## Orthogonal functions and matrix computations

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## Example 1: three-term recurrence relation

Consider inner product on the real line

$$
\langle p, q\rangle=\int_{\mathbb{R}} p(x) q(x) d \mu(x) .
$$

Orthogonal polynomials $p_{i}$ satisfy a three-term recurrence relation

$$
x p_{n}(x)=a_{n} p_{n+1}(x)+b_{n} p_{n}(x)+a_{n-1} p_{n-1}(x)
$$

In matrix language

$$
x\left[p_{0}(x), p_{1}(x), p_{2}(x), \ldots\right]=\left[p_{0}(x), p_{1}(x), p_{2}(x), \ldots\right] J
$$

with $J$ the Jacobi matrix

$$
J=\left[\begin{array}{llll}
b_{0} & a_{0} & & \\
a_{0} & b_{1} & a_{1} & \\
& a_{1} & b_{2} & \ddots \\
& & \ddots & \ddots
\end{array}\right]
$$

Many results involving orthogonal functions can be translated into matrix language and vice versa.

## Two examples:

- recurrence relations
- spectral transformations

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## Example 2: spectral transformation

Two related inner products $\langle\cdot\rangle_{\mu}$ and $\langle\cdot\rangle_{\nu}$

$$
\begin{aligned}
& \langle p, q\rangle_{\mu}=\int_{\mathbb{R}} p(x) q(x) d \mu(x) \\
& \langle p, q\rangle_{\nu}=\int_{\mathbb{R}} p(x) q(x) \mathbf{d} \nu(\mathbf{x})=\int_{\mathbb{R}} p(x) q(x)(\mathbf{x}-\boldsymbol{\beta}) \mathbf{d} \mu(\mathbf{x})
\end{aligned}
$$

Relation between Jacobi matrices $J_{\mu}$ and $J_{\nu}$ associated with $\mu$ and $\nu$, respectively:

$$
\begin{aligned}
& J_{\mu}-\beta I=L L^{T} \\
& J_{\nu}-\beta I=L^{T} L
\end{aligned}
$$

- Agrees with one step of a semi-infinite Cholesky LR algorithm
- $L$ is lower bidiagonal
[Bueno, Marcellán, Dopico]
[Galant], [Kautsky, Golub], [Watkins],


## Discrete inner product

Given the basis functions $p_{1}(z), p_{2}(z), \ldots$,
$n$ points $z_{i} \in \mathbb{C}$ and corresponding weights $w_{i}>0$,
define within the vector space $\mathcal{P}^{n}=\left\{\sum_{i=1}^{n} c_{i} p_{i}(z)\right\}$ the discrete inner product:

$$
\begin{equation*}
\langle p, q\rangle=\sum_{i=1}^{n} w_{i}^{2} \overline{p\left(z_{i}\right)} q\left(z_{i}\right) \tag{1}
\end{equation*}
$$

Let $\left\{a_{j}(z)\right\}$ be the orthonormal functions, i.e., $\left\langle a_{i}, a_{j}\right\rangle=\delta_{i j}$, such that $a_{j}(z) \in \mathcal{P}^{j} \backslash \mathcal{P}^{j-1}, j=1,2, \ldots, n$.
Discrete LS approximation with orthonormal functions
Given: a function $f(z)$
Find: the function $p(z)=\sum_{i=1}^{\alpha} c_{i} p_{i}(z)$ with $\alpha \leq n$ s.t.

$$
\sum_{i=1}^{n} w_{i}^{2}\left|f\left(z_{i}\right)-p\left(z_{i}\right)\right|^{2} \quad \text { is minimal. }
$$

Solution: The solution $p(z)$ can be represented as $p(z)=\sum_{j=1}^{\alpha} c_{j} a_{j}(z)$ with $c_{j}=\left\langle a_{j}, f\right\rangle$.
Thus LS problem is reduced to the problem of computing $a_{j}(z)$.

## General Scheme

Given: the nodes $z_{i}$ and the corresponding weights $w_{i}$, $i=1,2, \ldots, n$ of the discrete inner product.
Based on the nodes, one or more $n \times n$ diagonal matrices are derived: $D_{1}, D_{2}$
Definition (general IEP)
Given: $D_{1}, D_{2}, \ldots$ - diagonal matrices, $\mathbf{w}=\left(w_{i}\right)$ - weights
Find: Unitary matrix $Q$ and matrices $H_{j}$ having a specific structure such that

$$
Q^{H} \mathbf{w}=\|\mathbf{w}\| \mathbf{e}_{1}, \quad Q^{H} D_{j} Q=H_{j} .
$$

## Computing the recurrence parameters

For many important choices of the basis functions $p_{i}(z)$ of the vector space $\mathcal{P}^{n}=\left\{\sum_{i=1}^{n} c_{i} p_{i}(z)\right\}$, computing the recurrence parameters for the corresponding orthonormal functions $a_{i}(z)$

$$
\begin{aligned}
\left\langle a_{j}(z), a_{k}(z)\right\rangle & =\sum_{i=1}^{n} w_{i}^{2} \overline{a_{j}\left(z_{i}\right)} a_{k}\left(z_{i}\right) \\
& =Q_{j}^{H} Q_{k} \\
& =\delta_{j k}
\end{aligned}
$$

reduces to an inverse eigenvalue problem (IEP). Choices for the basis functions $p_{i}(z)$ :

- $1, z, z^{2}, z^{3}, \ldots$ (orthogonal polynomials, OP)
- $1, z, z^{-1}, z^{2}, z^{-2}, z^{3}, z^{-3}, \ldots$ (orthogonal Laurent polynomials)
- any sequence such that $z^{k}$ with $k>0$ comes after $z^{k-1}$ and $z^{-k}$ with $k>0$ comes after $z^{-(k-1)}$ (general orthogonal Laurent polynomials)
- $1, \frac{1}{z-y_{1}}, \frac{1}{z-y_{2}}, \ldots$ (orthogonal proper rational functions)
- ...

This can be extended to:

- multivariate orthogonal functions
- vector and matrix cases
- ...

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## Comments

- The structure of the matrix(ces) $H_{j}$ is determined by the recurrence relation for the orthonormal functions $a_{i}(z)$.
- The columns of the unitary matrix $Q_{i}$ are connected to the orthonormal functions as follows:

$$
Q_{k}=\operatorname{diag}\left(w_{i}\right)\left[a_{k}\left(z_{i}\right)\right]_{i=1,2, \ldots, n}
$$

- The columns $Q_{k}$ satisfy a corresponding recurrence relation.
- The orthonormality of the functions $a_{k}(z)$ is equivalent to $Q$ being unitary:


## Computing the recurrence parame recurrence paran by solving IEP

## Illustrations

In this talk:

- orthonormal polynomials in one variable [Reichel, Ammar, Gragg][Elhay, Golub, Kautsky][Golub, Meurant]
- orthonormal polynomials in two variables [VB, Chesnokov]
- orthonormal polyanalytic polynomials
- orthonormal Laurent polynomials
- orthonormal rational functions [VB, Fasino, Gemignani, Mastronardi]


## Algorithm for solving IEP - one-varaiable OP


$\rightarrow$ sequence of unitary similarity transformations using
Givens rotations $\rightarrow$


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One-variable orthogonal polynomials
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Basis functions: $p_{j}(z)=z^{j}, j=0,1, \ldots$
Recurrence relation: $z\left[a_{0}(z), a_{1}(z), \ldots\right]=\left[a_{0}(z), a_{1}(z), \ldots\right] H$ with $H$ upper Hessenberg
Hence: Computing (the recurrence relation coefficients of) the OPs $a_{j}(z)$ can be done by solving the following IEP.
Definition (IEP - one-variable OP)
Given: $D_{z}=\operatorname{diag}\left(z_{i}\right)$ - points, $\mathbf{w}=\left(w_{i}\right)$ - weights
Find: Unitary $Q$ and upper Hessenberg $H$ such that

$$
Q^{H} \mathbf{w}=\|\mathbf{w}\| \mathbf{e}_{1}, \quad Q^{H} D_{z} Q=H
$$

[Boley, Golub]

## 2 points



3 points (1)

$\xrightarrow{G_{w}(1,3)}$


6 points (1)

$\xrightarrow{G_{w}(1,6)}$


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$\xrightarrow{G(2,3)}$

. . skip some steps and jump to 6 points .

6 points (2)

$\xrightarrow{G(2,6)}$
3 points (2)

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One-variable orthogonal polynomials

6 points (3)


6 points (5)

$\xrightarrow{G(5,6)}$

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$\xrightarrow{G(4,6)}$

- in general: $\mathcal{O}\left(n^{3}\right)$ FLOPS
- $w_{i}$ real, $z_{i}$ real: $\mathcal{O}\left(n^{2}\right)$ FLOPS
- $z_{i}$ on the complex unit circle: $\mathcal{O}\left(n^{2}\right)$ FLOPS using Schur parametrization, $H$ is a unitary Hessenberg matrix
6 points (4)

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```

One-variable

$$
\begin{aligned}
& \text { One-variable } \\
& \text { orthogonal polynomials }
\end{aligned}
$$

Two-variable orthogonal polynomials
Consider the monomials in two variables as basis functions $p_{k}(x, y)=x^{i} y^{j}$.
Choose an ordering of these basis functions such that:

- $x p_{k}(x, y)=p_{m}(x, y)$ with $m>k$
- $y p_{k}(x, y)=p_{m^{\prime}}(x, y)$ with $m^{\prime}>k$.

Two examples:



Define the inner product:

$$
\langle p, q\rangle=\sum_{i=1}^{n} w_{i}^{2} \overline{p\left(\xi_{i}, \eta_{i}\right)} q\left(\xi_{i}, \eta_{i}\right), \quad \xi_{i}, \eta_{i} \in \mathbb{R} \text { or } \mathbb{C}
$$

## Example

We choose the following ordering:

leading to the following "recurrence relations" for the OP $a_{i}(z)$ :

$$
\begin{aligned}
x\left[a_{1}(x, y), a_{2}(x, y), \ldots\right] & =\left[a_{1}(x, y), a_{2}(x, y), \ldots\right] H_{x} \\
y\left[a_{1}(x, y), a_{2}(x, y), \ldots\right] & =\left[a_{1}(x, y), a_{2}(x, y), \ldots\right] H_{y}
\end{aligned}
$$

with the following "pivot" structure for $H_{x}$ and $H_{y}$ :

$$
H_{x}=\left[\begin{array}{r}
\times \times \times \times \cdots \\
\boxtimes \times \times \times \cdots \\
\times \times \times \cdots \\
\boxtimes \times \times \cdots \\
\boxtimes \times \cdots \\
\times \cdots \\
\boxtimes \cdots
\end{array}\right] \quad H_{y}=\left[\begin{array}{r}
\times \times \times \times \cdots \\
\times \times \times \times \cdots \\
\boxtimes \times \times \times \cdots \\
\times \times \times \cdots \\
\boxtimes \times \times \cdots \\
\boxtimes \times \cdots \\
\times \cdots \\
\boxtimes \cdots
\end{array}\right] .
$$

Goal: Construct such a basis, generalizing one-var. algorithm Idea: one-var. case: recurrence relation from multiplication by $x$ :

$$
x\left[a_{0}, a_{1}, \ldots, a_{n-1}\right]=\left[a_{0}, a_{1}, \ldots, a_{n-1}\right] H
$$

two-var. case: multiplications by $x$ and $y$, separately:

$$
\begin{aligned}
x\left[a_{1}, a_{2}, \ldots, a_{k}, \ldots\right] & =\left[a_{1}, a_{2}, \ldots, a_{k}, \ldots\right] H_{x} \\
y\left[a_{1}, a_{2}, \ldots, a_{k}, \ldots\right] & =\left[a_{1}, a_{2}, \ldots, a_{k}, \ldots\right] H_{y}
\end{aligned}
$$

$H_{x}$ and $H_{y}$ - generalized Hessenberg.
Some choice is left: i.e. $x y^{2}=x \cdot y^{2}$ or $x y^{2}=y \cdot x y$.
Recurrence coefficients: can be taken from the $H_{x}$ or the $H_{y}$ matrix.

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## Inverse eigenvalue problem

Computing (the recurrence relation coefficients of) the OPs $a_{j}(x, y)$ can be done by solving the following IEP.
Definition (IEP - two-variable OP)
Given: $D_{x}=\operatorname{diag}\left(x_{i}\right), D_{y}=\operatorname{diag}\left(y_{i}\right)-$ points, $\mathbf{w}=\left(w_{i}\right)$ - weights Find: Unitary $Q$ and "generalized" upper Hessenberg matrices $H_{x}$ and $H_{y}$ such that

$$
Q^{H} \mathbf{w}=\|\mathbf{w}\| \mathbf{e}_{1}, \quad Q^{H} D_{x} Q=H_{x}, \quad Q^{H} D_{y} Q=H_{y}
$$

where the pivot structure of $H_{x}$ and $H_{y}$ determines the degree structure of the sequence of orthonormal polynomials.


3 points (2)


4 points (1)
 $\xrightarrow{G_{w}(1,4)}$

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## Two-variable

orthogoriable polynomials
Numerical examples
Orthogonal
polyanalytic


Orthogonal Lau
polynomials
Orthogonal rational
functions
Appendix

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Order (4 points, 4 polynomials)


4 points (3)


5 points (1)


$$
\xrightarrow{\sigma_{w}(1,5)}
$$

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Order (5 points, 5 polynomials)
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5 points (2)


5 points (3)


Order (6 points, 6 polynomials)


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5 points (4)
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$\xrightarrow{G^{\mathrm{X}}(4,5)}$

$\xrightarrow{6^{x}(4,5)}$


6 points (2)


## 6 points (4)



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6 points (3)
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## 6 points (5)



Two-variable
orthogonal polynomials
Numerical example
$\xrightarrow{\sigma^{\curlyvee}(3,6)}$



We can now recover the polynomials:

$$
\begin{gathered}
b_{1}=\text { const } \\
x b_{1}=\left[b_{1}, b_{2}\right] \cdot H_{x}(1: 2,1) \\
y b_{1}=\left[b_{1}, b_{2}, b_{3}\right] \cdot H_{y}(1: 3,1) \\
x b_{2}=\left[b_{1}, b_{2}, b_{3}, b_{4}\right] \cdot H_{x}(1: 4,2) \\
y b_{2}=\left[b_{1}, b_{2}, b_{3}, b_{4}, b_{5}\right] \cdot H_{y}(1: 5,2) \\
x b_{3}=\left[b_{1}, b_{2}, b_{3}, b_{4}, b_{5}\right] \cdot H_{x}(1: 5,3) \\
y b_{3}=\left[b_{1}, b_{2}, b_{3}, b_{4}, b_{5}, b_{6}\right] \cdot H_{y}(1: 6,3)
\end{gathered}
$$

## The Padua points as the inner product points

The $n=(\delta+1)(\delta+2) / 2$ Padua points of degree $\delta$ are

$$
\operatorname{Pad}_{\delta}=\left\{\zeta=\left(\zeta_{1}, \zeta_{2}\right)\right\}=\left\{\gamma\left(\frac{k \pi}{\delta(\delta+1)}\right), \quad k=0, \ldots, \delta(\delta+1)\right\}
$$

where $\gamma(t)$ is their "generating curve"

$$
\gamma(t)=(-\cos ((\delta+1) t),-\cos (\delta t)) \quad t \in[0, \pi]
$$




More points, more polynomials...

and so on

Example 1: Orthogonality test

Recall the IEP: Find unitary $Q$ (transformation matrix) and generalized upper Hessenberg $H_{x}$ and $H_{y}$ such that $Q^{H} D_{x} Q=H_{x}$ and $Q^{H} D_{y} Q=H_{y}$.
Let $A=\left[a_{1}\left(\zeta_{i}\right) a_{2}\left(\zeta_{i}\right) \ldots a_{n}\left(\zeta_{i}\right)\right]_{i=1}^{n}$ and $W=\operatorname{diag}\left(w_{i}\right)$. Then

$$
W A=Q
$$

Test whether WA is orthogonal

- Take $n=5151$ Padua points $\zeta_{i}$ of degree $\delta=100, W=I$.
- Compute $H_{x}, H_{y}$.
- Compute values of OP's at all $\zeta_{i}$ 's by recurrence relations stored in $H_{x}, H_{y}$. Multiply them with the corresponding weight $w_{i}$ and store the results (columnwise) in $V=W A$. Let $R=\left|V^{H} V-I\right|$.
- For $k=10: 100: n$ compute $\max R(1: k, 1: k)$.

Example 1: Orthogonality test


Figure: Max orthogonality error for the first $k$ OPs, $n=5151$ points

Example 2: Least squares test


Figure: Franke function

## Example 2: Least squares test

Recall that the solution $p(z)$ to the discrete LS problem is

$$
p(z)=\sum_{j=1}^{\alpha} c_{j} a_{j}(z) .
$$

Then $\mathbf{c}=\left(c_{j}\right)$ is given by $\mathbf{c}=A^{H} W^{H} W \mathbf{f}=\left[\left\langle a_{i}, \mathbf{f}\right\rangle\right]$, so we perform the same row operations on $W \mathrm{f}$ as on w .
The LS solution

- Consider the Franke test function $F(\zeta)$ on $[0,1] \times[0,1]$ and transform the $n=5151$ Padua points to fit $[0,1]^{2}$.
- Compute $\mathbf{f}=F\left(\zeta_{i}\right)$
- Compute $A^{H} W^{H} W \mathbf{f}=\mathbf{c}$.
- Plot $\left|\mathbf{c}_{k}\right|$ for all $k$.

Example 2: Least squares test


Figure: LS solution coefficients for Franke function, $n=5151$ points
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Example 2: Least squares test, relative error


Figure: Relative error, appr poly of length $1000, n=5151$ points

Example 3: Polynomial that goes through the points

Choice of the points

- Consider the square $[0,1] \times[0,1]$.
- 20 equidistant points on the circle with center $(0.25 ; ~ 0.25)$ and radius 0.15 .
- The next 20 points similarly on a circle with center ( $0.75 ; 0.75$ ).
- The last 4 points are the 4 corners of the square.

Find: the polynomial having "least degree" that has zero value in the given points.
Solution: look for the first zero pivot appearing in the recurrence relation.
This happens for the 28th orthogonal polynomial (of degree 6)

Example 2: Least squares test, relative error


Figure: Relative error, appr poly of length 3000, $n=5151$ points

Example 3: Polynomial that goes through the points


Figure: Surface plot of an interpolating polynomial

Orthogonal polyanalytic polynomials
Consider the monomials in the two variables $z$ and $\bar{z}$ as basis functions $p_{k}(z, \bar{z})=z^{i} \bar{z}^{j}$.
Choose an ordering of these basis functions such that:

- $z p_{k}(z, \bar{z})=p_{m}(z, \bar{z})$ with $m>k$
- $\bar{z} p_{k}(z, \bar{z})=p_{m^{\prime}}(z, \bar{z})$ with $m^{\prime}>k$.

Two examples:



Define the inner product:

$$
\langle p, q\rangle=\sum_{i=1}^{n} w_{i}^{2} \overline{p\left(z_{i}, \bar{z}_{i}\right)} q\left(z_{i}, \bar{z}_{i}\right), \quad z_{i} \in \mathbb{C}
$$

## Example

We choose the following ordering:


leading to the following "recurrence relations" for the OP $a_{i}(z)$ :

$$
\begin{aligned}
z\left[a_{1}(z, \bar{z}), a_{2}(z, \bar{z}), \ldots\right] & =\left[a_{1}(z, \bar{z}), a_{2}(z, \bar{z}), \ldots\right] H_{z} \\
\bar{z}\left[a_{1}(z, \bar{z}), a_{2}(z, \bar{z}), \ldots\right] & =\left[a_{1}(z, \bar{z}), a_{2}(z, \bar{z}), \ldots\right] H_{\bar{z}}
\end{aligned}
$$

with the following "pivot" structure for $H_{z}$ and $H_{\bar{z}}$ :

$$
H_{z}=\left[\begin{array}{r}
\times \times \times \times \cdots \\
\boxtimes \times \times \times \cdots \\
\times \times \times \cdots \\
\boxtimes \times \times \cdots \\
\boxtimes \times \cdots \\
\times \cdots \\
\boxtimes \cdots
\end{array}\right] \quad H_{\bar{z}}=\left[\begin{array}{r}
\times \times \times \times \cdots \\
\times \times \times \times \cdots \\
\boxtimes \times \times \times \cdots \\
\times \times \times \cdots \\
\boxtimes \times \times \cdots \\
\boxtimes \times \cdots \\
\times \cdots \\
\boxtimes \cdots
\end{array}\right]
$$

Recurrence relation

Goal: Construct such a basis, generalizing one-var. algorithm Idea: one-var. case: recurrence relation from multiplication by $z$ :

$$
z\left[a_{0}, a_{1}, \ldots, a_{n-1}\right]=\left[a_{0}, a_{1}, \ldots, a_{n-1}\right] H
$$

two-var. case: multiplications by $z$ and $\bar{z}$, separately:

$$
\begin{aligned}
& z\left[a_{1}, a_{2}, \ldots, a_{k}, \ldots\right]=\left[a_{1}, a_{2}, \ldots, a_{k}, \ldots\right] H_{z} \\
& \bar{z}\left[a_{1}, a_{2}, \ldots, a_{k}, \ldots\right]=\left[a_{1}, a_{2}, \ldots, a_{k}, \ldots\right] H_{\bar{z}}
\end{aligned}
$$

$H_{z}$ and $H_{\bar{z}}$ - generalized Hessenberg.
Some choice is left: i.e. $z \bar{z}^{2}=z \cdot \bar{z}^{2}$ or $z \bar{z}^{2}=\bar{z} \cdot z \bar{z}$.
Recurrence coefficients: can be taken from the $H_{z}$ or the $H_{\bar{z}}$ matrix.

## Inverse eigenvalue problem

Computing (the recurrence relation coefficients of) the OPs $a_{j}(z, \bar{z})$ can be done by solving the following IEP.
Definition (IEP - two-variable OP)
Given: $D_{z}=\operatorname{diag}\left(z_{i}\right)$ - points, $\mathbf{w}=\left(w_{i}\right)$ - weights
Find: Unitary $Q$ and "generalized" upper Hessenberg matrices $H_{z}$ and $H_{\bar{z}}$ such that

$$
Q^{H} \mathbf{w}=\|\mathbf{w}\| \mathbf{e}_{1}, \quad Q^{H} D_{z} Q=H_{z}, \quad Q^{H} D_{\bar{z}} Q=H_{\bar{z}}
$$

where the pivot structure of $H_{z}$ and $H_{\bar{z}}$ determines the degree structure of the sequence of orthonormal polynomials.

## Vector space of Laurent polynomials

Consider the vector space of Laurent polynomials $\mathcal{P}^{\alpha}=\left\{\sum_{i=1}^{\alpha} c_{i} p_{i}(z)\right\}$, with the sequence of basis functions $p_{i}(z)$ :

$$
1, z, z^{-1}, z^{2}, z^{-2}, z^{3}, z^{-3}, \ldots
$$

The leading index of a nonzero Laurent polynomial $p(z)=\sum_{i=1}^{\alpha} c_{i} p_{i}(z)$ is defined as

$$
\text { I-index }(p)=\max \left\{i \mid c_{i} \neq 0\right\}
$$

## Inverse eigenvalue problem

Computing (the recurrence relation coefficients of) the OLPs $a_{j}(z)$ can be done by solving the following IEP.

Definition (IEP - orthogonal Laurent polynomials)
Given: $Z=\operatorname{diag}\left(z_{i}\right)$ - points, $\mathbf{w}=\left(w_{i}\right)$ - weights
Find: Unitary $Q$ and "generalized" upper Hessenberg matrices $H_{z}$ and $H_{z^{-1}}$ such that

$$
Q^{H} \mathbf{w}=\|\mathbf{w}\| \mathbf{e}_{1}, \quad Q^{H} Z Q=H_{z}, \quad Q^{H} Z^{-1} Q=H_{z^{-1}}
$$

where the pivot structure of $H_{z}$ and $H_{z^{-1}}$ determines the degree structure of the sequence of orthonormal Laurent polynomials.

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## Recurrence relation

Let us consider the leading indices of $z a_{i}(z)$ and $z^{-1} a_{i}(z)$ :

$$
\begin{aligned}
\text { l-index }\left(z\left[a_{1}(z), a_{2}(z), a_{3}(z), \ldots\right]\right) & =[2,4, \leq 4,6, \leq 6, \ldots] \\
\text { I-index }\left(z^{-1}\left[a_{1}(z), a_{2}(z), a_{3}(z), \ldots\right]\right) & =[3, \leq 3,5, \leq 5,7, \ldots]
\end{aligned}
$$

This leads to the following "recurrence relations" for the OLP $a_{i}(z)$ :

$$
\begin{aligned}
z\left[a_{1}(z), a_{2}(z), \ldots\right] & =\left[a_{1}(z), a_{2}(z), \ldots\right] H_{z} \\
z^{-1}\left[a_{1}(z), a_{2}(z), \ldots\right] & =\left[a_{1}(z), a_{2}(z), \ldots\right] H_{z_{-1}}
\end{aligned}
$$

with the following "pivot" structure for $H_{z}$ and $H_{z^{-1}}$ :

$$
H_{z}=\left[\begin{array}{r}
\times \times \times \times \cdots \\
\boxtimes \times \times \times \cdots \\
\times \times \times \cdots \\
\boxtimes \times \times \cdots \\
\times \cdots \\
\boxtimes \cdots \\
\vdots
\end{array}\right] \quad H_{z^{-1}}=\left[\begin{array}{r}
\times \times \times \times \cdots \\
\times \times \times \times \cdots \\
\boxtimes \times \times \times \cdots \\
\times \times \cdots \\
\boxtimes \times \cdots \\
\vdots \\
\end{array}\right]
$$

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## Vector space

Given the complex numbers $y_{1}, y_{2}, \ldots, y_{n}$ all different from each other. Let us consider the vector space $\mathcal{P}^{n}$ of all proper rational functions having possible poles in $y_{1}, y_{2}, \ldots, y_{n}$ :

$$
\mathcal{P}^{n}=\operatorname{span}\left\{1, \frac{1}{z-y_{1}}, \frac{1}{z-y_{2}}, \ldots, \frac{1}{z-y_{n}}\right\}
$$

Bilinear form

Given the complex numbers $z_{0}, z_{1}, \ldots, z_{n}$ which together with the numbers $y_{i}$ are all different from each other, and the "weights" $0 \leq w_{i}, i=0,1, \ldots, n$, we define the following bilinear form

$$
\langle p, q\rangle=\sum_{i=0}^{n} w_{i}^{2} \overline{p\left(z_{i}\right)} q\left(z_{i}\right) .
$$

This bilinear form defines an inner product in the space $\mathcal{P}^{n}$.

## Recurrence relation

The recurrence relation for the orthonormal rational functions $a_{j}(z)$ for $j=0,1, \ldots, n$ can be written as:

$$
\mathbf{a}(z)\left(z l-D_{y}\right) H=\mathbf{a}(z)+c a_{n+1}(z) \mathbf{e}_{n}
$$

with $a_{n+1}(z)=\frac{\prod_{j=0}^{n}\left(z-z_{j}\right)}{\prod_{j=1}^{n}\left(z-y_{j}\right)}$ and
$D_{y}=\operatorname{diag}\left(y_{0}, y_{1}, \ldots, y_{n}\right)$ with $y_{0}$ chosen arbitrarily. Multiply to the right by the inverse of the upper Hessenberg matrix $H: S=H^{-1}$ :

$$
z \mathbf{a}(z)=\mathbf{a}(z)\left(S+D_{y}\right)+c a_{n+1}(z) \mathbf{s}_{n}
$$

with $s_{n}$ the last row of the matrix $S$.
Note that the matrix $S$ is lower semiseparable, i.e., has rank 1 structure in the lower triangular part: tril $(S)=$ tril(rank 1 matrix).

Let us consider an orthonormal basis

$$
\mathbf{a}_{n}=\left[a_{0}, a_{1}, \ldots, a_{n}\right]
$$

for $\mathcal{P}^{n}$ satisfying the following properties

$$
\begin{aligned}
a_{j} \in \mathcal{P}^{j} \backslash \mathcal{P}^{j-1} & \left(\mathcal{P}^{-1}=\emptyset\right) \\
\left\langle a_{i}, a_{j}\right\rangle=\delta_{i, j} & (\text { Kronecker delta })
\end{aligned}
$$

for $i, j=0,1,2, \ldots, n$.

## Connecting the ORF $a_{j}(z)$ to the columns of $Q$

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Orthogonal rational functions

## Recurrence relation for $a_{j}(z)$ :

$$
z \mathbf{a}(z)=\mathbf{a}(z)\left(S+D_{y}\right)+c a_{n+1}(z) \mathbf{s}_{n}
$$

with $a_{n+1}(z)=\frac{\prod_{j=0}^{n}\left(z-z_{j}\right)}{\prod_{j=1}^{n}\left(z-y_{j}\right)}$.
Because

$$
Q_{j}=\left[w_{i} a_{j}\left(z_{i}\right)\right]_{i=0,1, \ldots, n}
$$

we derive the following relation for the unitary matrix $Q$ :

$$
D_{z} Q=Q\left(S+D_{y}\right), \quad \text { or } \quad Q^{H} D_{z} Q=S+D_{y}
$$

Later on we will design an algorithm to compute $Q$ and $S$. Now we will look at a recurrence relation between the columns $Q_{j}$ of $Q$. Then we will give the connection between the columns $Q_{j}$ and the values of rational functions satisfying a similar recurrence relation. Finally, we will show that these rational functions form a basis we are looking for.
notation: $H=S^{-1}$ is upper Hessenberg with subdiagonal elements $b_{0}, b_{1}, \ldots, b_{n-1}$. The $j$ th column $H_{j}$ of $H$ has the form

$$
H_{j}^{T}=:\left[\vec{h}_{j}^{T}, b_{j}, \overrightarrow{0}^{\top}\right]
$$

## Recurrence relation for $Q_{j}$

Let $Q=:\left[Q_{0}, Q_{1}, \ldots, Q_{n}\right]$.
The columns $Q_{j}$ satisfy the following recurrence relation

$$
\begin{aligned}
& \left(D_{z}-y_{j+1} I\right) b_{j} Q_{j+1}=Q_{j} \\
& \quad+\left(\left[Q_{0}, Q_{1}, \ldots, Q_{j}\right] D_{y, j}-D_{z}\left[Q_{0}, Q_{1}, \ldots, Q_{j}\right]\right) \vec{h}_{j} \\
& \quad j=0,1, \ldots, n-1
\end{aligned}
$$

with

$$
Q_{0}=\frac{\mathbf{w}}{\|\mathbf{w}\|}
$$

## Recurrence relation

Looking at the recurrence relation for $Q_{j}$

$$
\begin{aligned}
& \left(D_{z}-y_{j+1} I\right) b_{j} Q_{j+1}=Q_{j} \\
& +\left(\left[Q_{0}, Q_{1}, \ldots, Q_{j}\right] D_{y, j}-D_{z}\left[Q_{0}, Q_{1}, \ldots, Q_{j}\right]\right) \vec{h}_{j} \\
& \quad j=0,1, \ldots, n-1
\end{aligned}
$$

we can compute an orthonormal basis $\left[a_{0}, a_{1}, \ldots, a_{n}\right]$ for $\mathcal{P}^{n}$ using a similar recurrence relation

$$
\begin{aligned}
& \quad a_{j+1}(z)= \\
& \frac{a_{j}(z)+\left(\left[a_{0}, a_{1}, \ldots, a_{j}\right] D_{y, j}-z\left[a_{0}, a_{1}, \ldots, a_{j}\right]\right) \vec{h}_{j}}{\left(z-y_{j+1}\right) b_{j}}
\end{aligned}
$$

for $j=0,1, \ldots, n-1$ and
with $a_{0}(x)=1 / \sqrt{\sum\left|w_{i}\right|^{2}}$.

## Proof

Filling in $z_{i}$ for $z$ in the recurrence relation for $a_{j+1}(z)$, we get

$$
\begin{aligned}
& \left(D_{z}-y_{j+1} I\right) b_{j}\left[a_{j+1}\left(z_{i}\right)\right]=a_{j}\left(z_{i}\right) \\
& \quad+\left[a_{0}\left(z_{i}\right), a_{1}\left(z_{i}\right), \ldots, a_{j}\left(z_{i}\right)\right] D_{y, j} \vec{h}_{j} \\
& \quad-D_{z}\left[a_{0}\left(z_{i}\right), a_{1}\left(z_{i}\right), \ldots, a_{j}\left(z_{i}\right)\right] \vec{h}_{j}
\end{aligned}
$$

Because $Q_{0}=\operatorname{diag}(\mathbf{w})\left[a_{0}\left(z_{i}\right)\right]$ and because $\operatorname{diag}(\mathbf{w})$ is diagonal as well as all the other square matrices involved, first part of the theorem is proved.

For $j=0,1, \ldots, n$ we have that

$$
\begin{aligned}
& Q_{j}=\operatorname{diag}(\mathbf{w})\left[a_{j}\left(z_{i}\right)\right] \\
& a_{j} \in \mathcal{P}^{j} \backslash \mathcal{P}^{j-1}
\end{aligned}
$$

## Proof (continued)

We have to prove that $a_{j} \in \mathcal{P}^{j} \backslash \mathcal{P}^{j-1}$.
This is clearly true for $j=0$.
Suppose it is true for $j=0,1,2, \ldots, k<n$.
From the recurrence relation, we derive that $a_{k+1}(z)$ has the form

$$
\begin{aligned}
& a_{k+1}(z)= \\
& \frac{\text { rat. function with possible poles in } y_{0}, y_{1}, \ldots, y_{k}}{\left(z-y_{j+1}\right)} .
\end{aligned}
$$

Also $\lim _{z \rightarrow \infty} a_{k+1}(z) \in \mathbb{C}$.
Hence, $a_{j} \in \mathcal{P}^{j} \backslash \mathcal{P}^{j-1}$.

Orthonormality of $\vec{a}_{n}$
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The functions $\vec{a}_{n}=\left[a_{0}, a_{1}, \ldots, a_{n}\right]$ form an orthonormal basis for $P^{n}$ with respect to the inner product $\langle\cdot, \cdot\rangle$. Moreover, $z^{i} \in \mathcal{P}^{i} \backslash \mathcal{P}^{i-1}$.
Proof The only thing that remains to be proven is $\left\langle a_{i}, a_{j}\right\rangle=\delta_{i, j}$. This follows immediately from the fact that $Q=\operatorname{diag}(\mathbf{w})\left[a_{j}\left(z_{i}\right)\right]$ and $Q$ is unitary.

## Proof

The recurrence relation for the $a_{j}, j=0,1, \ldots, n$ can also be written as

$$
\mathbf{a}(z)\left(z l-D_{y}\right) H=\mathbf{a}(z)+c a_{n+1}(z) \mathbf{e}_{n} .
$$

Multiplying to the right by $S=\mathrm{H}^{-1}$, we derive the recurrence formula.
To determine $a_{n+1}$ we look at the last column of the previous relation. It follows that $a_{n+1}$ is a rational function having degree of numerator at most one more than the degree of the denominator and having possible poles in $y_{1}, y_{2}, \ldots, y_{n}$.

Recurrence relation

$$
z \mathbf{a}(z)=\mathbf{a}(z)\left(S+D_{y}\right)+c a_{n+1}(z) \mathbf{s}_{n}
$$

with $s_{n}$ the last row of the semiseparable matrix $S$ and $a_{n+1}(z)=\frac{\prod_{j=0}^{n}\left(z-z_{j}\right)}{\prod_{j=1}^{n}\left(z-y_{j}\right)}$.

## Proof (continued)

Let us evaluate the previous equation in the points $z_{i}$

$$
D_{z}\left[a_{j}\left(z_{i}\right)\right] H-\left[a_{j}\left(z_{i}\right)\right] D_{y} H=\left[a_{j}\left(z_{i}\right)\right]+c\left[a_{n+1}\left(z_{i}\right)\right] \mathbf{e}_{n}
$$

Multiplying to the left by $\operatorname{diag}(\mathbf{w})=D_{w}$ and because $D_{w} D_{z}=D_{z} D_{w}$, we obtain

$$
D_{z} Q H-Q D_{y} H=Q+c D_{w}\left[a_{n+1}\left(z_{i}\right)\right] \mathbf{e}_{n}
$$

From

$$
D_{z} Q H=Q\left(I+D_{y} H\right)
$$

it follows that $a_{n+1}$ has zeros in $z_{i}, i=0,1, \ldots, n$ and this proves the theorem.

- if all $z_{i}$ and $y_{i}$ are real: $Q^{\top} D_{z} Q=S+D_{y}$ Hence, $S$ is symmetric.
- if all $z_{i}$ are real but $y_{i}$ can be complex then the strictly upper triangular part $R$ is also of rank 1 .
- if $z_{i}$ are all on the unit circle, then the strictly upper triangular part $R$ is also of rank 1 .
In all these cases the computational complexity reduces to $\mathcal{O}\left(n^{2}\right)$.


## Connection between $Q, H$ and the OPs

Consider the following recurrence relation for $b_{i}$ :

$$
\begin{equation*}
b_{0}=\frac{1}{\|\mathbf{w}\|}, \quad z\left[b_{0}, b_{1}, \ldots, b_{n-1}\right]=\left[b_{0}, b_{1}, \ldots, b_{n-1}\right] H \tag{2}
\end{equation*}
$$

$H$ Hessenberg $\Rightarrow$ derive $b_{1}$ from 1st col, $b_{2}$ from 2 nd, $\ldots$

$$
\begin{gathered}
Q^{H} \mathbf{w}=\|\mathbf{w}\| \mathbf{e}_{1} \quad \Rightarrow \quad Q \mathbf{e}_{1}=\frac{\mathbf{w}}{\|\mathbf{w}\|}=\operatorname{diag}(\mathbf{w})\left[b_{0}\left(z_{i}\right)\right] \\
D_{z} Q=Q H \Rightarrow Q \mathbf{e}_{k}=\operatorname{diag}(\mathbf{w})\left[b_{k-1}\left(z_{i}\right)\right], k=1,2, \ldots, n .
\end{gathered}
$$

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## Vector case of discrete LS approximation

Discrete vector LS approximation
Given: $L$ functions $f_{j}(z), z \in \mathcal{D} \subset \mathbb{C}, N$ points $z_{i} \in \mathcal{D} \subset \mathbb{C}$ and corresponding weights $w_{i}$
Find: $L$ polynomials $p_{j}(z): \operatorname{deg} p_{j} \leqslant \alpha_{j}$ and a normalization s.t.

$$
\left.\sum_{i=1}^{N} w_{i}^{2}\left|\left[f_{1} f_{2} \cdots f_{L}\right]\left[\begin{array}{c}
p_{1} \\
\vdots \\
p_{L}
\end{array}\right]\right|_{z=z_{i}}\right|^{2} \quad \text { is minimal. }
$$

Special case: discrete LS approximation
Find: the polynomial $p(z)$ of degree $\leqslant \alpha$ s.t.
$\sum_{i=1}^{N} w_{i}^{2}\left|f\left(x_{i}\right)-p\left(z_{i}\right)\right|^{2}=\sum_{i=1}^{N} w_{i}^{2}\left|\left[\begin{array}{ll}f\left(z_{i}\right) & -1\end{array}\right]\left[\begin{array}{c}1 \\ p\left(z_{i}\right)\end{array}\right]\right|^{2} \quad$ is min.
[VB, Bultheel]

