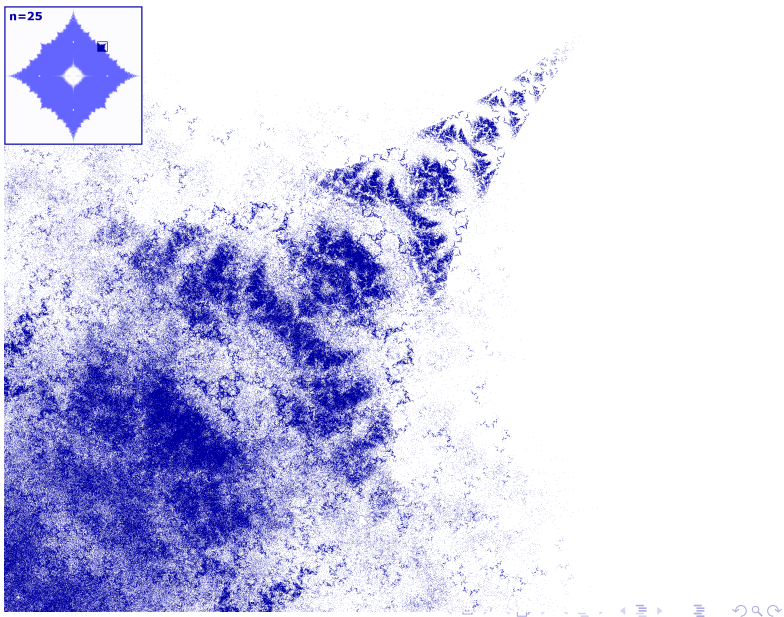


Size 29



Size 30

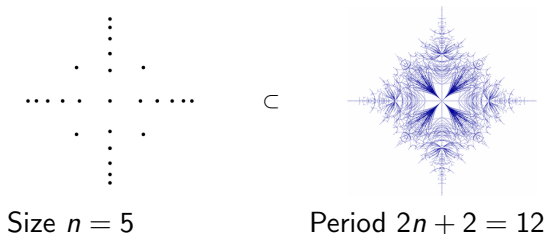
Zoom into Region $1 + i$ of σ_{25}^0



Example [Feinberg/Zee 1999]

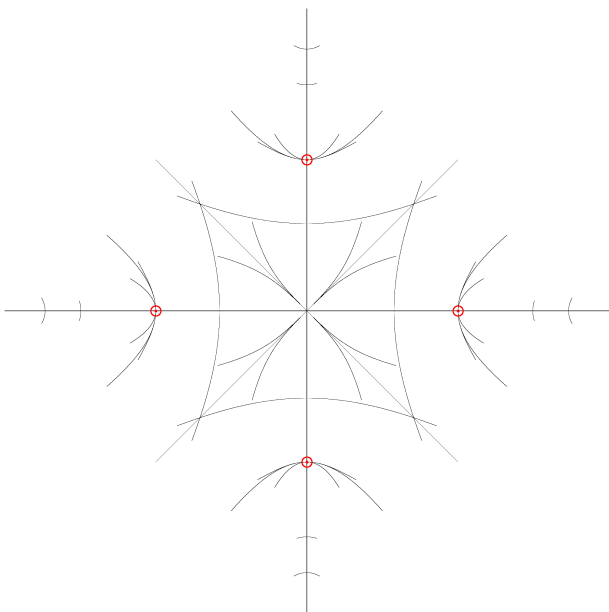
The **finite** matrix spectra σ_n^0 are even **contained** in the **periodic** (infinite) matrix spectra shown before.

More precisely, the spectra of all $n \times n$ principal submatrices are **contained** in the set of all $(2n+2)$ -periodic matrices:

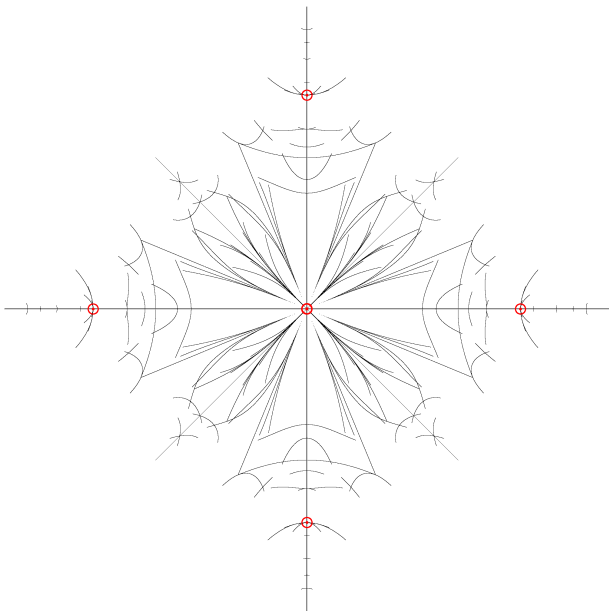


Here we demonstrate this inclusion for some values of n .

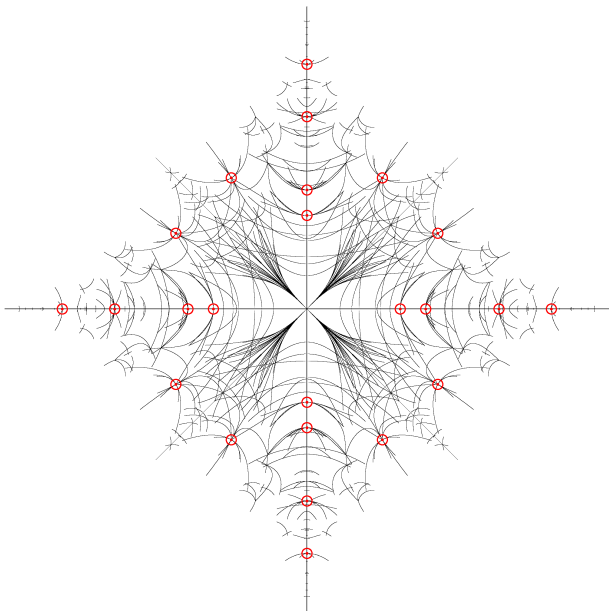
Finite Matrix Spectra in Periodic Matrix Spectra: $n = 2$



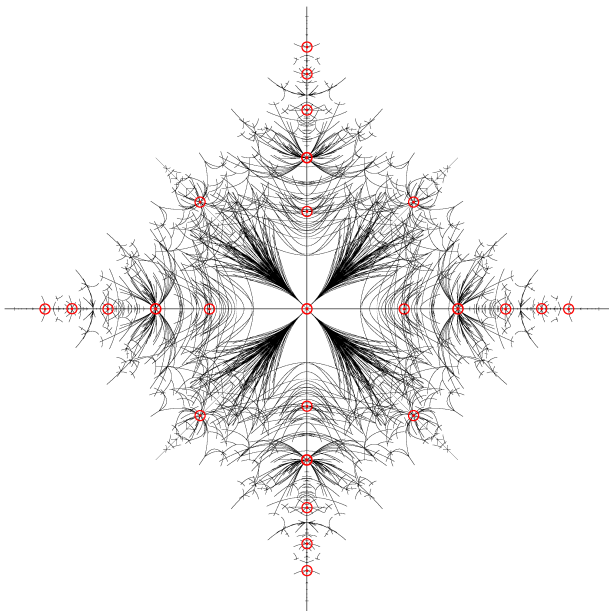
Finite Matrix Spectra in Periodic Matrix Spectra: $n = 3$



Finite Matrix Spectra in Periodic Matrix Spectra: $n = 4$



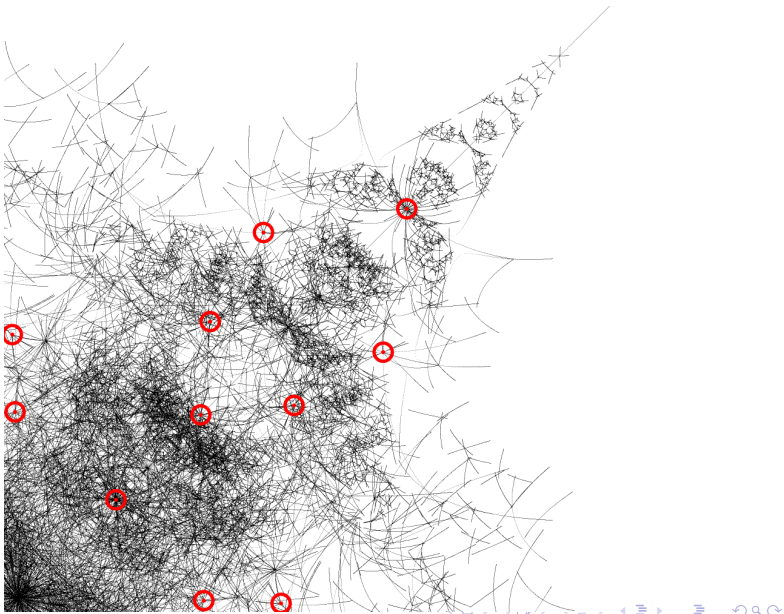
Finite Matrix Spectra in Periodic Matrix Spectra: $n = 5$



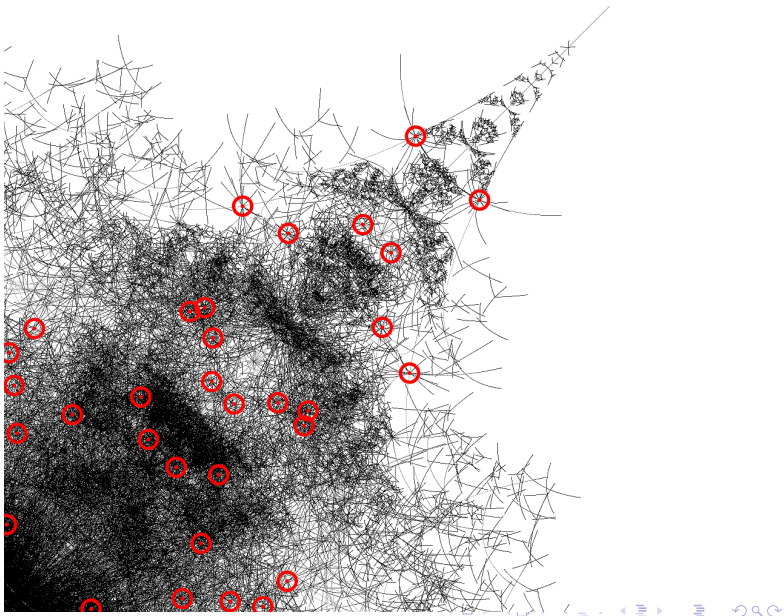
Finite Matrix Spectra in Periodic Matrix Spectra: $n = 8$



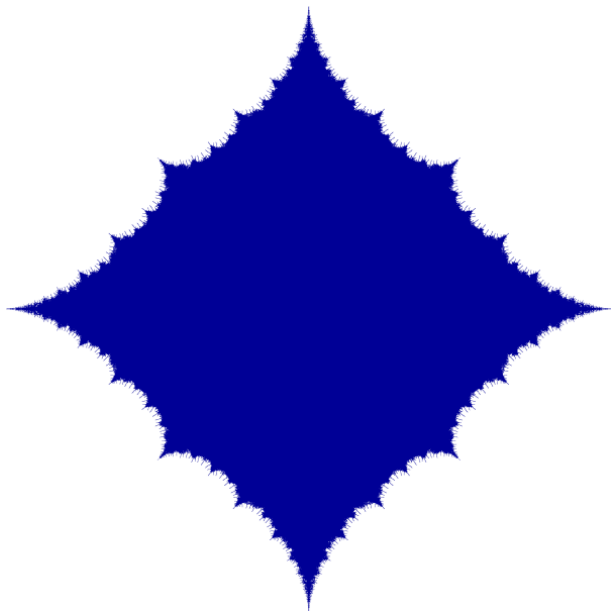
Finite Matrix Spectra in Periodic Matrix Spectra: $n = 9$



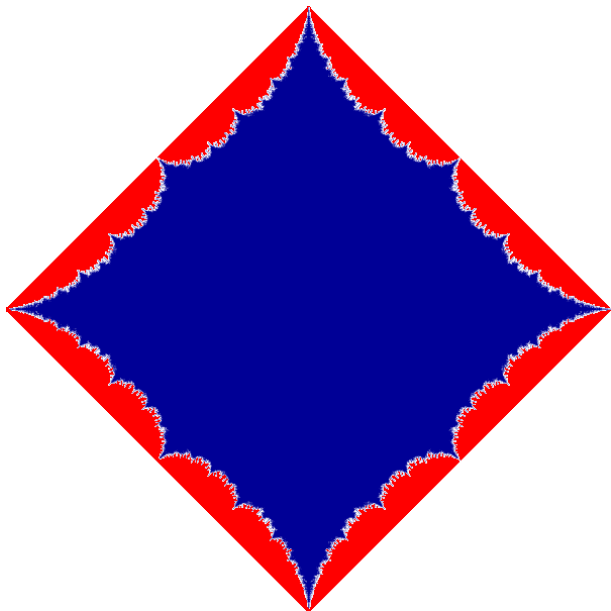
Finite Matrix Spectra in Periodic Matrix Spectra: $n = 10$



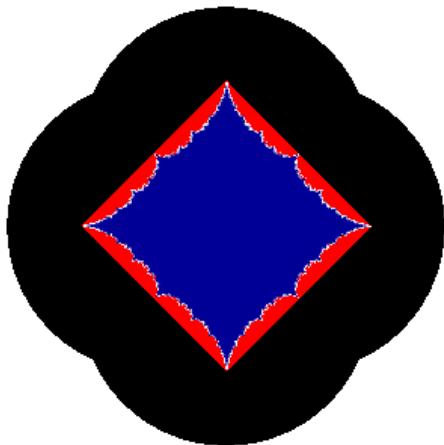
Conjecture: $\text{spec } A^b$ if b is pseudoergodic



Upper bound on $\text{spec } A^b$ by the closed numerical range

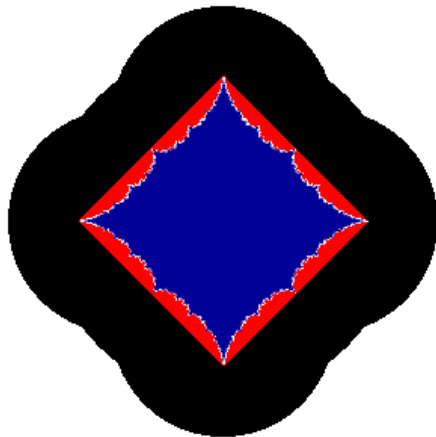


...in comparison with plots of $\Sigma_n^0 = \sigma_n^{0+\varepsilon_n}$



$n = 2$

...in comparison with plots of $\Sigma_n^0 = \sigma_n^{0+\varepsilon_n}$



$n = 3$

...in comparison with plots of $\Sigma_n^0 = \sigma_n^{0+\varepsilon_n}$



$n = 4$

...in comparison with plots of $\Sigma_n^0 = \sigma_n^{0+\varepsilon_n}$



$n = 5$

...in comparison with plots of $\Sigma_n^0 = \sigma_n^{0+\varepsilon_n}$



$n = 6$

...in comparison with plots of $\Sigma_n^0 = \sigma_n^{0+\varepsilon_n}$



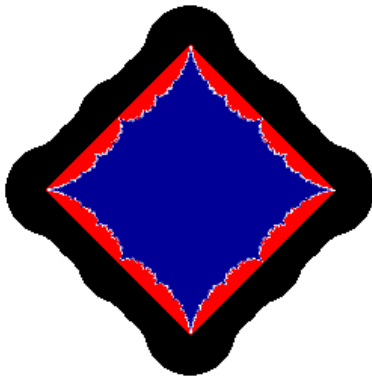
$n = 7$

...in comparison with plots of $\Sigma_n^0 = \sigma_n^{0+\varepsilon_n}$



$n = 8$

...in comparison with plots of $\Sigma_n^0 = \sigma_n^{0+\varepsilon_n}$



$n = 9$

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$n = 11$

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$n = 12$

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$n = 14$

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$n = 15$

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$n = 16$

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$n = 17$

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$n = 18$

Where does \sum_n^0 go as $n \rightarrow \infty$?

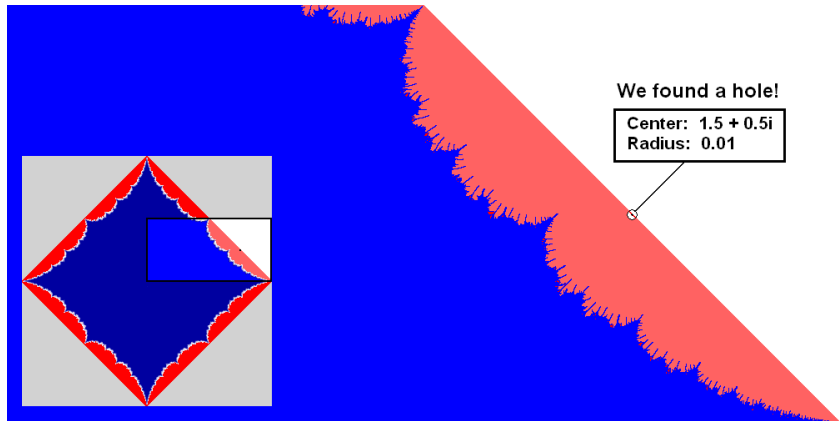
Computational cost for these pics: $n \cdot 2^{n-1} \times$ number of pixels.

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So let us focus on just **one** point (pixel) λ :

Where does Σ_n^0 go as $n \rightarrow \infty$?

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So let us focus on just **one** point (pixel) λ :



$$\lambda = 1.5 + 0.5i \notin \Sigma_{36}^0 \supset \text{spec } A^b$$

The Finite Section Method, Part II

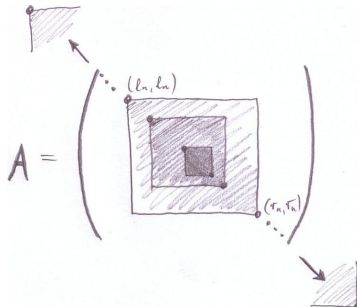
- 1 Classes of Infinite Matrices
- 2 The Finite Section Method, Part I
- 3 Limit Operators
- 4 The Spectrum: Formulas and Bounds
- 5 Spectral Bounds: An Example
- 6 The Finite Section Method, Part II**

Now we come back to the FSM
and bring in our knowledge
on Fredholm indices.

The FSM and the Fredholm index

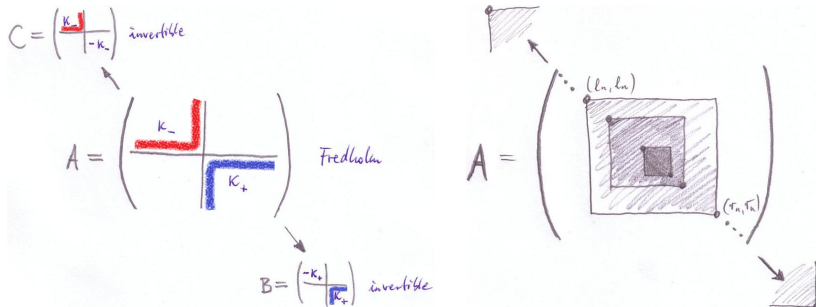
Recall the following two facts in the **bi-infinite** case:

$$C = \begin{pmatrix} K_- & \\ & -K_- \end{pmatrix} \text{ invertible}$$
$$A = \begin{pmatrix} K_- & \\ & K_+ \end{pmatrix} \text{ Fredholm}$$
$$B = \begin{pmatrix} -K_+ & \\ & K_+ \end{pmatrix} \text{ invertible}$$



The FSM and the Fredholm index

Recall the following two facts in the **bi-infinite case**:



So $\text{ind } A_+ = 0 = \text{ind } A_-$ is **necessary** for applicability of the FSM!

The FSM and the Fredholm index

Abbreviate $\text{ind } A_+ =: \kappa_+$ and $\text{ind } A_- =: \kappa_-$.

$\kappa_+ = 0 = \kappa_-$ is **necessary** for applicability of the FSM!

Otherwise: **Shift** the system up or down accordingly, i.e. place the corners of your finite sections A_n on another (the κ_-^{th}) diagonal!

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Otherwise: **Shift** the system up or down accordingly, i.e. place the corners of your finite sections A_n on another (the κ_-^{th}) diagonal!

This means: Replace $Ax = b$ by $V_{\kappa_+}Ax = V_{\kappa_+}b$.

Reason: The new system has plus-index zero:

$$\text{ind } (V_{\kappa_+}A)_+ = \text{ind } (V_{\kappa_+})_+ + \text{ind } A_+ = -\kappa_+ + \kappa_+ = 0$$

This process is called **index cancellation**.

⇒ We have found the (from the FSM perspective)
“proper” main diagonal of A !

Example: FSM for slowly oscillating operators

Suppose $A \in BDO(E)$ has slowly oscillating diagonals.
We want to solve $Ax = b$ by the FSM.

Assumption (minimal): A be invertible.

Step 1: Compute the plus-index $\kappa_+ := \text{ind } A_+$.

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Therefore, take an arbitrary limop B of A at $+\infty$ and recall that $\text{ind } B_+ = \text{ind } A_+$.

A is slowly oscillating $\Rightarrow B_+$ is Toeplitz $\Rightarrow \kappa_+ = -\text{wind}(a(\mathbb{T}), 0)$

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Step 3: Truncate.

Remarkable fact: We can truncate at **arbitrary** points l_n and r_n !
Reason: All limops B and C (w.r.t. subsequences of r and l , resp) are Laurent operators. So all B_- and all C_+ are Toeplitz operators that are Fredholm of index zero (after index cancellation).

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Coburn's lemma \Rightarrow all B_- and all C_+ are invertible, as well as A .

FSM theorem \Rightarrow The FSM applies.

Finite Sections of Random Jacobi Operators

Assumption (minimal): $A \in \Psi E(U, V, W)$ is invertible.
How do we truncate A to get an applicable FSM?

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Let's take a Laurent operator B . So pick **arbitrary** $u \in U$, $v \in V$ and $w \in W$. Then $\kappa_+ = \text{ind } A_+ = \text{ind } B_+ = -\text{wind}(E_{u,w}, v)$.

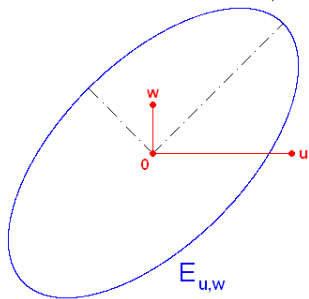
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It's very simple:

- if v is outside $E_{u,w}$: $\kappa_+ = 0$
- if v is inside $E_{u,w}$: $\kappa_+ = \pm 1$
 - if $|u| > |w|$: $\kappa_+ = -1$
 - if $|u| < |w|$: $\kappa_+ = +1$

Note: The result κ_+ is independent of $u \in U$, $v \in V$, $w \in W$!

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- $\kappa_+ = -1$: shift **up**

$$\begin{pmatrix} \diagdown \\ \diagdown \\ \diagdown \end{pmatrix} \begin{pmatrix} x \\ \\ \end{pmatrix} = \begin{pmatrix} b \\ \uparrow \\ \end{pmatrix}$$

- $\kappa_+ = 0$: leave as it is

$$\begin{pmatrix} \diagdown \\ \diagdown \\ \diagdown \end{pmatrix} \begin{pmatrix} x \\ \\ \end{pmatrix} = \begin{pmatrix} \\ b \\ \end{pmatrix}$$

- $\kappa_+ = +1$: shift **down**

$$\begin{pmatrix} \diagdown \\ \diagdown \\ \diagdown \end{pmatrix} \begin{pmatrix} x \\ \\ \end{pmatrix} = \begin{pmatrix} \downarrow \\ b \\ \end{pmatrix}$$

In either case, the new system $\tilde{A}x = \tilde{b}$ has $\text{ind } \tilde{A}_+ = 0 = \text{ind } \tilde{A}_-$.

Finite Sections of Random Jacobi Operators

Step 3. Do the truncations.

Choose the truncation points $\dots < l_2 < l_1 < r_1 < r_2 < \dots$ so that

$$\begin{aligned} & \begin{pmatrix} v_{l_n} & w_{l_n+1} & & \\ u_{l_n} & \ddots & \ddots & \\ & \ddots & \ddots & \\ & & & \end{pmatrix} \longrightarrow \begin{pmatrix} v & w & & \\ u & \ddots & \ddots & \\ & \ddots & \ddots & \\ & & & \end{pmatrix} =: C_+ \\ \text{and} \quad & \begin{pmatrix} \ddots & \ddots & & \\ \ddots & \ddots & & w_{r_n} \\ & u_{r_n-1} & v_{r_n} & \\ & & & \end{pmatrix} \longrightarrow \begin{pmatrix} \ddots & \ddots & & \\ \ddots & \ddots & & w \\ & u & v & \\ & & & \end{pmatrix} =: B_- \end{aligned}$$

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as $n \rightarrow \infty$, for some fixed $u \in U$, $v \in V$ and $w \in W$.

Both **Toeplitz** operators C_+ and B_- are Fredholm of index 0 (because $\text{ind } A_+ = 0 = \text{ind } A_-$) so they are **invertible** (Coburn).

But how about the 'full' FSM?

The previous truncation pattern was specially adapted to the operator $A \in \Psi E(U, V, W)$ at hand.

The **full (or usual) FSM** uses the cut-off sequences $l = (-1, -2, \dots)$ and $r = (1, 2, \dots)$.

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- 2 If $0 \in U, W$ and A is invertible then the full FSM is applicable.
- 3 If $\kappa_+ := \text{ind } A_+ = \pm 1$ and A is invertible then, after index cancellation, the full FSM is applicable.

Thank you!



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M. LINDNER:
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